



# I/O Management



# Goals of this Lecture

Help you to learn about:

- The C/Unix **file** abstraction
- Standard C I/O
  - Data structures & functions
- Unix I/O
  - Data structures & functions
- The implementation of Standard C I/O using Unix I/O
- Programmatic redirection of stdin, stdout, and stderr
- Pipes



# Agenda

## The C/Unix file abstraction

Unix I/O system calls

C's Standard IO library (FILE \*)

Implementing standard C I/O using Unix I/O

Redirecting standard files

Pipes



# C/Unix File Abstraction

## Problem:

- At the physical level...
- Code that **reads** from **keyboard** is very different from code that reads from **disk**, etc.
- Code that **writes** to **video screen** is very different from code that writes to **disk**, etc.
- Would be nice if application programmer didn't need to worry about such details

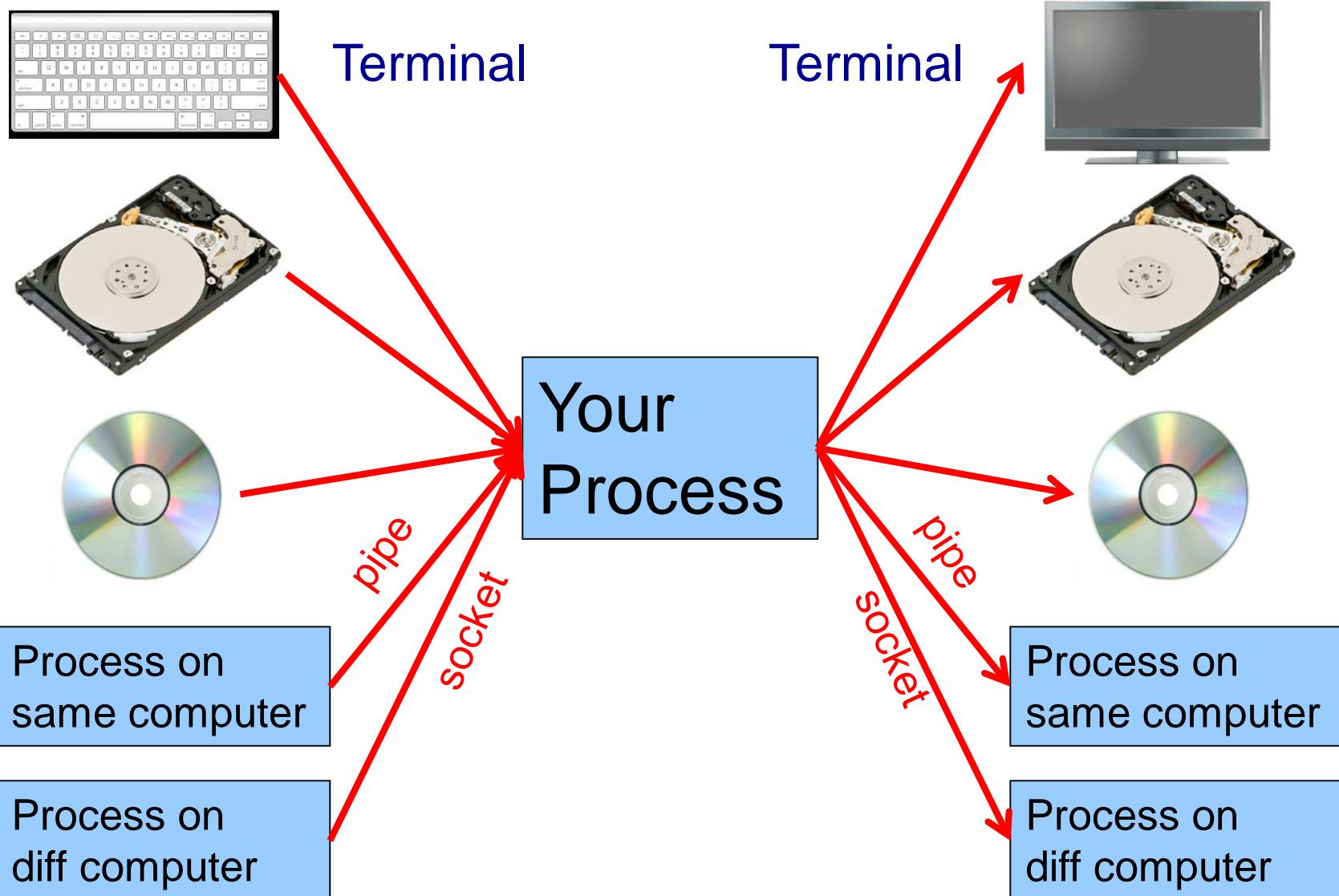
## Solution:

- **File**: a sequence of bytes
- C and Unix allow application program to treat any data source/destination as a **file**

Commentary: **Beautiful** abstraction!



# Data Sources and Destinations

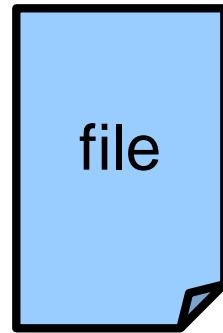




# C/Unix File Abstraction

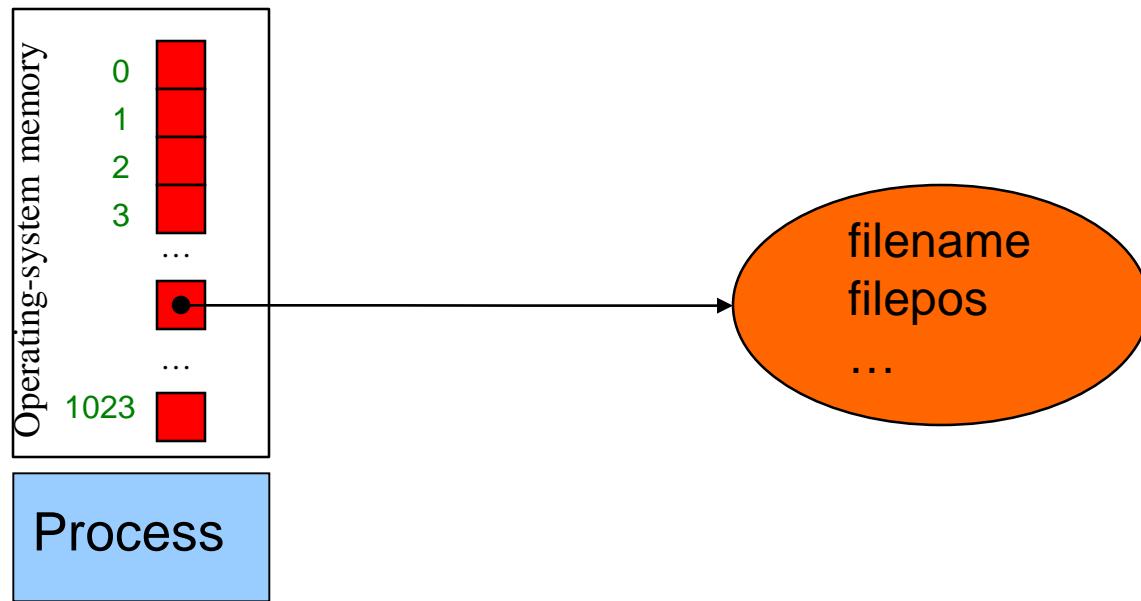
Each file has an associated **file position**

- Starts at beginning of file (if opened to read or write)
- Starts at end of file (if opened to append)





# Unix I/O Data Structures



**File descriptor:** Integer that uniquely identifies an open file

**File descriptor table:** an array

Indices are file descriptors; elements are pointers to file tables

One unique file descriptor table for each process

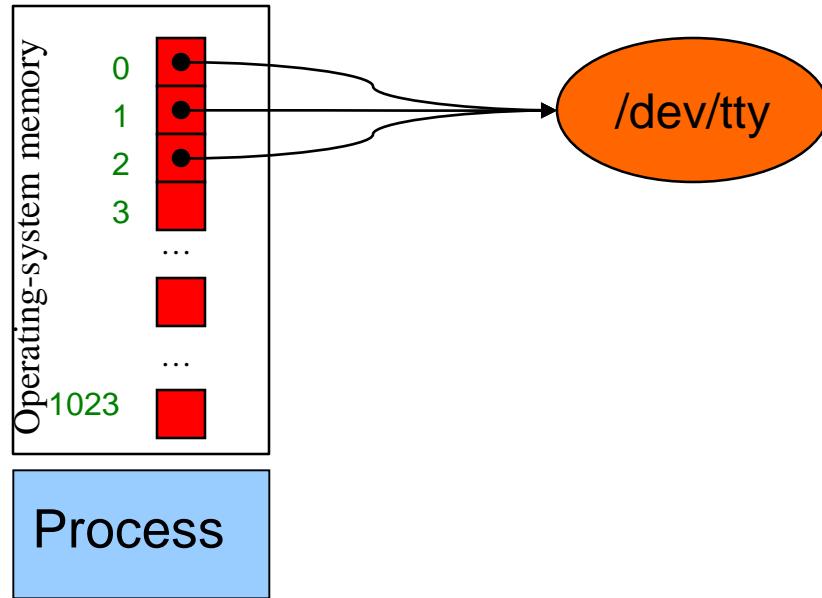
**File table:** a structure

In-memory surrogate for an open file

Created when process opens file; maintains file position



# Unix I/O Data Structures



At process start-up files with fd 0, 1, 2 are open automatically  
(By default) each references file table for a file named /dev/tty

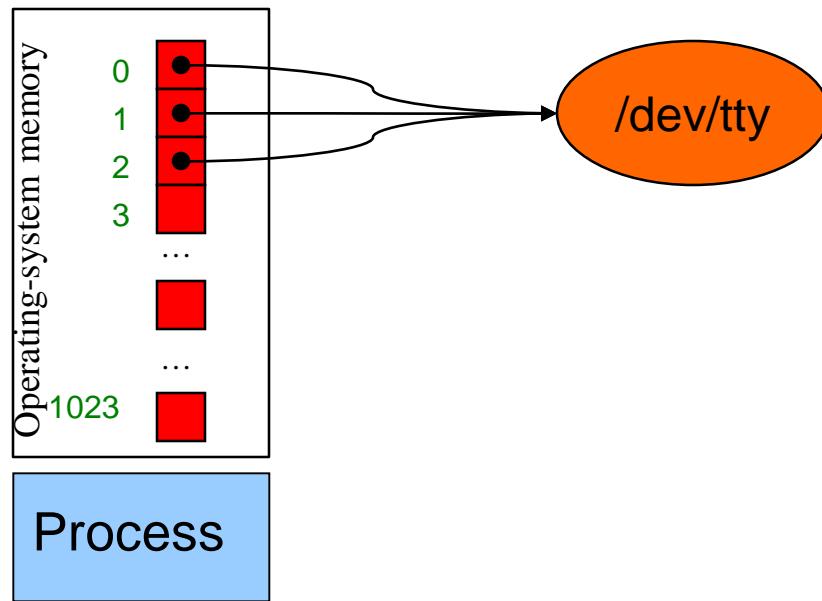
**/dev/tty**

In-memory surrogate for the terminal  
**Terminal**

Combination keyboard/video screen



# Unix I/O Data Structures



Read from stdin  $\Rightarrow$  read from fd 0

Write to stdout  $\Rightarrow$  write to fd 1

Write to stderr  $\Rightarrow$  write to fd 2



# System-Level Functions

As noted in the ***Exceptions and Processes*** lecture...

Linux system-level functions for **I/O management**

Function	Description
read()	Read data from file descriptor Called by getchar(), scanf(), etc.
write()	Write data to file descriptor Called by putchar(), printf(), etc.
open()	Open file or device Called by fopen(..., "r")
close()	Close file descriptor Called by fclose()
creat()	Open file or device for writing Called by fopen(..., "w")
lseek()	Change file position Called by fseek()



# System-Level Functions

As noted in the ***Exceptions and Processes*** lecture..

Linux system-level functions for **I/O redirection** and **inter-process communication**

Function	Description
dup()	Duplicate an open file descriptor
pipe()	Create a channel of communication between processes



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Implementing standard C I/O using Unix I/O

Redirecting standard files

Pipes



# Unix I/O Functions

```
int creat(char *filename, mode_t mode);
```

- Create a new empty file named `filename`
  - `mode` indicates permissions of new file
- Implementation:
  - Create new empty file on disk
  - Create file table
  - Set first unused file descriptor to point to file table
  - Return file descriptor used, -1 upon failure



# Unix I/O Functions

```
int open(char *filename, int flags, ...);
```

- Open the file whose name is `filename`
  - `flags` often is `O_RDONLY`
- Implementation (assuming `O_RDONLY`):
  - Find existing file on disk
  - Create file table
  - Set first unused file descriptor to point to file table
  - Return file descriptor used, -1 upon failure



# Unix I/O Functions

`int close(int fd);`

- Close the file `fd`
- Implementation:
  - Destroy file table referenced by element `fd` of file descriptor table
    - As long as no other process is pointing to it!
  - Set element `fd` of file descriptor table to `NULL`



# Unix I/O Functions

```
int read(int fd, void *buf, int count);
```

- Read into `buf` up to `count` bytes from file `fd`
- Return the number of bytes read; 0 indicates end-of-file

```
int write(int fd, void *buf, int count);
```

- Writes up to `count` bytes from `buf` to file `fd`
- Return the number of bytes written; -1 indicates error

```
int lseek(int fd, int offset, int whence);
```

- Set the file position of file `fd` to file position `offset`. `whence` indicates if the file position is measured from the beginning of the file (`SEEK_SET`), from the current file position (`SEEK_CUR`), or from the end of the file (`SEEK_END`)
- Return the file position from the beginning of the file



# Unix I/O Functions

## Note

- Only 6 system-level functions support all I/O from all kinds of devices!

Commentary: **Beautiful** interface!



# Unix I/O Example 0

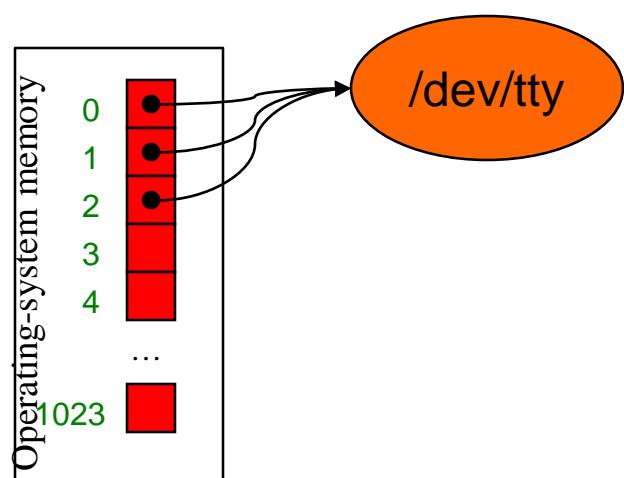
## Proto-getchar()

```
#include <string.h>
#include <unistd.h>

int proto_getchar(void)
{   char buf[1];
    int n;

    n = read(0, buf, 1);
    if (n==1)
        return buf[0];
    else return EOF;
}
```

# of bytes to (try to) read  
0 is the file descriptor  
of the standard input



and the problem is . . . too slow.

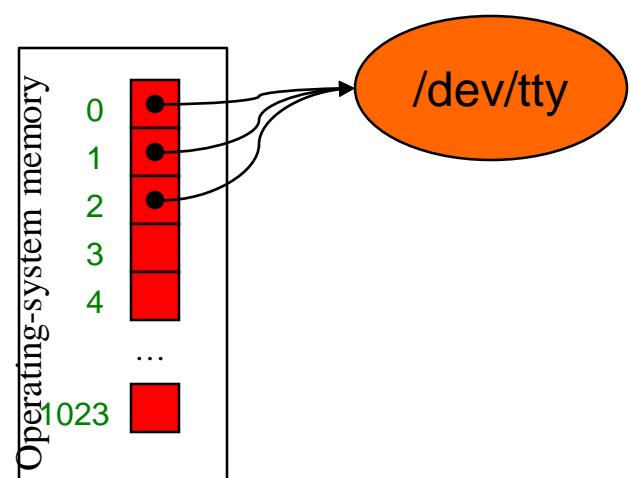
Does a system call for every character.



# Unix I/O Example 1

Write “hello, world\n” to /dev/tty

```
#include <string.h>
#include <unistd.h>
int main(void)
{
    char hi[] = "hello, world\n";
    size_t countWritten = 0;
    size_t countToWrite = strlen(hi);
    while (countWritten < countToWrite)
        countWritten +=
            write(1, hi + countWritten,
                  countToWrite - countWritten);
    return 0;
}
```



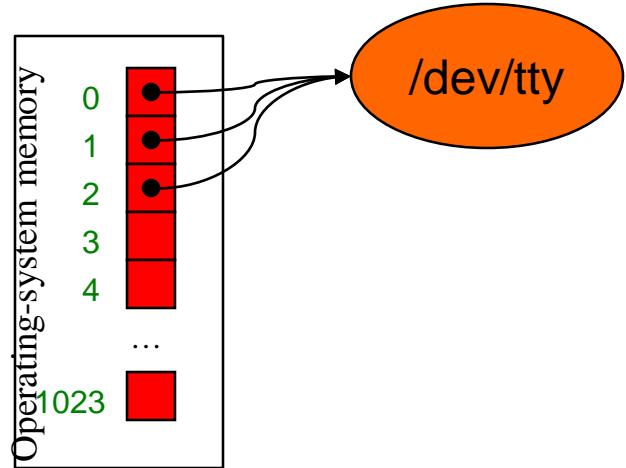
To save space,  
no error handling  
code is shown



# Unix I/O Example 2

```
#include <fcntl.h>
#include <unistd.h>
int main(void)
{ enum {BUFFERSIZE = 10};
  int fdIn, fdOut;
  int countRead, countWritten;
  char buf[BUFFERSIZE];
  fdIn = open("infile", O_RDONLY);
  fdOut = creat("outfile", 0600);
  for (;;)
  { countRead =
      read(fdIn, buf, BUFFERSIZE);
    if (countRead == 0) break;
    countWritten = 0;
    while (countWritten < countRead)
      countWritten +=
        write(fdOut,
              buf + countWritten,
              countRead - countWritten);
  }
  close(fdOut);
  close(fdIn);
  return 0;
}
```

Copy all bytes  
from infile to outfile



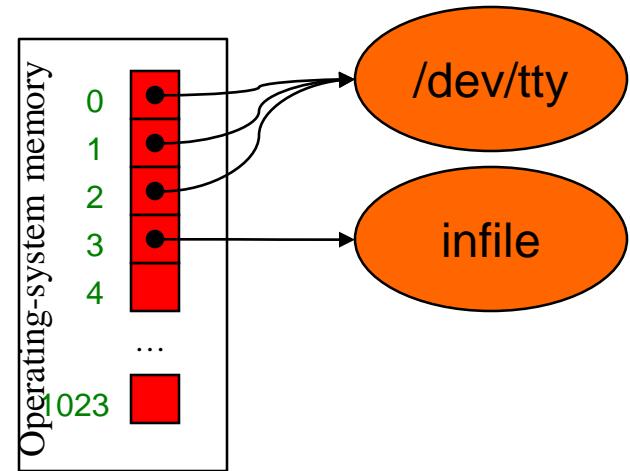
To save space,  
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# Unix I/O Example 2

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#include <unistd.h>
int main(void)
{ enum {BUFFERSIZE = 10};
  int fdIn, fdOut;
  int countRead, countWritten;
  char buf[BUFFERSIZE];
  fdIn<= open("infile", O_RDONLY);
  fdOut = creat("outfile", 0600);
  for (;;)
  { countRead =
      read(fdIn, buf, BUFFERSIZE);
    if (countRead == 0) break;
    countWritten = 0;
    while (countWritten < countRead)
      countWritten +=
        write(fdOut,
              buf + countWritten,
              countRead - countWritten);
  }
  close(fdOut);
  close(fdIn);
  return 0;
}
```

3

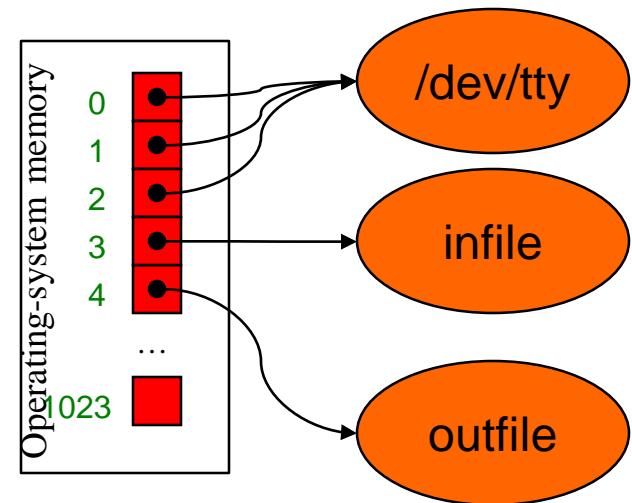




# Unix I/O Example 2

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int main(void)
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  int fdIn, fdOut;
  int countRead, countWritten;
  char buf[BUFFERSIZE];
  fdIn<= open("infile", O_RDONLY);
  fdOut<- creat("outfile", 0600);
  for (;;)
  { countRead =
      read(fdIn, buf, BUFFERSIZE);
    if (countRead == 0) break;
    countWritten = 0;
    while (countWritten < countRead)
      countWritten +=
        write(fdOut,
              buf + countWritten,
              countRead - countWritten);
  }
  close(fdOut);
  close(fdIn);
  return 0;
}
```

3  
4

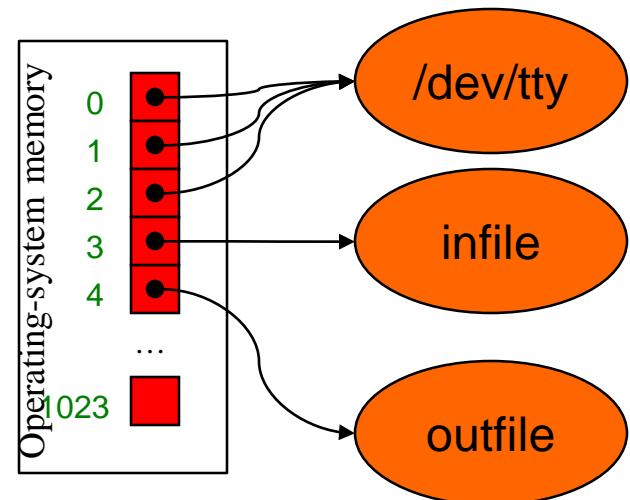




# Unix I/O Example 2

```
#include <fcntl.h>
#include <unistd.h>
int main(void)
{ enum {BUFFERSIZE = 10};
  int fdIn, fdOut;
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  char buf[BUFFERSIZE];
  fdIn<= open("infile", O_RDONLY);
  fdOut<= creat("outfile", 0600);
  for (;;)
  { countRead =
      read(fdIn, buf, BUFFERSIZE);
    if (countRead == 0) break;
    countWritten = 0;
    while (countWritten < countRead)
      countWritten +=
        write(fdOut,
              buf + countWritten,
              countRead - countWritten);
  }
  close(fdOut);
  close(fdIn);
  return 0;
}
```

3  
4

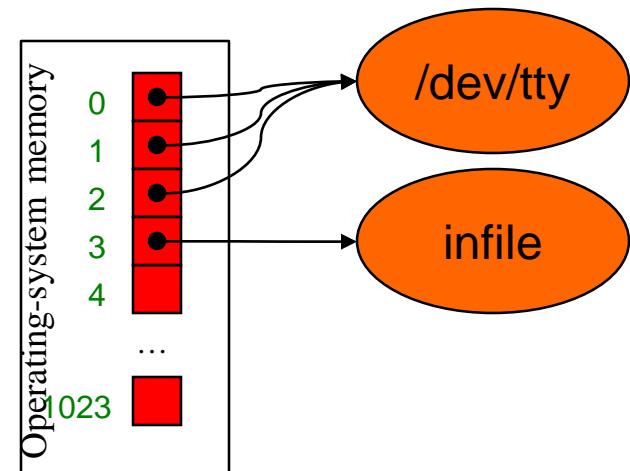




# Unix I/O Example 2

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#include <fcntl.h>
#include <unistd.h>
int main(void)
{ enum {BUFFERSIZE = 10};
  int fdIn, fdOut;
  int countRead, countWritten;
  char buf[BUFFERSIZE];
  fdIn<= open("infile", O_RDONLY);
  fdOut<= creat("outfile", 0600);
  for (;;)
  { countRead =
      read(fdIn, buf, BUFFERSIZE);
    if (countRead == 0) break;
    countWritten = 0;
    while (countWritten < countRead)
      countWritten +=
        write(fdOut,
              buf + countWritten,
              countRead - countWritten);
  }
  close(fdOut);
  close(fdIn);
  return 0;
}
```

3  
4

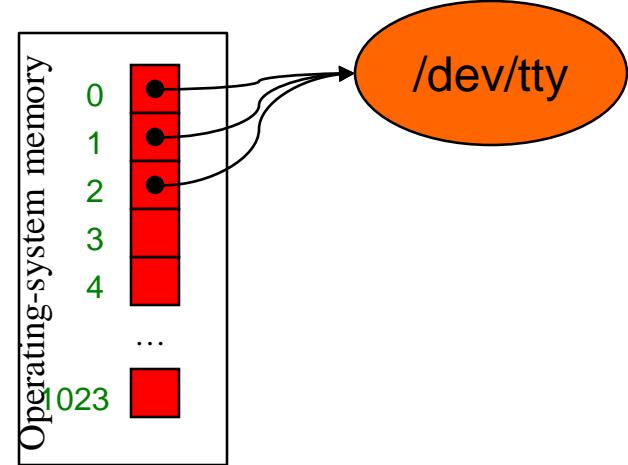




# Unix I/O Example 2

```
#include <fcntl.h>
#include <unistd.h>
int main(void)
{ enum {BUFFERSIZE = 10};
  int fdIn, fdOut;
  int countRead, countWritten;
  char buf[BUFFERSIZE];
  fdIn<= open("infile", O_RDONLY);
  fdOut<= creat("outfile", 0600);
  for (;;)
  { countRead =
      read(fdIn, buf, BUFFERSIZE);
    if (countRead == 0) break;
    countWritten = 0;
    while (countWritten < countRead)
      countWritten +=
        write(fdOut,
              buf + countWritten,
              countRead - countWritten);
  }
  close(fdOut);
  close(fdIn);
  return 0;
}
```

3  
4





# Agenda

The C/Unix file abstraction

Unix I/O system calls

**C's Standard IO library (FILE \*)**

Implementing standard C I/O using Unix I/O

Redirecting standard files

Pipes



# Standard C I/O Data Structure

We want 1-character-at-a-time I/O (`getc()`, `putc()`)

We want a-few-characters-at-a-time I/O (`scanf`, `printf`)

We *could* do this with `read()` and `write()` system calls,

BUT IT WOULD BE TOO SLOW to do 1 syscall per byte

Solution: Buffered input/output as an Abstract Data Type

## The **FILE** ADT

- A **FILE** object is an in-memory surrogate for an opened file
  - Created by `fopen()`
  - Destroyed by `fclose()`
  - Used by reading/writing functions



# Standard C I/O Functions

Some of the most popular:

`FILE *fopen(const char *filename, const char *mode);`

- Open the file named `filename` for reading or writing
- `mode` indicates data flow direction
  - “r” means read; “w” means write, “a” means append)
- Creates `FILE` structure
- Returns address of `FILE` structure

`int fclose(FILE *file);`

- Close the file identified by `file`
- Destroys `FILE` structure whose address is `file`
- Returns 0 on success, EOF on failure



# Standard C Input Functions

Some of the most popular:

`int fgetc(FILE *file);`

- Read a char from the file identified by `file`
- Return the char on success, `EOF` on failure

`int getchar(void);`

- Same as `fgetc(stdin)`

`char *fgets(char *s, int n, FILE *file);`

- Read at most `n` characters from `file` into array `s`
- Returns `s` on success, `NULL` on failure

`char *gets(char *s);`

- Essentially same as `fgets(s, INT_MAX, stdin)`
- *Buffer overflow waiting to happen*



# Standard C Input Functions

Some of the most popular:

```
int fscanf(FILE *file, const char *format, ...);
```

- Read chars from the file identified by `file`
- Convert to values, as directed by `format`
- Copy values to memory
- Return count of values successfully scanned

```
int scanf(const char *format, ...);
```

- Same as `fscanf(stdin, format, ...)`



# Standard C Output Functions

Some of the most popular:

```
int fputc(int c, FILE *file);
```

- Write **c** (converted to a char) to file
- Return **c** on success, **EOF** on failure

```
int putchar(int c);
```

- Same as **fputc(c, stdout)**

```
int fputs(const char *s, FILE *file);
```

- Write string **s** to **file**
- Return non-negative on success, **EOF** on error

```
int puts(const char *s);
```

- Essentially same as **fputs(s, stdout)**



# Standard C Output Functions

Some of the most popular:

```
int fprintf(FILE *file, const char *format, ...);
```

- Write chars to the file identified by `file`
- Convert values to chars, as directed by `format`
- Return count of chars successfully written
- Works by calling `fputc()` repeatedly

```
int printf(const char *format, ...);
```

- Same as `fprintf(stdout, format, ...)`



# Standard C I/O Functions

Some of the most popular:

`int fflush(FILE *file);`

- On an output file: write any buffered chars to `file`
- On an input file: behavior undefined
- `file == NULL` ⇒ flush buffers of **all** open files

`int fseek(FILE *file, long offset, int origin);`

- Set the file position of `file`
- Subsequent read/write accesses data starting at that position
- Origin: `SEEK_SET`, `SEEK_CUR`, `SEEK_END`

`int ftell(FILE *file);`

- Return file position of `file` on success, -1 on error



# Standard C I/O Example 1

Write “hello, world\n” to **stdout**

```
#include <stdio.h>
int main(void)
{   char hi[] = "hello world\n";
    size_t i = 0;
    while (hi[i] != '\0')
    {   putchar(hi[i]);
        i++;
    }
    return 0;
}
```

Simple  
Portable  
Efficient (via buffering)

```
#include <stdio.h>
int main(void)
{   puts("hello, world");
    return 0;
}
```

```
#include <stdio.h>
int main(void)
{   printf("hello, world\n");
    return 0;
}
```



# Standard C I/O Example 2

Copy all bytes from infile to outfile

```
#include <stdio.h>
int main(void)
{  int c;
   FILE *inFile;
   FILE *outFile;
   inFile = fopen("infile", "r");
   outFile = fopen("outfile", "w");
   while ((c = fgetc(inFile)) != EOF)
      fputc(c, outFile);
   fclose(outFile);
   fclose(inFile);
   return 0;
}
```

Simple  
Portable  
Efficient (via buffering)



# Standard C Buffering

Question: Exactly when are buffers flushed?

Answers:

If reading from a file

- (1) When buffer is empty



# Standard C Buffering

Question: Exactly when are buffers flushed?

Answers:

If writing to an ordinary file

- (1) File's buffer becomes full
- (2) Process calls `fflush()` on that file
- (3) Process terminates normally

If writing to `stdout` (in addition to previous)

- (4) `stdout` is bound to terminal and '`\n`' is appended to buffer
- (5) `stdin` and `stdout` are bound to terminal  
*and* read from `stdin` occurs

If writing to `stderr`

- Irrelevant; `stderr` is unbuffered



# Standard C Buffering Example

```
#include <stdio.h>
int main(void)
{  int dividend, divisor, quotient;

    printf("Dividend: ");
    scanf("%d", &dividend); ← Buffer flushed

    printf("Divisor: ");
    scanf("%d", &divisor); ← Buffer flushed

    printf("The quotient is ");
    quotient = dividend / divisor;
    printf("%d\n", quotient); ← Buffer flushed

    return 0;
}
```

```
$ pgm
Dividend: 6
Divisor: 2
The quotient is 3
$
```

```
$ pgm
Dividend: 6
Divisor: 0
Floating point exception
$
```

Output buffered  
Buffer flushed

Output buffered  
Buffer flushed

Output buffered  
Buffer flushed



# Agenda

The C/Unix file abstraction

Unix I/O system calls

C's Standard IO library (FILE \*)

## Implementing standard C I/O using Unix I/O

Redirecting standard files

Pipes



# Standard C I/O

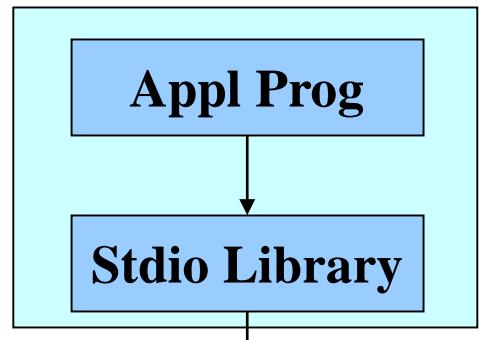
## Question:

- How to implement standard C I/O data structure and functions using Unix I/O data structures and functions?

## Answer:

- In principle...
- In stages...

User process



`FILE *fp`

`int fd`

**File descriptor:**  
An integer that uniquely identifies an open file



# Implementing getchar and putchar

`getchar()` calls `read()` to read one byte from fd 0

`putchar()` calls `write()` to write one byte to fd 1

```
int getchar(void)
{  unsigned char c;
   if (read(0, &c, 1) == 1)
      return (int)c;
   else
      return EOF;
}
```

```
int putchar(int c)
{  if (write(1, &c, 1) == 1)
   return c;
else
   return EOF;
}
```



# Implementing Buffering

## Problem: poor performance

- `read()` and `write()` access a physical device (e.g., a disk)
- Reading/writing one char at a time can be time consuming
- Better to read and write in larger blocks
  - Recall ***Storage Management*** lecture

## Solution: buffered I/O

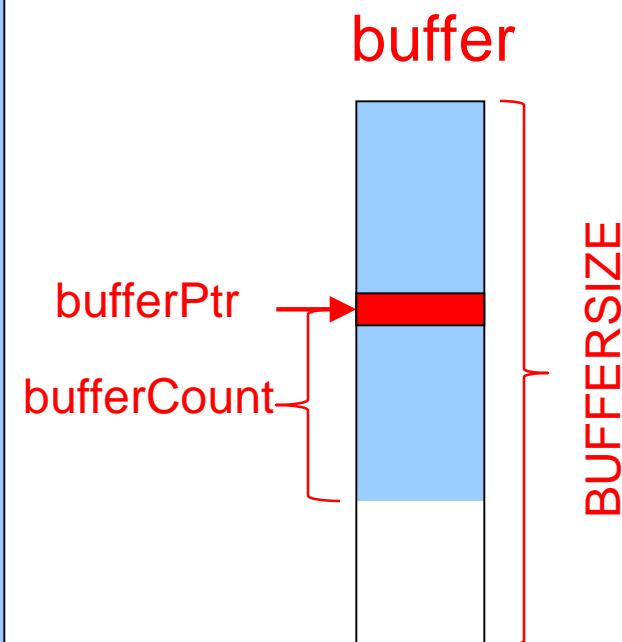
- **Read** a large block of chars from source device into a buffer
  - Provide chars from buffer to the client as needed
- **Write** individual chars to a buffer
  - “Flush” buffer contents to destination device when buffer is full, or when file is closed, or upon client request



# Implementing getchar Version 2

getchar( ) calls read( ) to read multiple chars from fd 0 into buffer

```
int getchar(void)
{
    enum {BUFFERSIZE = 4096}; /*arbitrary*/
    static unsigned char buffer[BUFFERSIZE];
    static unsigned char *bufferPtr;
    static int bufferCount = 0;
    if (bufferCount == 0) /* must read */
    {
        bufferCount =
            read(0, buffer, BUFFERSIZE);
        if (bufferCount <= 0) return EOF;
        bufferPtr = buffer;
    }
    bufferCount--;
    bufferPtr++;
    return (int)(*(bufferPtr-1));
}
```





# Implementing putchar Version 2

putchar( ) calls write( ) to write multiple chars from buffer to fd 1

```
int putchar(int c)
{  enum {BUFFERSIZE = 4096};
   static char buffer[BUFFERSIZE];
   static int bufferCount = 0;
   if (bufferCount == BUFFERSIZE) /* must write */
   {  int countWritten = 0;
      while (countWritten < bufferCount)
      {  int count =
          write(1, buffer+countWritten, BUFFERSIZE-countWritten);
          if (count <= 0) return EOF;
          countWritten += count;
      }
      bufferCount = 0;
   }
   buffer[bufferCount] = (char)c;
   bufferCount++;
   return c;
}
```

Real implementation  
also flushes buffer  
at other times



# Implementing the FILE ADT

## Observation:

- `getchar()` reads from `stdin` (fd 0)
- `putchar()` writes to `stdout` (fd 1)

## Problem:

- How to read/write from/to files other than `stdin` (fd 0) and `stdout` (fd 1)?
- Example: How to define `fgetc()` and `fputc()`?

## Solution:

- Use `FILE` structure



# Implementing the FILE ADT

```
enum {BUFFERSIZE = 4096};

struct File
{ unsigned char buffer[BUFFERSIZE]; /* buffer */
  int bufferCount; /* num chars left in buffer */
  unsigned char *bufferPtr; /* ptr to next char in buffer */
  int flags; /* open mode flags, etc. */
  int fd; /* file descriptor */
};

typedef struct File FILE;

/* Initialize standard files. */
FILE *stdin = ...
FILE *stdout = ...
FILE *stderr = ...
```

Derived from  
K&R Section 8.5

More complex  
on our system



# Implementing fopen and fclose

```
f = fopen(filename, "r")
```

- Create new **FILE** structure; set **f** to point to it
- Initialize all fields
- **f->fd = open(filename, ...)**
- Return **f**

```
f = fopen(filename, "w")
```

- Create new **FILE** structure; set **f** to point to it
- Initialize all fields
- **f->fd = creat(filename, ...)**
- Return **f**

```
fclose(f)
```

- **close(f->fd)**
- Destroy **FILE** structure



# Implementing fgetc

```
int fgetc(FILE *f)
{
    if (f->bufferCount == 0) /* must read */
    {
        f->bufferCount =
            read(f->fd, f->buffer, BUFFERSIZE);
        if (f->bufferCount <= 0) return EOF;
        f->bufferPtr = f->buffer;
    }
    f->bufferCount--;
    f->bufferPtr++;
    return (int)(*(f->bufferPtr-1));
}
```

- Accepts FILE pointer f as parameter
- Uses fields within f
- Reads from f->fd instead of 0



# Implementing fputc

```
int fputc(int c, FILE *f)
{  if (f->bufferCount == BUFFERSIZE) /* must write */
{   int countWritten = 0;
    while (countWritten < f->bufferCount)
    {   int count =
        write(f->fd, f->buffer+countWritten,
              BUFFERSIZE-countWritten);
        if (count <= 0) return EOF;
        countWritten += count;
    }
    f->bufferCount = 0;
}
f->buffer[f->bufferCount] = (char)c;
f->bufferCount++;
return c;
}
```

Real implementation  
also flushes buffer  
at other times

- Accepts FILE pointer f as parameter
- Uses fields within f
- Writes to f->fd instead of 1



# Implementing Standard C I/O Functions

Standard C Function	In Unix Implemented by Calling
fopen()	<b>open()</b> or <b>creat()</b>
fclose()	<b>close()</b>



# Implementing Standard C I/O Functions

Standard C Function	In Unix Implemented by Calling
fgetc()	read()
getchar()	fgetc()
fgets()	fgetc()
gets()	fgets()
fscanf()	fgetc()
scanf()	fscanf()



# Implementing Standard C I/O Functions

Standard C Function	In Unix Implemented by Calling
fputc()	<b>write()</b>
putchar()	fputc()
fputs()	fputc()
puts()	fputs()
fprintf()	fputc()
printf()	fprintf()



# Implementing Standard C I/O Functions

Standard C Function	In Unix Implemented by Calling
fflush()	<b>write()</b>
fseek()	<b>lseek()</b>
ftell()	<b>lseek()</b>



# Agenda

The C/Unix file abstraction

Unix I/O system calls

C's Standard IO library (FILE \*)

Implementing standard C I/O using Unix I/O

**Redirecting standard files**

Pipes



# Redirection

Unix allows programmatic redirection of **stdin**, **stdout**, or **stderr**

How?

- Use `open()`, `creat()`, and `close()` system-level functions
- Use `dup()` system-level function

`int dup(int oldfd);`

- Create a copy of file descriptor `oldfd`
- Old and new file descriptors may be used interchangeably; they refer to the same open file table and thus share file position and file status flags
- Uses the **lowest-numbered** unused descriptor for the new descriptor
- Returns the new descriptor, or -1 if an error occurred.

Paraphrasing man page



# Redirection Example

How does shell implement `somepgm > somefile`?

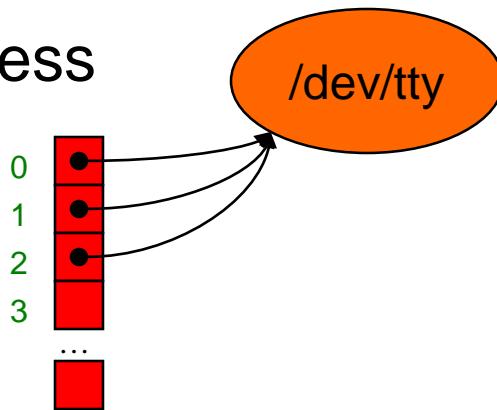
```
pid = fork();
if (pid == 0)
{ /* in child */
    fd = creat("somefile", 0600);
    close(1);
    dup(fd);
    close(fd);
    execvp(somepgm, someargv);
    fprintf(stderr, "exec failed\n");
    exit(EXIT_FAILURE);
}
/* in parent */
wait(NULL);
```



# Redirection Example Trace (1)

## Parent Process

File descriptor table



```
pid = fork();
if (pid == 0)
{ /* in child */
    fd = creat("somefile", 0600);
    close(1);
    dup(fd);
    close(fd);
    execvp(somepgm, someargv);
    fprintf(stderr, "exec failed\n");
    exit(EXIT_FAILURE);
}
/* in parent */
wait(NULL);
```

Parent has file descriptor table; first three point to “terminal”



# Redirection Example Trace (2)

Parent Process

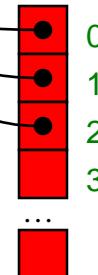
File descriptor table



/dev/tty

Child Process

File descriptor table



```
pid = fork();
if (pid == 0)
{ /* in child */
    fd = creat("somefile", 0600);
    close(1);
    dup(fd);
    close(fd);
    execvp(somepgm, someargv);
    fprintf(stderr, "exec failed\n");
    exit(EXIT_FAILURE);
}
/* in parent */
wait(NULL);
```

```
pid = fork();
if (pid == 0)
{ /* in child */
    fd = creat("somefile", 0600);
    close(1);
    dup(fd);
    close(fd);
    execvp(somepgm, someargv);
    fprintf(stderr, "exec failed\n");
    exit(EXIT_FAILURE);
}
/* in parent */
wait(NULL);
```

Parent forks child; child has identical-but distinct file descriptor table



# Redirection Example Trace (3)

Parent Process

File descriptor table



/dev/tty

Child Process

File descriptor table



```
pid = fork();
if (pid == 0)
{ /* in child */
    fd = creat("somefile", 0600);
    close(1);
    dup(fd);
    close(fd);
    execvp(somepgm, someargv);
    fprintf(stderr, "exec failed\n");
    exit(EXIT_FAILURE);
}
/* in parent */
wait(NULL);
```

```
pid = fork();
if (pid == 0)
{ /* in child */
    fd = creat("somefile", 060);
    close(1);
    dup(fd);
    close(fd);
    execvp(somepgm, someargv);
    fprintf(stderr, "exec failed\n");
    exit(EXIT_FAILURE);
}
/* in parent */
wait(NULL);
```

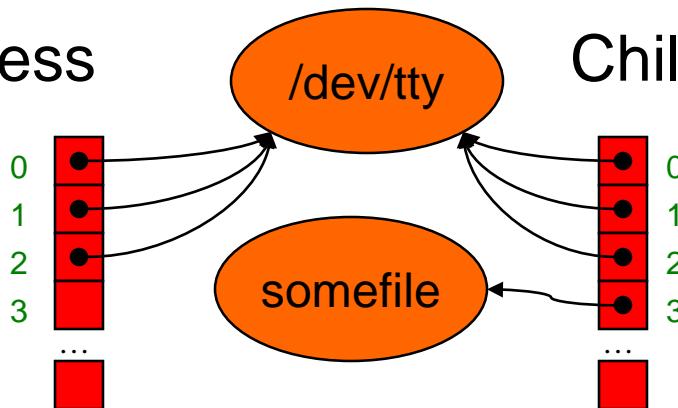
Let's say OS gives CPU to parent; parent waits



# Redirection Example Trace (4)

Parent Process

File descriptor table



Child Process

File descriptor table

```
pid = fork();
if (pid == 0)
{ /* in child */
    fd = creat("somefile", 0600);
    close(1);
    dup(fd);
    close(fd);
    execvp(somepgm, someargv);
    fprintf(stderr, "exec failed\n");
    exit(EXIT_FAILURE);
}
/* in parent */
wait(NULL);
```

```
pid = fork();
if (pid == 0)
{ /* in child */
    fd = creat("somefile", 0600);
    close(1);
    dup(fd);
    close(fd);
    execvp(somepgm, someargv);
    fprintf(stderr, "exec failed\n");
    exit(EXIT_FAILURE);
}
/* in parent */
wait(NULL);
```

3

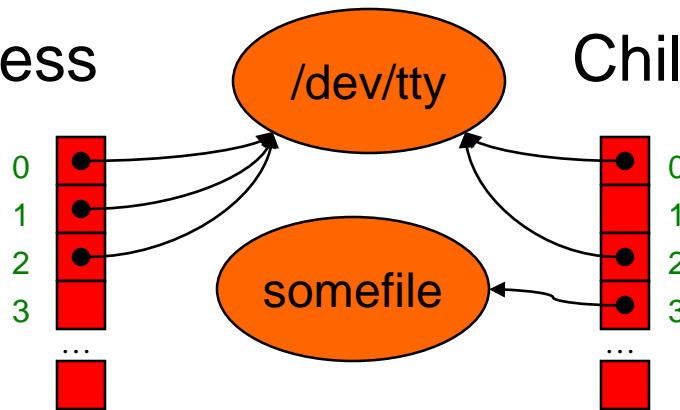
OS gives CPU to child; child creates somefile



# Redirection Example Trace (5)

Parent Process

File descriptor table



Child Process

File descriptor table

```
pid = fork();
if (pid == 0)
{ /* in child */
    fd = creat("somefile", 0600);
    close(1);
    dup(fd);
    close(fd);
    execvp(somepgm, someargv);
    fprintf(stderr, "exec failed\n");
    exit(EXIT_FAILURE);
}
/* in parent */
wait(NULL);
```

```
pid = fork();
if (pid == 0)
{ /* in child */
    fd = creat("somefile", 0600);
    close(1);
    dup(fd);
    close(fd);
    execvp(somepgm, someargv);
    fprintf(stderr, "exec failed\n");
    exit(EXIT_FAILURE);
}
/* in parent */
wait(NULL);
```

3

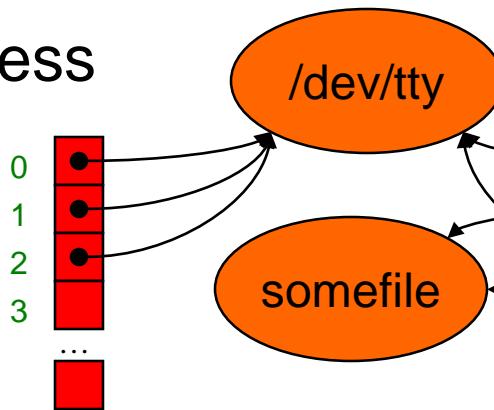
Child closes file descriptor 1 (stdout)



# Redirection Example Trace (6)

Parent Process

File descriptor table



Child Process

File descriptor table

```
pid = fork();
if (pid == 0)
{ /* in child */
    fd = creat("somefile", 0600);
    close(1);
    dup(fd);
    close(fd);
    execvp(somepgm, someargv);
    fprintf(stderr, "exec failed\n");
    exit(EXIT_FAILURE);
}
/* in parent */
wait(NULL);
```

```
pid = fork();
if (pid == 0)
{ /* in child */
    fd = creat("somefile", 0600);
    close(1);
    dup(fd);
    close(fd);
    execvp(somepgm, someargv);
    fprintf(stderr, "exec failed\n");
    exit(EXIT_FAILURE);
}
/* in parent */
wait(NULL);
```

3

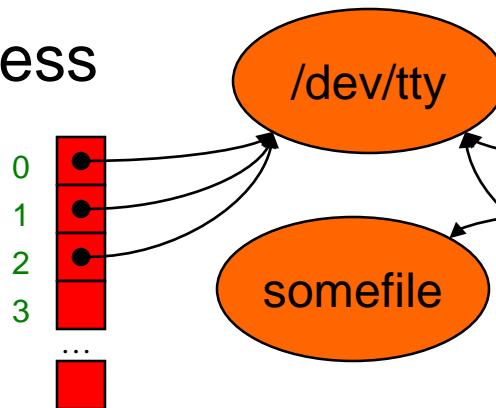
Child duplicates file descriptor 3 into first unused spot



# Redirection Example Trace (7)

Parent Process

File descriptor table



Child Process

File descriptor table

```
pid = fork();
if (pid == 0)
{ /* in child */
    fd = creat("somefile", 0600);
    close(1);
    dup(fd);
    close(fd);
    execvp(somepgm, someargv);
    fprintf(stderr, "exec failed\n");
    exit(EXIT_FAILURE);
}
/* in parent */
wait(NULL);
```

```
pid = fork();
if (pid == 0)
{ /* in child */
    fd = creat("somefile", 0600);
    close(1);
    dup(fd);
    close(fd);
    execvp(somepgm, someargv);
    fprintf(stderr, "exec failed\n");
    exit(EXIT_FAILURE);
}
/* in parent */
wait(NULL);
```

3

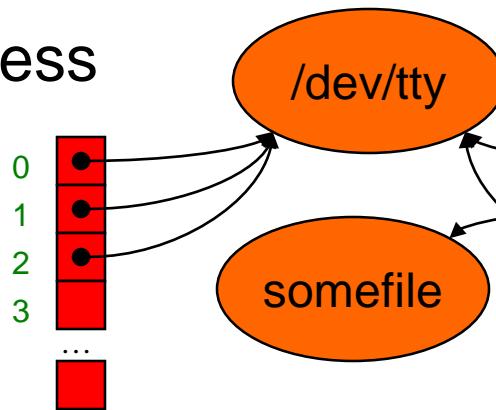
Child closes file descriptor 3



# Redirection Example Trace (8)

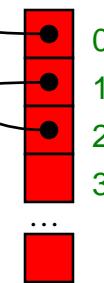
Parent Process

File descriptor table



Child Process

File descriptor table



```
pid = fork();
if (pid == 0)
{ /* in child */
    fd = creat("somefile", 0600);
    close(1);
    dup(fd);
    close(fd);
    execvp(somepgm, someargv);
    fprintf(stderr, "exec failed\n");
    exit(EXIT_FAILURE);
}
/* in parent */
wait(NULL);
```

Child calls `execvp()`

3

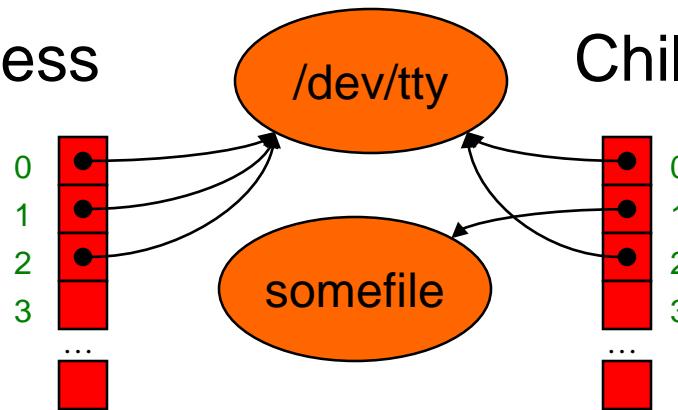
```
pid = fork();
if (pid == 0)
{ /* in child */
    fd = creat("somefile", 0600);
    close(1);
    dup(fd);
    close(fd);
    execvp(somepgm, someargv);
    fprintf(stderr, "exec failed\n");
    exit(EXIT_FAILURE);
}
/* in parent */
wait(NULL);
```



# Redirection Example Trace (9)

Parent Process

File descriptor table



Child Process

File descriptor table

```
pid = fork();
if (pid == 0)
{ /* in child */
    fd = creat("somefile", 0600);
    close(1);
    dup(fd);
    close(fd);
    execvp(somepfm, someargv);
    fprintf(stderr, "exec failed\n");
    exit(EXIT_FAILURE);
}
/* in parent */
wait(NULL);
```

*somepgm*

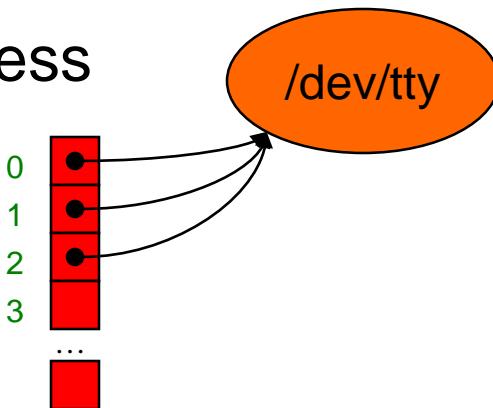
Somepgm executes with stdout redirected to somefile



# Redirection Example Trace (10)

## Parent Process

File descriptor table



```
pid = fork();
if (pid == 0)
{ /* in child */
    fd = creat("somefile", 0600);
    close(1);
    dup(fd);
    close(fd);
    execvp(somefile, someargv);
    fprintf(stderr, "exec failed\n");
    exit(EXIT_FAILURE);
}
/* in parent */
wait(NULL);
```

Somepgm exits; parent returns from `wait()` and proceeds



# Agenda

The C/Unix file abstraction

Unix I/O system calls

C's Standard IO library (FILE \*)

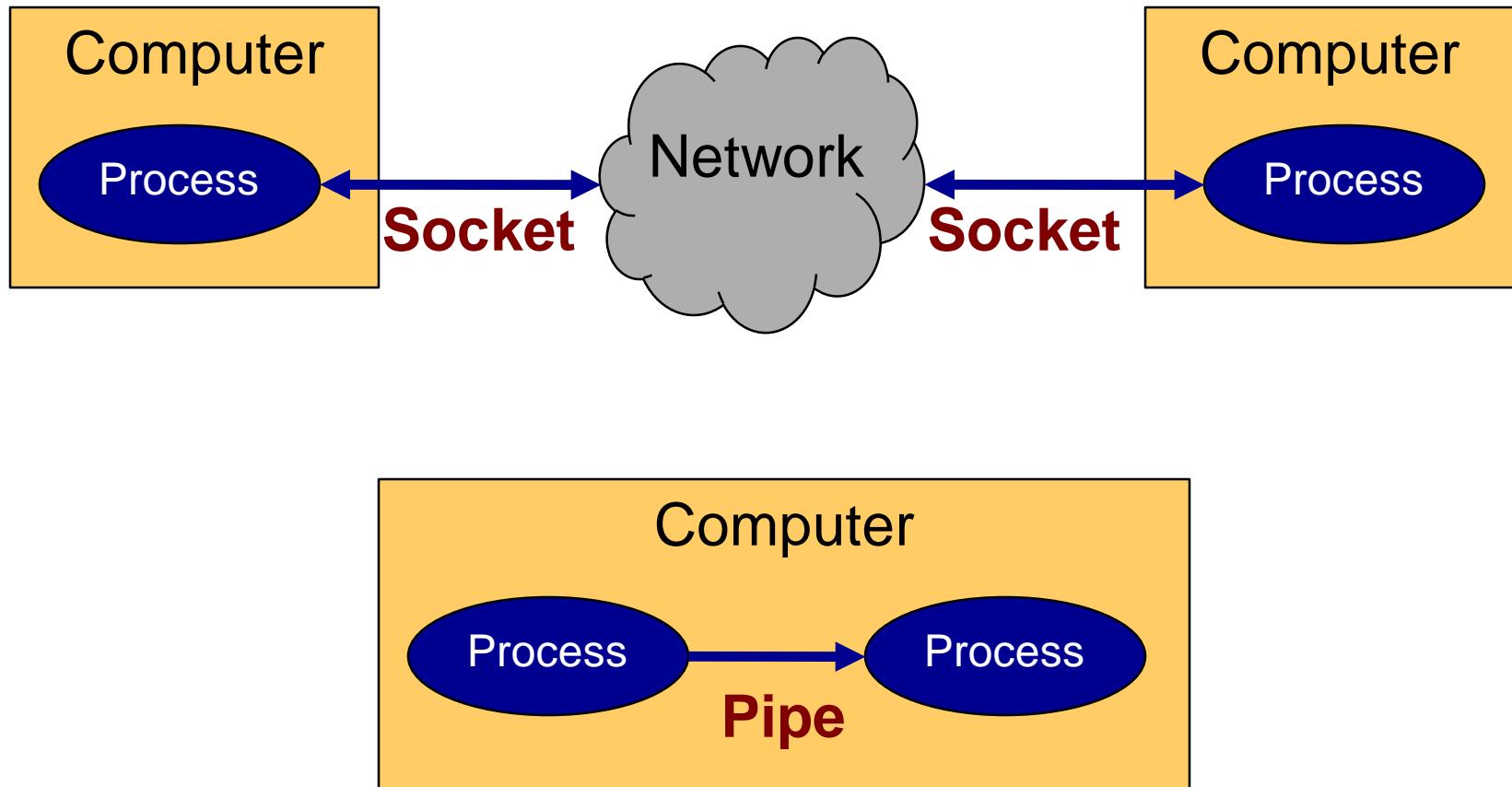
Implementing standard C I/O using Unix I/O

Redirecting standard files

**Pipes**



# Inter-Process Communication (IPC)





# IPC Mechanisms

## Socket

- Mechanism for **two-way** communication between processes on **any computers** on same network
- Processes created independently
- Used for client/server communication (e.g., Web)

## Pipe

- Mechanism for **one-way** communication between processes on the **same computer**
- Allows parent process to communicate with child process
- Allows two “sibling” processes to communicate
- Used mostly for a **pipeline** of **filters**

Both support **file** abstraction



# Pipes, Filters, and Pipelines

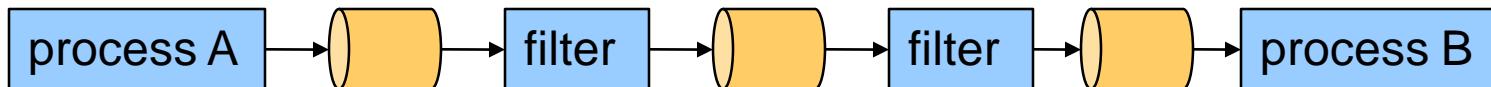
## Pipe



**Filter:** Program that reads from stdin and writes to stdout



**Pipeline:** Combination of pipes and filters





# Pipeline Examples

When debugging your shell program...

```
grep alloc *.c
```

- In all of the .c files in the working directory, display all lines that contain “alloc”

```
cat *.c | decomment | grep alloc
```

- In all of the .c files in the working directory, display all non-comment lines that contain “alloc”

```
cat *.c | decomment | grep alloc | more
```

- In all of the .c files in the working directory, display all non-comment lines that contain “alloc”, one screen at a time



# Creating a Pipe

```
int pipe(int pipefd[2])
```

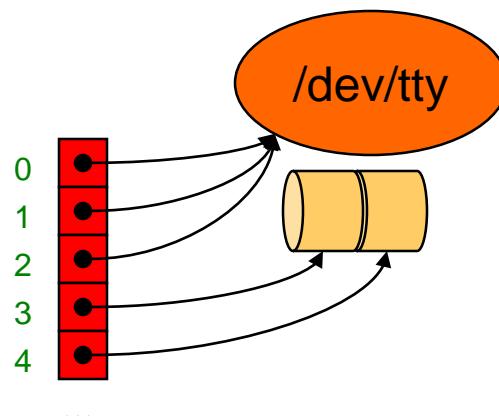
- `pipe()` creates a pipe, a unidirectional data channel that can be used for interprocess communication
- The array `pipefd` is used to return two file descriptors referring to the ends of the pipe
- `pipefd[0]` refers to the read end of the pipe
- `pipefd[1]` refers to the write end of the pipe
- Data written to the write end of the pipe is buffered by the kernel until it is read from the read end of the pipe
- Quoting `man -s2 pipe`



# Pipe Example 1 (1)

Parent process sends data to child process

```
int p[2];
...
pipe(p)
pid = fork();
if (pid == 0)
{ /* in child */
    close(p[1]);
    /* Read from fd p[0] */
    exit(0);
}
/* in parent */
close(p[0]);
/* Write to fd p[1] */
wait(NULL);
```



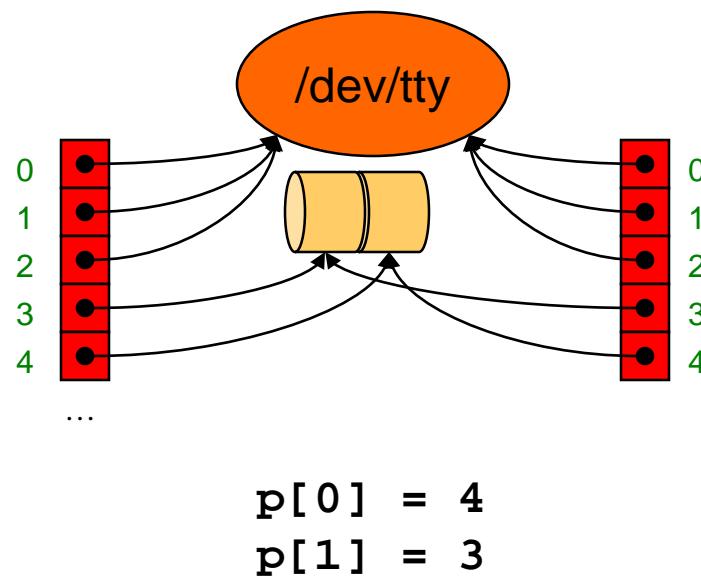
**p[0] = 4**  
**p[1] = 3**



# Pipe Example 1 (2)

Parent process sends data to child process

```
int p[2];
...
pipe(p)
pid = fork();
if (pid == 0)
{ /* in child */
close(p[1]);
/* Read from fd p[0] */
exit(0);
}
/* in parent */
close(p[0]);
/* Write to fd p[1] */
wait(NULL);
```



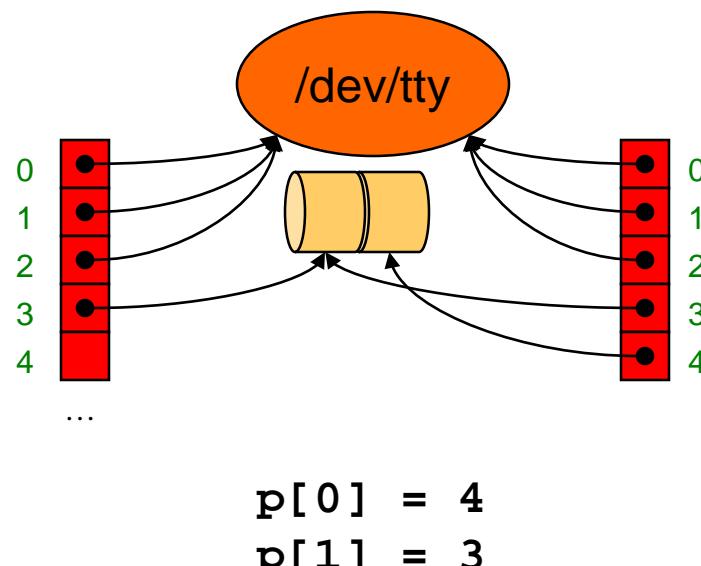
```
int p[2];
...
pipe(p)
pid = fork();
if (pid == 0)
{ /* in child */
close(p[1]);
/* Read from fd p[0] */
exit(0);
}
/* in parent */
close(p[0]);
/* Write to fd p[1] */
wait(NULL);
```



# Pipe Example 1 (3)

Parent process sends data to child process

```
int p[2];
...
pipe(p)
pid = fork();
if (pid == 0)
{ /* in child */
close(p[1]);
/* Read from fd p[0] */
exit(0);
}
/* in parent */
close(p[0]);
/* Write to fd p[1] */
wait(NULL);
```



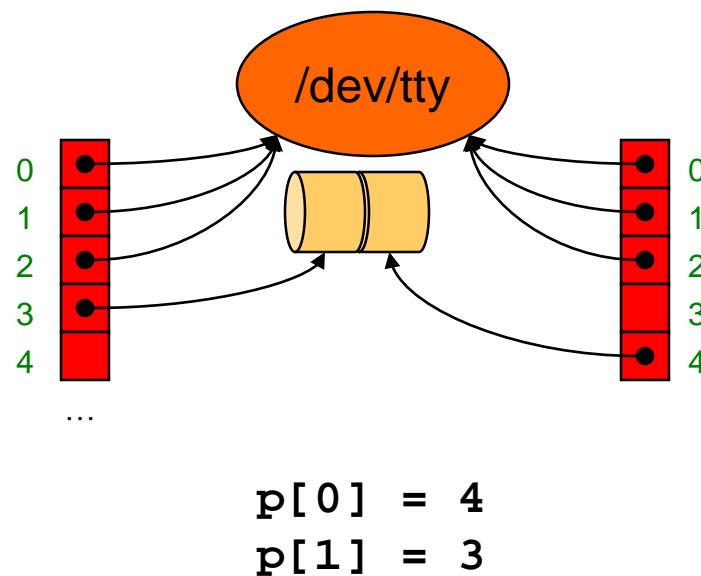
```
int p[2];
...
pipe(p)
pid = fork();
if (pid == 0)
{ /* in child */
close(p[1]);
/* Read from fd p[0] */
exit(0);
}
/* in parent */
close(p[0]);
/* Write to fd p[1] */
wait(NULL);
```



# Pipe Example 1 (4)

Parent process sends data to child process

```
int p[2];
...
pipe(p)
pid = fork();
if (pid == 0)
{ /* in child */
close(p[1]);
/* Read from fd p[0] */
exit(0);
}
/* in parent */
close(p[0]);
/* Write to fd p[1] */
wait(NULL);
```



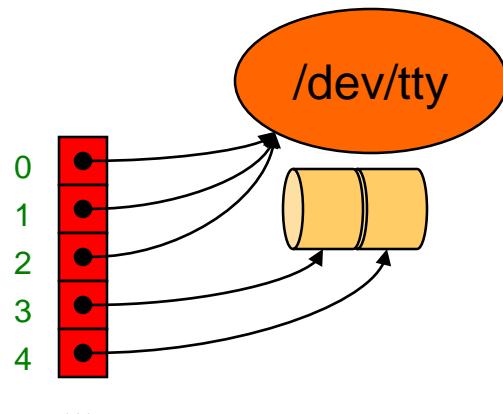
```
int p[2];
...
pipe(p)
pid = fork();
if (pid == 0)
{ /* in child */
close(p[1]);
/* Read from fd p[0] */
exit(0);
}
/* in parent */
close(p[0]);
/* Write to fd p[1] */
wait(NULL);
```



# Pipe Example 2 (1)

Parent sends data to child through stdin/stdout

```
int p[2];
...
pipe(p)
pid = fork();
if (pid == 0)
{ /* in child */
    close(0);
    dup(p[0]);
    close(p[0]);
    close(p[1]);
    /* Read from stdin */
    exit(0);
}
/* in parent */
close(1);
dup(p[1]);
close(p[1]);
close(p[0]);
/* write to stdout */
wait(NULL);
```



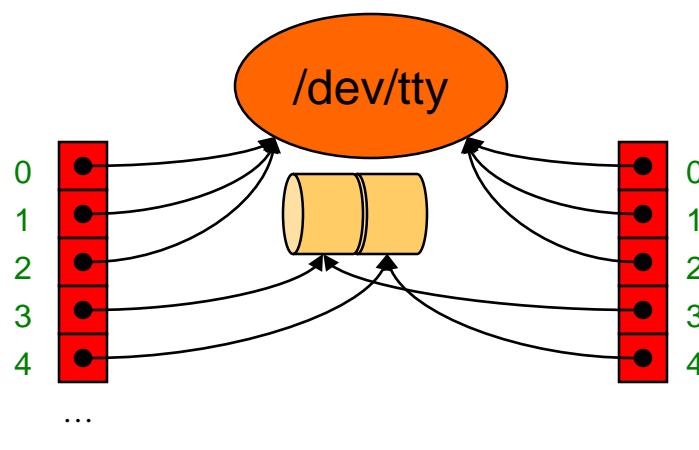
**p[0] = 4**  
**p[1] = 3**



# Pipe Example 2 (2)

Parent sends data to child through stdin/stdout

```
int p[2];
...
pipe(p)
pid = fork();
if (pid == 0)
{ /* in child */
    close(0);
    dup(p[0]);
    close(p[0]);
    close(p[1]);
    /* Read from stdin */
    exit(0);
}
/* in parent */
close(1);
dup(p[1])
close(p[1]);
close(p[0]);
/* write to stdout */
wait(NULL);
```



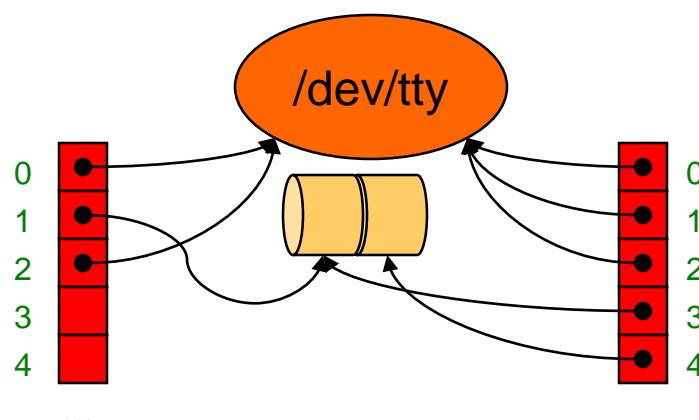
```
int p[2];
...
pipe(p)
pid = fork();
if (pid == 0)
{ /* in child */
    close(0);
    dup(p[0]);
    close(p[0]);
    close(p[1]);
    /* Read from stdin */
    exit(0);
}
/* in parent */
close(1);
dup(p[1])
close(p[1]);
close(p[0]);
/* write to stdout */
wait(NULL);
```



# Pipe Example 2 (3)

Parent sends data to child through stdin/stdout

```
int p[2];
...
pipe(p)
pid = fork();
if (pid == 0)
{ /* in child */
    close(0);
    dup(p[0]);
    close(p[0]);
    close(p[1]);
    /* Read from stdin */
    exit(0);
}
/* in parent */
close(1);
dup(p[1])
close(p[1]);
close(p[0]);
/* write to stdout */
wait(NULL);
```



`p[0] = 4`  
`p[1] = 3`

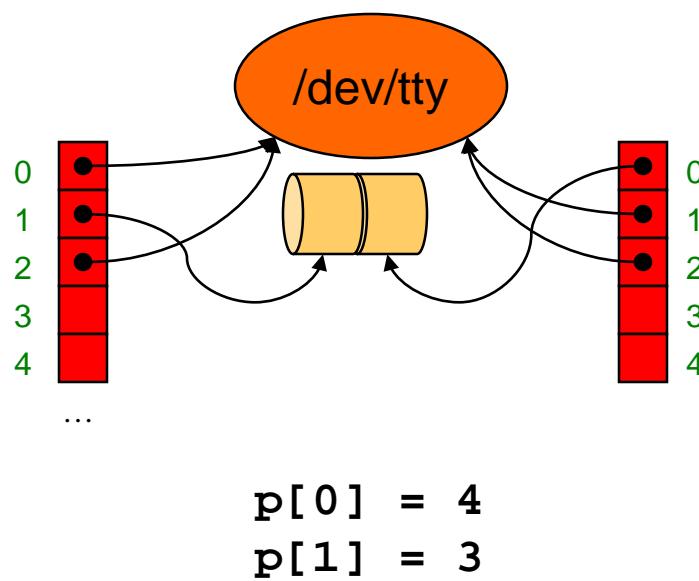
```
int p[2];
...
pipe(p)
pid = fork();
if (pid == 0)
{ /* in child */
    close(0);
    dup(p[0]);
    close(p[0]);
    close(p[1]);
    /* Read from stdin*/
    exit(0);
}
/* in parent */
close(1);
dup(p[1])
close(p[1]);
close(p[0]);
/* write to stdout */
wait(NULL);
```



# Pipe Example 2 (4)

Parent sends data to child through stdin/stdout

```
int p[2];
...
pipe(p)
pid = fork();
if (pid == 0)
{ /* in child */
    close(0);
    dup(p[0]);
    close(p[0]);
    close(p[1]);
    /* Read from stdin */
    exit(0);
}
/* in parent */
close(1);
dup(p[1]);
close(p[1]);
close(p[0]);
/* write to stdout */
wait(NULL);
```



Now add in execs, and  
you get the shell's  
implementation of pipes!

```
int p[2];
...
pipe(p)
pid = fork();
if (pid == 0)
{ /* in child */
    close(0);
    dup(p[0]);
    close(p[0]);
    close(p[1]);
    /* Read from stdin */
    exit(0);
}
/* in parent */
close(1);
dup(p[1]);
close(p[1]);
close(p[0]);
/* write to stdout */
wait(NULL);
```



# Summary

The C/Unix file abstraction

Unix I/O

- File descriptors, file descriptor tables, file tables
- `creat()`, `open()`, `close()`, `read()`, `write()`, `lseek()`

C's Standard I/O

- `FILE` structure
- `fopen()`, `fclose()`, `fgetc()`, `fputc()`, ...

Implementing standard C I/O using Unix I/O

- Buffering

Redirecting standard files

- `dup()`

Pipes

- `pipe()`