

Using and storing the index

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Review: Inverted Index

- For each term, keep list of document entries, one for each document in which it appears: a **postings list**
 - Document entry is list of positions at which term occurs and attributes for each occurrence: a **posting**
- Keep **summary term information**
- Keep **summary document information**

meta-data



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Consider “advanced search” queries

Content Coordination

- Phrases
- Numeric range
- NOT
- OR

Document Meta-data

- Language
- Geographic region
- File format
- Date published
- From specific domain
- Specific licensing rights
- Filtered by “safe search”

Issue of **efficient retrieval**

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Basic retrieval algorithms

- One term
- AND of several terms
- OR of several terms
- NOT term
- proximity

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Basic postings list processing: **Merging** posting lists

- Have two lists must **coordinate**
 - Find shared entries and do “something”
 - “something” changes for different operations
 - Set operations UNION? INTERSECTION? DIFFERENCE? ...
 - Filter with document meta-data as process

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Basic retrieval algorithms: using merging of postings lists

HOW?

- AND of several terms
- OR of several terms
- NOT term
- proximity

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Basic retrieval algorithms

- One term:
 - look up posting list in (inverted) index
- AND of several terms:
 - **Intersect** posting lists of the terms: **a list merge**
- OR of several terms:
 - **Union** posting lists of the terms
 - eliminate **duplicates**: **a list merge**
- NOT term
 - If *terms* AND NOT(*other terms*), take a difference
 - **a list merge** (similar to AND)
- Proximity
 - **a list merge** (similar to AND)

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Algorithms for Merging Postings Lists: two unsorted lists

How?

Role call

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Merging two unsorted lists

- **Read 2nd list over and over** - once for each entry on 1st list
 - computationally expensive
 - time $O(|L_1|*|L_2|)$ where $|L|$ length list L
- **Build hash table** on entry values;
insert entries of one list, then other;
look for collisions
 - must have good hash table
 - unwanted collisions expensive
 - often can't fit in memory: disk version
- **Sort lists**; use algorithm for sorted lists
 - often lists on disk: external sort
 - can sort in $O(|L| \log |L|)$ operations

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Sorted lists

- Lists sorted by some identifier
 - same identifier both lists; not nec. unique
- Read both lists in “parallel”
 - **Classic list merge**:
(sorted list₁, sorted list₂) ⇒ sorted set union
 - General merge: if no duplicates, get time $|L_1|+|L_2|$
- Build lists so sorted
 - pay cost at most once
 - maybe get sorted order “naturally”
- If only one list sorted, can do binary search of sorted list for entries of other list
 - Must be able to binary search! - **rare!**
 - can't binary search disk

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Keys for documents

For posting lists, entries are documents
What value is used to sort?

- Unique **document IDs**
 - can still be duplicate documents
 - consider for Web when consider crawling
- document **scoring function** that is **independent of query**
 - PageRank, HITS authority
 - sort on document IDs as secondary key
 - allows for **approximate “highest k” retrieval**
 - approx. k highest ranking doc.s for a query

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Keys within document list

Processing within document posting

- Proximity of terms
 - merge lists of terms occurrences within same doc.
- Sort on term position

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Computing document score

- “On fly”- as find each satisfying document
- Separate phase after build list of satisfying documents
- For either, must sort doc.s by score

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Web query processing: limiting size

- For Web-scale collections, may **not process complete posting list** for each term in query
 - at least not initially
- Need docs **sorted first on global (static) quantity**
 - why not by term frequency for doc?
- Only **take first k doc.s** on each term list
 - k depends on query - how?
 - k depends on how many want to be able to return
 - Google: 1000 max returns
 - Flaws w/ partial retrieval from each list?
- Other limits? **query size**
 - Google: 32 words max query size

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Limiting size with term-based sorting

- Can sort doc.s on postings list by score of term
 - term frequency + ...
- Lose linear merge - salvage any?
- Tiered index:
 - tier 1: docs with highest term-based scores, sorted by ID or global quantity
 - tier 2: docs in next bracket of score quality, sorted
 - etc.
 - need to decide size or range of brackets
- If give up AND of query terms, can use idf too
 - only consider terms with high idf = rarer terms

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Data structure for inverted index?

How access individual terms and each associated postings list?

Assume an entry for each term points to its posting list

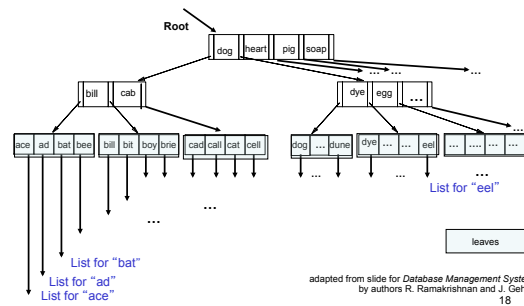
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Data structure for inverted index?

- Sorted array:
 - binary search IF can keep in memory
 - High overhead for additions
- Hashing
 - Fast look-up
 - Collisions
- Search trees: B+-trees
 - Maintain balance - always log look-up time
 - Can insert and delete

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Example B+ Tree
order = 2: 2 to 4 search keys per interior node



B+- trees

- All index entries are at leaves
- Order m B+ tree has $m+1$ to $2m+1$ children for each interior node
 - except root can have as few as 2 children
- Look up: follow root to leaf by keys in interior nodes
- Insert:
 - find leaf in which belongs
 - If leaf full, split
 - Split can propagate up tree
- Delete:
 - Merge or redistribute from too-empty leaf
 - Merge can propagate up tree

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Disk-based B+ trees for large data sets

- Each leaf is file page (block) on disk
- Each interior node is file page on disk
- Keep top of tree in buffer (RAM)
- Typical sizes:
 - $m \sim 200$;
 - average fanout ~ 267
 - Height 4 gives ~ 5 billion entries

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prefix key B+ trees

- Save space
- Each interior node key is shortest prefix of word needed to distinguish which child pointer to follow
 - Allows more keys per interior node
 - higher fanout
 - fanout determined by what can fit
 - keep at least 1/2 full

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Revisit hashing - on disk

- hash of term gives address of bucket on disk
- bucket contains pairs (term, address of first page of postings list)
- bucket occupies one file page



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