

## Finding near-duplicate documents

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## Duplicate versus near duplicate documents

- Duplicate = **identical**?
- Near duplicate:  
**small structural differences**
  - not just content similarity
- define “small”
  - date change?
  - small edits?
  - metadata change?
  - other?

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## Applications

- Crawling network
- Indexing
- Returning query results
  - cluster near duplicates; return 1
- Plagiarism

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## Framework

- Algorithm to assign quantitative degree of similarity between documents
- Issues
  - What is basic token for documents?
    - character
    - word/term
  - What is threshold for “near duplicate”?
  - What are computational costs?

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## Classic document comparison

- Edit distance
  - count deletions, additions, substitutions to convert  $Doc_1$  into  $Doc_2$
  - each action can have different cost
  - applications
    - UNIX “diff”
    - similarity of genetic sequences
- Edit distance algorithm
  - dynamic programming
  - time  $O(m*n)$  for strings length  $m$  and  $n$

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## Edit distance for collections

- token = word
  - compare other applications
- Cost is  $O(\sum_{i,j} |Doc_i| * |Doc_j|)$
- Right sense of similarity?

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## Addressing computation cost

A general paradigm to find duplicates in N docs:

1. Define function  $f$  capturing contents of each document in **one number**  
"Hash function", "signature", "fingerprint"
2. Create  $\langle f(\text{doc}_i), \text{ID of doc}_i \rangle$  pairs
3. Sort the pairs
4. Recognize duplicate or near-duplicate documents as having the same  $f$  value or  $f$  values within a **small threshold**

Compare: computing a similarity score on pairs of documents

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## Optimistic cost

A general paradigm to find duplicates in N docs:

1. Define function  $f$  capturing contents of each document in **one number**  $O(|\text{doc}|)$   
"Hash function", "signature", "fingerprint"
2. Create  $\langle f(\text{doc}_i), \text{ID of doc}_i \rangle$  pairs  $O(\sum_{i=1..N} (|\text{doc}_i|))$
3. Sort the pairs  $O(N \log N)$
4. Recognize duplicate or near-duplicate documents as having the same  $f$  value or  $f$  values within a **small threshold**  $O(N)$

Compare: computing a similarity score on pairs of documents

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## General paradigm: details

1. Define function  $f$  capturing contents of each document in one number  
"Hash function", "signature", "sketch", "fingerprint"
2. Create  $\langle f(\text{doc}_i), \text{ID of doc}_i \rangle$  pairs
3. Sort the pairs
4. Recognize duplicate or near-duplicate documents as having the same  $f$  value or  $f$  values within a small threshold
  - recognize exact duplicates:
    - threshold = 0
    - examine documents to verify duplicates
  - recognize near-duplicates  
Use small "small threshold"  
=> "near duplicate" not transitive

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## "Syntactic clustering"

We will look at this one example:

Andrei Z. Broder, Steven C. Glassman, Mark S. Manasse, and Geoffrey Zweig, [Syntactic Clustering of the Web](#)  
Sixth International WWW Conference, 1997.

- "syntactic similarity" versus semantic  
Sequences of words
- Finding near duplicates
- Doc = sequence of words  
Word = Token
- Uses **sampling**
- Similarity based on **shingles**
- Does compare documents

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## Shingles

- A **w-shingle** is a contiguous subsequence of  $w$  words
- The **w-shingling of doc D**,  $S(D, w)$  is the set of **unique** w-shingles of D

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## Similarity of docs with shingles

- For **fixed w**, **resemblance** of docs A and B :  
$$r(A, B) = \frac{|S(A) \cap S(B)|}{|S(A) \cup S(B)|}$$

Jaccard coefficient
- For **fixed w**, **containment** of doc A in doc B :  
$$C(A, B) = \frac{|S(A) \cap S(B)|}{|S(A)|}$$
- For **fixed w**, **resemblance distance** between docs A and B :  
$$D(A, B) = 1 - r(A, B)$$

Is a metric (triangle inequality)

**Note we are now comparing documents!**

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## Example

A: "a rose is red a rose is white"

4-shingles:

"a rose is red"  
 "rose is red a"  
 "is red a rose"  
 "red a rose is"  
 "a rose is white"

B: "a rose is white a rose is red"

4-shingles:

"a rose is white"  
 "rose is white a"  
 "is white a rose"  
 "white a rose is"  
 "a rose is red"

$$r(A, B) = 0.25$$

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## Compare

A: "a rose is red a rose is white"

3-shingles:

"a rose is"  
 "rose is red"  
 "is red a"  
 "red a rose"  
 "rose is white"

B: "a rose is white a rose is red"

3-shingles:

"a rose is"  
 "rose is white"  
 "is white a"  
 "white a rose"  
 "rose is red"

$$r(A, B) = 0.43$$

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## Sample of shingles

Want to **estimate**  $r$  and/or  $c$

Do this by calculating **approximation on a sample of shingles for fixed  $w$**

- 1-to-1 map each shingle to integer in fixed, large range  $R$   
 - 64-bit hash,  $R=[0, 2^{64}-1]$
- Let  $\Pi$  be a random permutation from  $R$  to  $R$
- For any  $S(D)$  define:  
 $H(D)$  = Set of **integer hash values** corresponding to shingles in  $S(D)$   
 $\Pi(D)$  = Set of permuted values in  $H(D)$   
 $x(\Pi, D)$  = **smallest integer in  $\Pi(D)$**

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## Sketch of shingles

- Let  $\Pi_1, \dots, \Pi_m$  be  $m$  random permutations  $R \rightarrow R$   
 - text:  $m=20$

The sketch of doc  $D$  for  $\Pi_1, \dots, \Pi_m$  is

$$\psi(D) = \{x(\Pi_i, D) \mid 1 \leq i \leq m\}$$

- doc  $\rightarrow$  set shingles  $\rightarrow$  set integers  
 $\rightarrow m$  sets permuted integers  
 $\rightarrow m$  smallest integers: one per permutation

Sketch is a **sampling**

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## Approximation of resemblance

Theorem:

For random permutation  $\Pi$ :

$$r(A, B) = P(x(\Pi, A) = x(\Pi, B))$$

Estimate  $P(x(\Pi, A) = x(\Pi, B))$  as

$$|\psi(A) \cap \psi(B)| / m$$

recall  $m$  is # permutations

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## Last time

- Defining and detecting near-duplicate documents

## Today

- Example Syntactic Clustering algorithm
- Finish detecting near-duplicates
- Non-text retrieval

## Next

- Finish non-text retrieval

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## Example: compare

A: "a rose is red a rose is white"

3-shingles:

1 "a rose is"      B: "a rose is white a rose is red"  
 2 "rose is red"      3-shingles:  
 3 "is red a"      "a rose is" 1  
 4 "red a rose"      "rose is white" 5  
 5 "rose is white"      "is white a" 6  
    "white a rose" 7  
    "rose is red" 2

$r(A, B) = 0.43$

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## Example mappings

- $R = [0, 10000]$
- Let  $H(i) = i * 1000; 1 \leq i \leq 7$
- Let  $m=5$
- Define a permutation
  - Example
    - Get  $\text{randval} = \text{Math.random}()$
    - Compute function of  $\text{randval}$  and  $H(i)$  to get  $\Pi(i)$
- Do 5 times for 5 permutations

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$\psi(A) = \{x(\Pi, A) \mid 1 \leq i \leq m\} = \{568, 1150, 6119, 6880, 1905\}$

$\Pi_1$ :	<u>568</u> 1136 1705 2273 2842 3410 3979	$\Pi_2$ :	<u>1150</u> 2301 3452 4602 5753	$\Pi_3$ :	<u>9223</u> 8447 7671 6895 <u>6119</u> 5343 4567
$\Pi_4$ :	9376 8752 8128 <u>7504</u> <u>6880</u> 6256 5633	$\Pi_5$ :	2976 5952 8929 <u>1905</u> 4881 7858 834		

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$\psi(B) = \{x(\Pi, B) \mid 1 \leq i \leq m\} = \{568, 1150, 4567, 5633, 834\}$

$\Pi_1$ :	<u>568</u> 1136 1705 2273 2842 3410 3979	$\Pi_2$ :	<u>1150</u> 2301 3452 4602 5753 6904 8054	$\Pi_3$ :	<u>9223</u> 8447 7671 6895 <u>6119</u> 5343 <u>4567</u>
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$\psi(A) = \{x(\Pi, A) \mid 1 \leq i \leq m\} = \{568, 1150, 6119, 6880, 1905\}$   
 $\psi(B) = \{x(\Pi, B) \mid 1 \leq i \leq m\} = \{568, 1150, 4567, 5633, 834\}$

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$\Pi_4$ :	9376 8752 8128 7504 6880 6256 5633	$\Pi_5$ :	2976 5952 8929 1905 4881 7858 834		

Resemblance estimate:  
 $|\psi(A) \cap \psi(B)| / m$   
 $= 2/5 = .4$   
 Actual resemblance  
 $= 3/7 = .43$

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## Algorithm used (text's version)

1. Calculate *sketch*  $\psi(D_i)$  for every doc  $D_i$
2. Calculate  $|\psi(D_i) \cap \psi(D_j)| = ct_{ij}$  for each non-empty intersection:
  - i. Produce list of <shingle value, docID> pairs for all shingle values  $x(\Pi_k, D_i)$  in the sketch for each doc.
  - ii. Sort the list by shingle value
  - iii. Produce all triples <ID( $D_i$ ), ID( $D_j$ ),  $ct_{ij}$ > for which  $ct_{ij} > 0$   
 This *not linear-time* for the list of docs for one shingle value
3. Recognize duplicate, near-duplicate documents: **resemblance**  $ct_{ij}/m$  above a large **threshold**

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### Algorithm cost

1. Calculate *sketch*  $\psi(D_i)$  for every  $D_i$   $O(\sum_i m|D_i|)$
2. Calculate  $|\psi(D_i) \cap \psi(D_j)| = ct_{ij}$  for each non-empty intersection:
  - i. Produce list of <shingle value, docID> pairs for all shingle values  $x(\Pi_k, D_i)$  in the sketch for each doc.
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### Algorithm cost

1. Calculate *sketch*  $\psi(D_i)$  for every  $D_i$   $O(\sum_i m|D_i|)$
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  - i. Produce list of <shingle value, docID> pairs for all shingle values  $x(\Pi_k, D_i)$  in the sketch for each doc.
  - ii. Sort the list by shingle value  $O(mN \log(mN))$
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This *not linear-time* for the list of docs for one shingle value  $O(mN^2)$
3. Recognize duplicate, near-duplicate documents: *resemblance*  $ct_{i,j}/m$  above a large *threshold*

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### Algorithm cost

1. Calculate *sketch*  $\psi(D_i)$  for every  $D_i$   $O(\sum_i m|D_i|)$
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This *not linear-time* for the list of docs for one shingle value  $O(mN^2)$
3. Recognize duplicate, near-duplicate documents: *resemblance*  $ct_{i,j}/m$  above a large *threshold*  $O(N^2)$

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### Revisit the original paradigm

A general paradigm to find duplicates in N docs:

1. Define function  $f$  capturing contents of each document in one number  $O(|doc|)$   
"Hash function", "signature", "fingerprint"
2. Create < $f(doc)$ , ID of doc> pairs  $O(\sum_{i=1}^N (|doc_i|))$
3. Sort the pairs  $O(N \log N)$
4. Recognize duplicate or near-duplicate documents as having the same  $f$  value or  $f$  values within a small *threshold*  $O(N)$

Compare: computing a similarity score on pairs of documents

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### Syntactic Clustering Paradigm

- Does compare docs, so not same as paradigm we started with, but uses ideas
- Contents of doc captured by sketch – a set of shingle values
- Similarity of docs scored by count of common shingle values for docs
- Don't look at all doc pairs, look at all doc pairs that share a shingle value
- Text clusters by similarity threshold

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## More efficient : supershingles

### “meta-sketch”

1. Sort shingle values of a sketch
2. Compute the shingling of the sequence of shingle values
  - Each original shingle value now a token
  - Gives “supershingles”
3. “meta-sketch” = set of supershingles

**One supershingle in common =>  
sequences of shingles in common  
Documents with  $\geq 1$  supershingle in common => similar**

- Each supershingle for a doc. characterizes the doc
- Sort <supershingle, docID> pairs: docs sharing a supershingle are similar => our first paradigm

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## Pros and Cons of Supershingles

- + Faster
- Problems with small documents – not enough shingles
- Can't do containment
  - Shingles of superset that are not in subset break up sequence of shingle values

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## Using with Web Crawling

- Want know if new doc. too similar to ones seen
- What this calculation look like?

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## Using with Web Crawling

- Want know if new doc. too similar to ones seen
- No clustering required
- calculate sketch or supershingle of new document
- Look up to see if have similar document
  - or similar document that is fresh enough
  - Need efficient look-up

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## Variations of shingling

- Can define different ways to do sampling
- Studies in original paper used modular arithmetic
  - sketch formed by taking shingle hash values mod some selected m

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## Original experiments (1996) by Broder et. al.

- 30 million HTML and text docs (150GB) from Web crawl
- 10-word shingles
- 600 million shingles (3GB)
- 40-bit shingle “fingerprints”
- Sketch using 4% shingles (variation of alg. we've seen)
- Used count of shingles for similarity
- Using threshold  $t = 50\%$ , found
  - 3.6 million clusters of 12.3 million docs
  - 2.1 million clusters of identical docs – 5.3 million docs
  - remaining 1.5 million clusters mixture:  
“exact duplicates and similar”

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