



Procedure Call

CS 217



Procedure Call

- Involves following actions
 - pass arguments
 - save a return address
 - transfer control to callee
 - transfer control back to caller
 - return results

```
int add(int a, int b)
{
    return a + b;
}

int main()
{
    int c = add(3, 4);
    printf("%d\n", c);
    return 0;
}
```

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Jump Instruction



`jmp1 address, reg`

10	reg	111000	rs1	0	0	rs2
10	reg	111000	rs1	1	simm13	
31	29	24	18	13	12	4

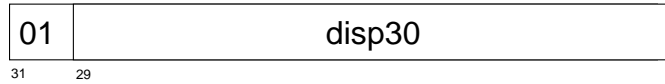
leaves PC in *reg*

```
reg = PC; /* return address */
PC = nPC;
nPC = rs1 + op2;
```

Call Instruction



`call label`



leaves PC (location of `call`) in `%o7` (`%r15`)

```
%o7 = PC; /* return address */
PC = nPC;
nPC = PC + sign_extend(disps30) << 2;
```

Like: `jmp1 label, %o7`

Procedure Call Example



- Simplest example: leaf procedure (`c=a+b`)

```
ld  a,%o0      ld  a,%o0
ld  b,%o1      call _add
call _add      ld  b,%o1
nop           st  %o0,c
st  %o0,c
```

Return Instruction



```
jmp1 %o7+8,%g0
```

- transfers control from caller to callee
- synthetic instructions: `ret` and `ret1`
- why +8?

```
_add: save %sp, ..., %sp
      add %o0, %o1, %o1
      ret
      restore
```

Calls with Function Pointers



```
jmp1 reg,%o7
```

- jumps to the 32-bit address specified in *reg*
- leaves PC (return address) in %o7 (%r15)
- example: `c = (*apply)(a,b);`

```
ld    b,%o0
ld    c,%o1
ld    apply,%o3
jmp1  %o3,%o7; nop
st    %o0,a
```

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Procedure Call

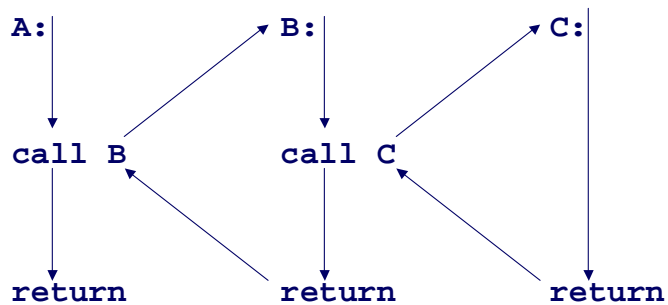


- Requirements
 - Pass a variable number of arguments
 - Pass and return structures
 - Allocate and deallocate space for local variables
 - Save and restore caller's registers
 - Handle nested procedure calls
- Entry and exit sequences collaborate to implement these requirements

Nested/Recursive Calls



- A calls B, which calls C



Must work when B is A

Arguments and Return Values



- By convention
 - caller places arguments in the “out” registers
 - callee finds its arguments in the “in” registers
 - only the first 6 arguments are passed in registers
 - the rest are passed on the stack

Arguments and Return Value (cont)

- Registers at call time

<u>caller</u>	<u>callee</u>	
%o7	%i7	return address -8 (%r15)
%o6	%i6	stack/frame pointer (%r14)
%o5	%i5	sixth argument
...	...	
%o0	%i0	first argument

Arguments and Return Value (cont)

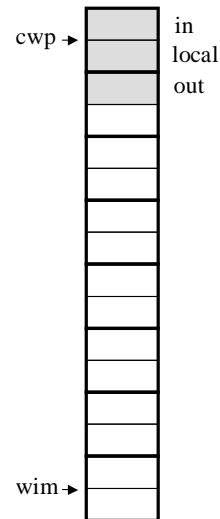
- Registers at return time

<u>caller</u>	<u>callee</u>	
%o5	%i5	sixth return value
%o4	%i4	fifth return value
...	...	
%o0	%i0	first return value

Register Windows



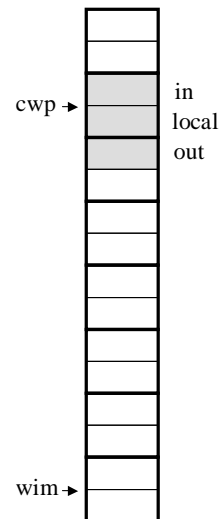
- Machine has more than 32 registers
 - Each procedure gets 16 “new” registers
 - All procedures can use globals
- The window “slides” at call time
 - caller’s out registers become callee’s in registers
- Instructions
 - **save** slides the window forward
 - **restore** slides the window backwards
 - decrement/increments CWP register
- Finite number of windows (usually 8)



Register Windows



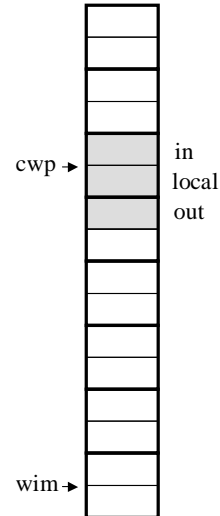
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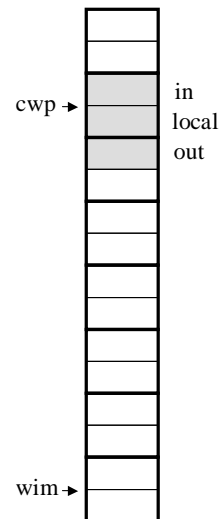
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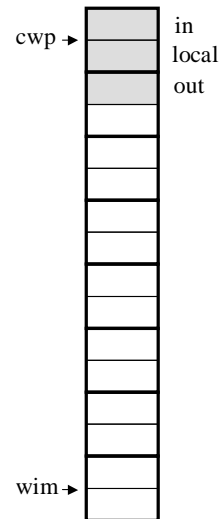
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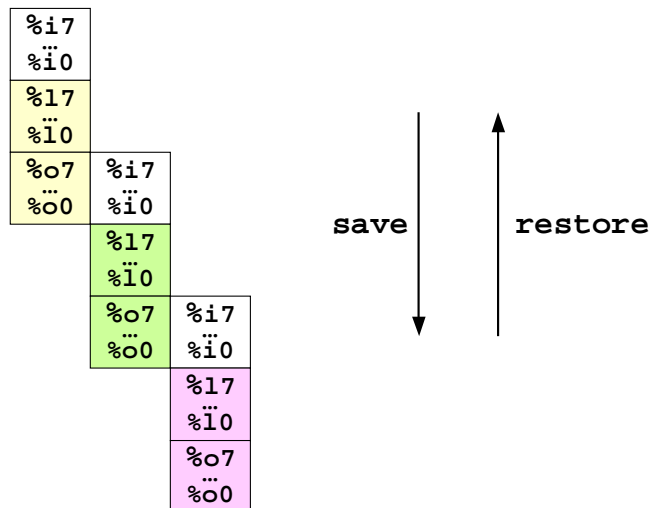
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Register Windows (cont)



Window Management



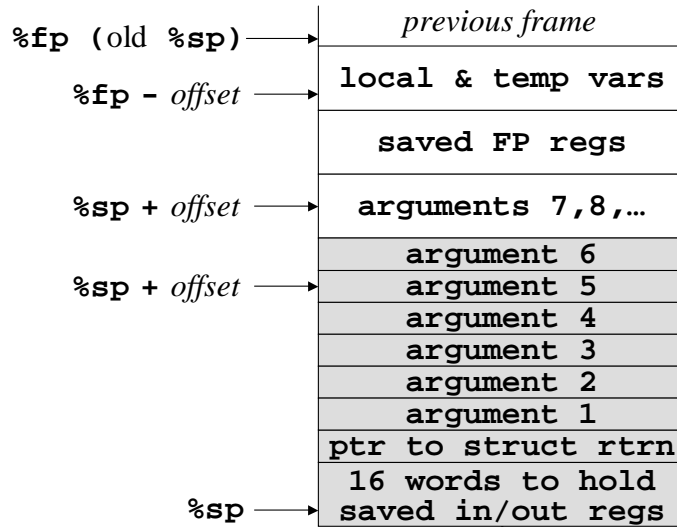
- Call time (`save`)
 - `save %sp,N,%sp`
 - current window becomes previous window
 - decrements CWP and checks for overflow
 - adds N to the stack pointer (allocates N bytes if $N < 0$)
 - if overflow occurs, save registers on the stack (must be enough stack space)
- Return time (`restore`)
 - previous window becomes current window
 - increments CWP and checks for underflow

Stack



- Procedure call information stored on stack
 - locals, including compiler temporaries
 - caller's registers, if necessary
 - callee's arguments, if necessary
- Sparc's stack grows "down" from high to low address
- The stack pointer (`%sp`) points to top word on the stack (must be multiple of 8)

Stack Frame



Example Stack Frames



```
main() {
    t(1,2,3,4,5,6,7,8);
}

t(int a1, int a2, int a3, int a4,
  int a5, int a6, int a7, int a8) {
    int b1 = a1;
    return s(b1,a8);
}

s(int c1, int c2) {
    return c1 + c2;
}
```

Example (cont)



```
_main: save %sp,-104,%sp
       set 1,%o0
       set 2,%o1
       set 3,%o2
       set 4,%o3
       set 5,%o4
       set 6,%o5
       set 7,%i5
       st %i5,[%sp+4*6+68]
       set 8,%i5
       st %i5,[%sp+4*7+68]
       call _t; nop
       ret; restore
```

Example (cont)



```
_t: save %sp,-96,%sp
    st %i0,[%fp-4]
    ld [%fp-4],%o0
    ld [%fp+96],%o1
    call _s; nop
    mov %o0,%i0
    ret; restore

_s: add %o0,%o1,%o0
    retl; nop
```

