

Lecture 18. Elementary Systems Programming

- **Software tools** are programs that **manipulate programs**, each in potentially different **languages**

	Name	Input	Output
Macro preprocessor	cpp	C	C
Compiler	rcc	C	assembly code
Assembler	as	assembly code	object code
Linker	ld -r	object code, libraries	object code
Loader	ld	object code	executable code
Operating system	UNIX	executable code	

- ‘Driver’ programs, like `lcc`, hide many of these steps

```
% lcc -v hello.c
/usr/local/lcc/lib/cpp ... hello.c hello.i
/usr/local/lcc/lib/rcc -target=sparc-solaris hello.i hello.s
/bin/as -o hello.o hello.s
/bin/ld -o a.out ... hello.o -lm -lc
% a.out
Hello world!
```

Compilation Pipeline

```
% cat hello.c
/* Everyone's first
   C program. */
#include <stdio.h>

int main(void) {
    printf("Hello world!\n");
    return 0;
}
```

- The macro **preprocessor** strips comments, expands macro definitions, processes conditional compilation directives, and injects include files

```
% lcc -E hello.c >hello.i; cat hello.i
#line 1 "hello.c"
...
#line 1 "/usr/local/lib/lcc/include/stdio.h"
...
extern int printf(const char *, ...);
...
#line 4 "hello.c"

int main(void) {
    printf("Hello world!\n");
    return 0;
}
```

See Chap. 13 in Deitel and Deitel for details

Compilation and Assembly

- The **compiler** translates C to symbolic assembly language, à symbolic TOY instructions

```
% lcc -S hello.i; cat hello.s
...
_main:
save %sp,-96,%sp
set L2,%o0
call _printf; nop
mov %g0,%i0
L1:
ret; restore
L2: ...
```

- The **assembler** translates symbolic assembly language to relocatable object code, à TOY instruction encodings

```
% lcc -c hello.s; dis hello.o
0:  9d e3 bf a0      save    %sp, -96, %sp
4:  11 00 00 00      sethi   %hi(printf), %o0
8:  90 12 20 00      or     %o0, printf, %o0
c:  40 00 00 00      call   0xc
10: 01 00 00 00      nop
14: b0 10 00 00      clr    %i0
18: 81 c7 e0 08      ret
1c: 81 e8 00 00      restore
```

Linking

- The *linker* combines object code files and libraries in a new object code file

```
% ld -r -o foo.o hello.o -lc; dis foo.o
main      0:  9d e3 bf a0      save    %sp, -96, %sp
          4:  11 00 00 00      sethi   %hi(.L350), %o0
          8:  90 12 20 00      or      %o0, .L350, %o0
          c:  40 00 00 00      call   (.L350+12)
          10: 01 00 00 00      nop
          14: b0 10 00 00      clr     %i0
          18: 81 c7 e0 08      ret
          1c: 81 e8 00 00      restore
printf    20: 9d e3 bf a0      save    %sp, -96, %sp
          24: 15 00 00 00      sethi   %hi(0x0), %o2
          28: f2 27 a0 48      st      %i1, [%fp + 72]
          2c: d4 02 a0 00      ld      [%o2], %o2
          30: 11 00 00 00      sethi   %hi(0x0), %o0
          34: f4 27 a0 4c      st      %i2, [%fp + 76]
```

...

Loading

- The ***loader*** translates object code to executable code

```
% ld foo.o; dis a.out
```

```
...
```

```
15148:  9d e3 bf a0      save    %sp, -96, %sp
1514c:  11 00 00 8a      sethi  %hi(0x22800), %o0
15150:  90 12 22 2c      or     %o0, 0x22c, %o0
15154:  40 00 00 05      call   0x15168
15158:  01 00 00 00      nop
1515c:  b0 10 00 00      clr   %i0
15160:  81 c7 e0 08      ret
15164:  81 e8 00 00      restore
15168:  9d e3 bf a0      save    %sp, -96, %sp
1516c:  13 00 00 e8      sethi  %hi(0x3a000), %o1
15170:  f2 27 a0 48      st     %i1, [%fp + 72]
15174:  d2 0a 62 b4      ldub  [%o1 + 692], %o1
15178:  f4 27 a0 4c      st     %i2, [%fp + 76]
```

```
...
```

- The operating system loads the executable code into memory and jumps to it

```
% a.out
```

```
Hello world!
```

Assembly Language

- An assembly language is a symbolic representation for machine language
 - Mnemonic names for opcodes and registers; usually terse
 - Symbolic names for addresses — data locations and jump ‘targets’
 - Easy to delete, insert, and rearrange instructions
- TAL: TOY Assembly Language

HALT		halt	
ADD	R,R ₁ ,R ₂	add	$R \leftarrow R_1 + R_2$
SUB	R,R ₁ ,R ₂	subtract	$R \leftarrow R_1 - R_2$
MUL	R,R ₁ ,R ₂	multiply	$R \leftarrow R_1 \times R_2$
XOR	R,R ₁ ,R ₂	exclusive OR	$R \leftarrow R_1 \wedge R_2$
AND	R,R ₁ ,R ₂	logical AND	$R \leftarrow R_1 \& R_2$
SHL	R,R ₁ ,R ₂	shift left	$R \leftarrow R_1 \ll R_2$
SHR	R,R ₁ ,R ₂	shift right	$R \leftarrow R_1 \gg R_2$
LI	R, <i>const8</i>	load immediate	$R \leftarrow \textit{const8}$ (8-bit constant)
LD	R,(R ₁ + <i>const8</i>)	load	$R \leftarrow M[R_1 + \textit{const8}]$
ST	R,(R ₁ + <i>const8</i>)	store	$M[R_1 + \textit{const8}] \leftarrow R$
SYS	R, <i>const8</i>	system call	system call <i>const8</i> , arg in R
J	<i>label</i>	jump	$PC \leftarrow \textit{label}$
JLT	R, <i>label</i>	jump if less	$PC \leftarrow \textit{label}$ if $R < 0$
JI	(R)	jump indirect	$PC \leftarrow R$
JAL	R, <i>label</i>	jump and link	$R \leftarrow PC, PC \leftarrow \textit{label}$

Programming in TAL

power.t: (see page 9-6)

```

POWER    LI    R4,1           initialize R4
         LI    R3,1           initialize z
LOOP     SUB   R2,R2,R4       decrement exponent
         JLT   R2,DONE        quit when done
         MUL   R3,R3,R1       set z to z*x
         J     LOOP           do it again
DONE     JI    (R5)           return to caller

```

main.t: computes $A^4 + B^5$ (see page 9-7)

```

MAIN     LI    R0,0           initialize R0
         LI    R1,A           load A into R1
         LD    R1,(R1+0)
         LI    R2,4           want  $A^4$ 
         JAL   R5,POWER       call POWER
         ADD   R6,R3,R0       copy  $A^4$  to R6
         LI    R1,B           do it again for  $B^5$ 
         LD    (R1+0)
         LI    R2,5
         JAL   R5,POWER
         ADD   R6,R6,R3       R6 now holds  $A^4 + B^5$ 
         SYS   R6,2           print it
         HALT
A        3
B        2

```

Object Code

- The assembler reads TAL and emits relocatable object code

power.o:

```

00:      B401      =POWER      initialize R4
01:      B301      initialize z
02:      2224      decrement exponent
03:      6206     +start address  quit when done
04:      3331      set z to z*x
05:      5002     +start address  do it again
06:      7500      return to caller

```

- Relocation information tells the linker

The definitions of symbols

How to adjust jump targets relative to the ultimate starting address of the module

Which symbols are defined in other separately compiled modules

- Object code is usually a compact, binary format, not text as suggested above

Object Code, cont'd

main.o:

```

00:      B100      =MAIN          initialize R0
01:      B100     +A             load A into R1
02:      9110
03:      B204
04:      8500     +POWER        call POWER
05:      1630
06:      B100     +B             do it again for B5
07:      9110
08:      B205
09:      8500     +POWER        call POWER
0A:      1663
0B:      4602
0C:      0000
0D:      0003      =A
0E:      0002      =B

```

- **Assemblers maintain symbol tables: Sets of (symbol,value) pairs used to map**

Mnemonics to values

LI → B_{16} , R6 → 6, ...

Labels to offsets

LOOP → 2_{16} , DONE → 6_{16} , POWER → 0, A → $0D_{16}$, ...

power.o symbol table:

(POWER, 0)

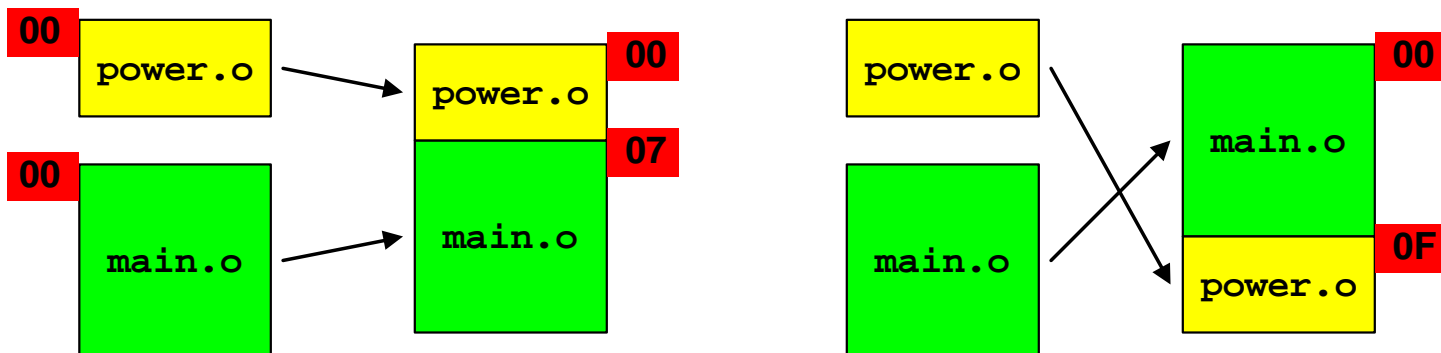
main.o symbol table:

(MAIN, 0) (A, $0D_{16}$) (POWER, ?) (B, $0E_{16}$)

Can implement symbol tables with binary trees

Linking

- The linker reads several object files and emits one relocatable object code file
 - Concatenates object code from input files
 - Relocates symbol definitions and instructions based on starting addresses in output object code



Resolves references to undefined symbols

Merges symbol tables

Linking, cont'd

- Linking `main.o` and `power.o` resolves references to `POWER`, `A`, `B`, and adjusts offsets to jump targets

```

00:      B100      =MAIN          initialize R0
01:      B10D      +start address load A into R1
02:      9110
03:      B204
04:      850F      +start address want A4
05:      1630      call POWER
06:      B10E      +start address copy A4 to R6
07:      9110      do it again for B5
08:      B205
09:      850F      +start address call POWER
0A:      1663      R6 now holds A4 + B5
0B:      4602      print it
0C:      0000
0D:      0003      =A
0E:      0002      =B

0F:      B401      =POWER          initialize R4
10:      B301      initialize z
11:      2224      decrement exponent
12:      6215      +start address quit when done
13:      3331      set z to z*x
14:      5011      +start address do it again
15:      7500      return to caller

```

Output includes relocation information for additional linking

Loading

- The loader reads a starting address and object code with no undefined symbols and emits executable code, adding in the starting address where necessary

```

20
20:      B100          initialize R0
21:      B12D          load A into R1
22:      9110
23:      B204          want A4
24:      852F          call POWER
25:      1630          copy A4 to R6
26:      B12E          do it again for B5
27:      9110
28:      B205
29:      852E          call POWER
2A:      1663          R6 now holds A4 + B5
2B:      4602          print it
2C:      0000
2D:      0003
2E:      0002

2F:      B401          initialize R4
30:      B301          initialize z
31:      2224          decrement exponent
32:      6235          quit when done
33:      3331          set z to z*x
34:      5031          do it again
35:      7500          return to caller

```

Separate Compilation

- A ***program*** is made up of many small ***modules***
 - A ‘few’ application-specific modules
 - ‘Many’ general-purpose modules, e.g., standard I/O functions like `printf`
- Compile general-purpose modules ***separately***, collect their object code in ***libraries***
- Compile application-specific modules separately, keep their object code
- To build a program
 1. Link together the application-specific object code modules
 2. Search the libraries for the general-purpose modules used by (1)
- Advantages
 - Avoid recompiling infrequently changed modules
 - Share libraries of well-tested general-purpose modules — don’t reinvent, reuse
- Designing and implementing general-purpose modules sounds easy, but it’s ***not***
 - Take COS 217, Introduction to Programming Systems
 - Read D. R. Hanson, *C Interfaces and Implementations: Techniques for Creating Reusable Software*, Addison-Wesley, 1997 (used in COS 217)