Consistency



COS 316: Principles of Computer System Design

Lecture 12

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Why Do We Build Systems?

- ...
- Abstract away complexity

Distributed Systems are Highly Complex Internally



Concurrent access by many client

Distributed Systems are Highly Complex Internally Sharding, Geo-Replication, Concurrency



Distributed Systems are Highly Complex Internally

Sharding, Geo-Replication, Concurrency



Consistency Models

- Contract between a (distributed) system and the applications that run on it
- A consistency model is a set of guarantees made by the distributed system

Stronger vs Weaker Consistency

Application Code

Strongly Consistent Distributed System

Application Code

Weakly Consistent Distributed System

Stronger vs Weaker Consistency

- Stronger consistency models
 - + Easier to write applications
 - System must hide many behaviors
- Fundamental tradeoffs between consistency & performance
 - (Discuss CAP, PRAM, SNOW in 418!)
- Weaker consistency models
 - Harder to write applications
 - Cannot (reasonably) write some applications
 - + System needs to hide few behaviors

Consistency Hierarchy

Linearizability Behaves like a single machine
Causal+ Consistency Everyone sees related
operations in the same order
Eventual Consistency Anything goes

Linearizability == " "Appears to be a Single Machine"

- External client submitting requests and getting responses from the system can't tell this is not a single machine!
- There is some total order over all operations
 - Processes all requests one by one
- Order preserves the real-time ordering between operations
 - If operation A completes before operation B begins, then A is ordered before B in real-time
 - If neither A nor B completes before the other begins, then there is no real-time order
 - (But there must be *some* total order)

Real-Time Ordering Examples



Real-Time Ordering Examples





w₁, w₂, r₂, w₃, r₃



W₁, **r**₁, **W**₂, **r**₂, **W**₃

Linearizable? P_A - w(x=1) -- w(x=2) - P_B - w(x=3) - \mathbf{P}_{C} -r(x)=2 -r(x)=3 -r(x)=3 \mathbf{P}_{D} -r(x)=1 -r(x)=2 -r(x)=2 \mathbf{P}_{D} r(x)=2 r(x)=2 \mathbf{P}_{D}



W₁, **W**₂, **r**₂, **r**₂, **W**₃



Linearizable? $P_A \rightarrow w(x=1)$ – w(x=2) – P_B - w(x=3) - P_{C} -r(x)=2 -r(x)=3 -r(x)=3 \mathbf{P}_{D} r(x)=1 r(x)=2 -1 \mathbf{P}_{D} -r(x)=2 -r(x)=2 -r(x)=2 P_{D} - r(x)=1 - r(x)=3 - 1 P_{D} r(x)=2 r(x)=1 r(x)=1 \mathbf{P}_{D}

Х

Linearizable? $P_A - w(x=1) -$ - w(x=2) - P_B – w(x=3) – Pc - w(x=4) - w(x=5) - w(x=5) \mathbf{P}_{D} ⊢ w(x=6) − P_E $| r(x)=2 - | r(x)=3 - | r(x)=6 - | r(x)=5 - | \sqrt{2}$ P_{F}

W₁, W₂, r₂, W₄, W₃, r₃, W₆, r₆, W₅, r₅

OR

 $W_1, W_4, W_2, r_2, W_3, r_3, W_6, r_6, W_5, r_5$

OR

w₁, w₂, r₂, w₃, r₃, w₄, w₆, r₆, w₅, r₅

Linearizable? P_A - w(x=1) -– w(x=2) – \mathbf{P}_{B} - w(x=3) - \mathbf{P}_{C} - w(x=4) - w(x=5) - w(x=5) P_{D} - w(x=6) - P_E | r(x)=2 - | r(x)=5 - | r(x)=6 - | r(x)=5 - | X $\mathbf{P}_{\mathbf{G}}$

Linearizable? P_A - w(x=1) -- w(x=2) - P_B – w(x=3) – Pc - w(x=4) - w(x=5) - w(x=5) P_{D} - w(x=6) - P_E $| r(x)=4 - | r(x)=2 - | r(x)=3 - | r(x)=6 - | \sqrt{2}$ P_H

 $W_1, W_4, r_4, W_2, r_2, W_3, r_3, W_5, W_6, r_6$

Linearizability == "Appears to be a Single Machine"

- There is some total order over all operations
 - Processes all requests one by one
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How to Provide Linearizability?

- 1. Use a single machine \bigcirc
- 2. Use "state-machine replication" on top of a consensus protocol like Paxos
 - Distributed system appears to be single machine that does not fail!!
 - Covered extensively in 418

3. ...

Consistency Hierarchy



Consistency Hierarchy



Causal+ Consistency Informally

- 1. Writes that are potentially causally related must be seen by everyone in the same order.
- 2. Concurrent writes may be seen in a different order by different entities.
 - Concurrent: Writes not causally related
- Potential causality: event a could have a causal effect on event b.
 - Think: is there a path of information from *a* to *b*?
 - a and b done by the same entity (e.g., me)
 - *a* is a write and *b* is a read of that write
 - + transitivity

Causal+ Sufficient



Causal+ Sufficient













Causal+ Not Sufficient

(Need Linearizability)

- Need a total order of operations
 - e.g., Alice's bank account ≥ 0
- Need a real-time ordering of operations
 - e.g., Alice changes her password, Bob cannot login with old password

Consistency Hierarchy



Eventual Consistency

- Anything goes for now...
 - (If updates stop, eventually all copies of the data are the same)
- But, eventually consistent systems often try to provide consistency and often do
 - e.g., Facebook's TAO system provided linearizable results 99.9994% of the time [Lu et al. SOSP '15]
- "Good enough" sometimes
 - e.g., 99 vs 100 likes

Consistency Model Summary

- Consistency model specifies strength of abstraction
 - Linearizability \rightarrow Causal+ \rightarrow Eventual
 - Stronger hides more, but has worse performance
- When building an application, what do you need?
 - Select system(s) with necessary consistency
 - Always safe to pick stronger
- When building a system, what are your guarantees?
 - Must design system such that they always hold
 - Must confront fundamental tradeoffs with performance
 - What is more important?