

# Class Meeting: Lectures 21-23

COS 461: Computer Networks

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# Today

- Network Security
- Datacenter Networks
- Course Summary & Wrap-Up

# BGP Security Today

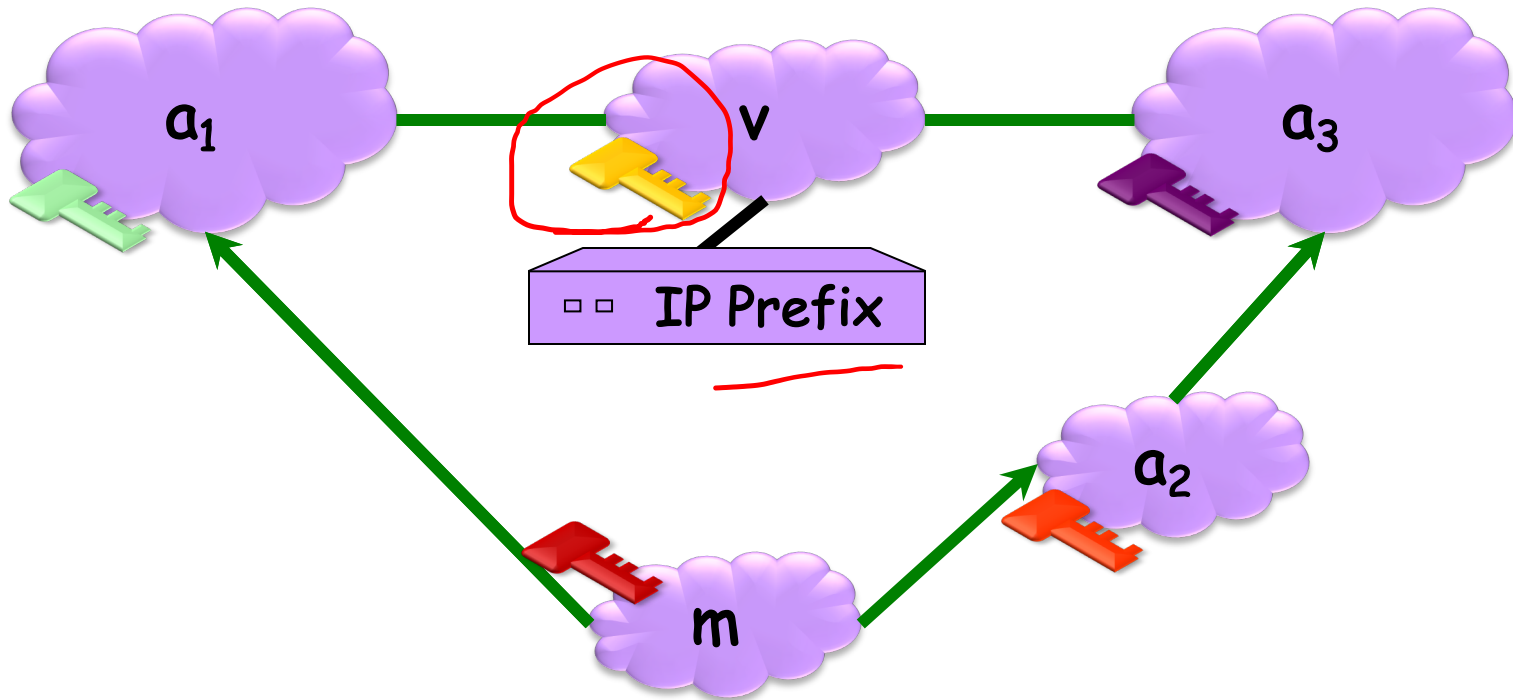
- Applying “best common practices”
  - Securing the session (authentication, encryption)
  - Filtering routes by prefix and AS path
  - Packet filters to block unexpected control traffic
- This is not good enough
  - Depends on vigilant application of practices
  - Doesn't address fundamental problems
    - Can't tell who owns the IP address block
    - Can't tell if the AS path is bogus or invalid
    - Can't be sure data packets follow the chosen route

# Proposed Enhancements to BGP

# Secure BGP



Origin Authentication + cryptographic signatures



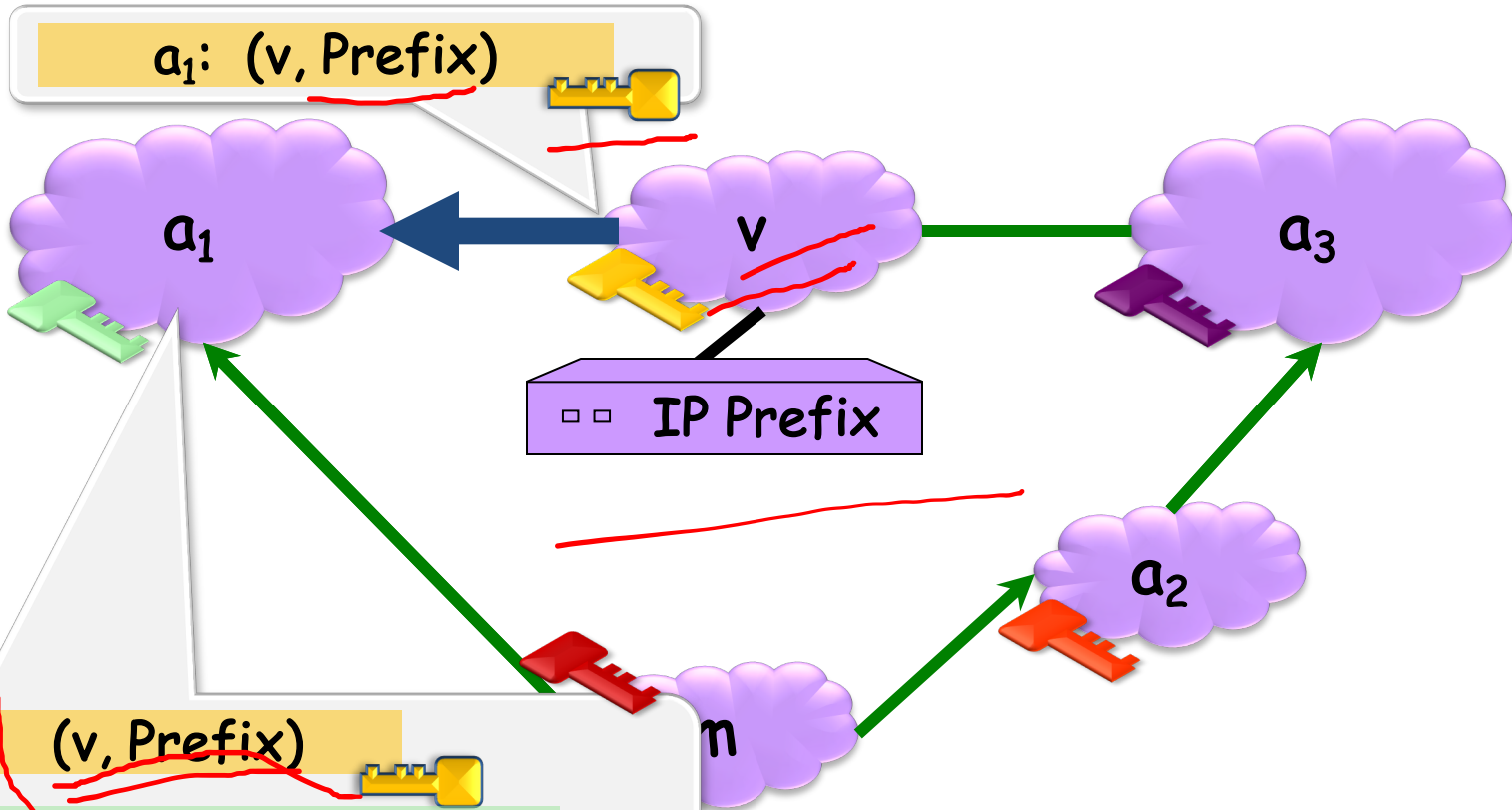
**Public Key Signature:** Anyone who knows  $v$ 's public key can verify that the message was sent by  $v$ .



# Secure BGP



Origin Authentication + cryptographic signatures



...one who knows  $v$ 's public key can verify that the message was sent by  $v$ .



# "Secure BGP"

- **Route attestations**

- Distributed as an attribute in BGP update message
- Signed by each AS as route traverses the network

- **Address attestations**

- Claim the right to originate a prefix
- Signed and distributed out-of-band
- Checked through delegation chain from ICANN

- **S-BGP can validate**

- AS path indicates the order ASes were traversed
- No intermediate ASes were added or removed
- Proper ASes originate prefixes

# S-BGP Deployment Challenges

- Complete, accurate registries of prefix "owner"
- Public Key Infrastructure
  - To know the public key for any given AS
- Cryptographic operations
  - E.g., digital signatures on BGP messages
- Need to perform operations quickly
  - To avoid delaying response to routing changes
- Difficulty of incremental deployment
  - Hard to have a "flag day" to deploy S-BGP



# Detecting Suspicious Routes

- **Monitoring BGP update messages**
  - Use past history as an implicit registry
- **E.g., AS that announces each address block**
  - Prefix 18.0.0.0/8 usually originated by AS 3
- **E.g., AS-level edges and paths**
  - Never seen the subpath "7018 88 1785"
- **Out-of-band detection mechanism**
  - Generate reports and alerts
  - Internet Alert Registry: <http://iar.cs.unm.edu/>
  - Prefix Hijack Alert System: <http://phas.netsec.colostate.edu/>

# BGP Security: Conclusions

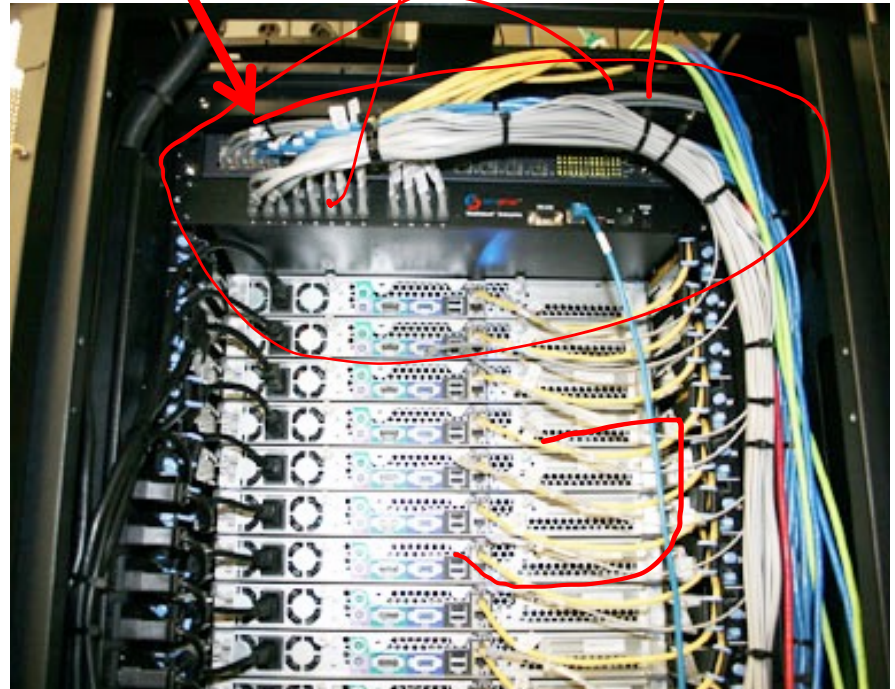
- Internet protocols designed based on trust
  - Insiders are good actors, bad actors on the outside
- Border Gateway Protocol is very vulnerable
  - Glue that holds the Internet together
  - Hard for an AS to locally identify bogus routes
  - Attacks can have serious global consequences
- Proposed solutions/approaches
  - Secure variants of the Border Gateway Protocol
  - Anomaly detection, with automated response
  - Broader focus on data-plane availability

# Today

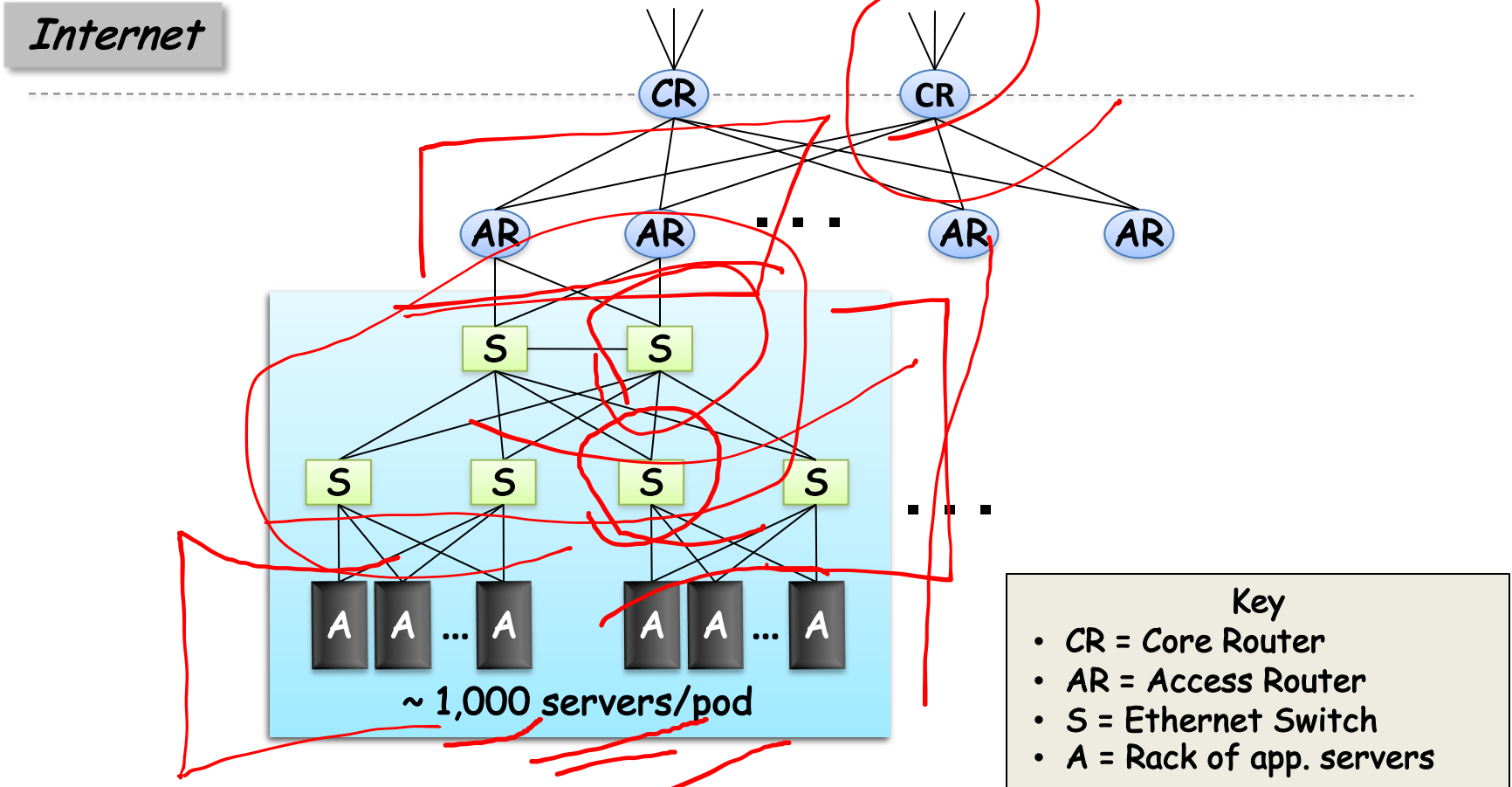
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# Top-of-Rack Architecture

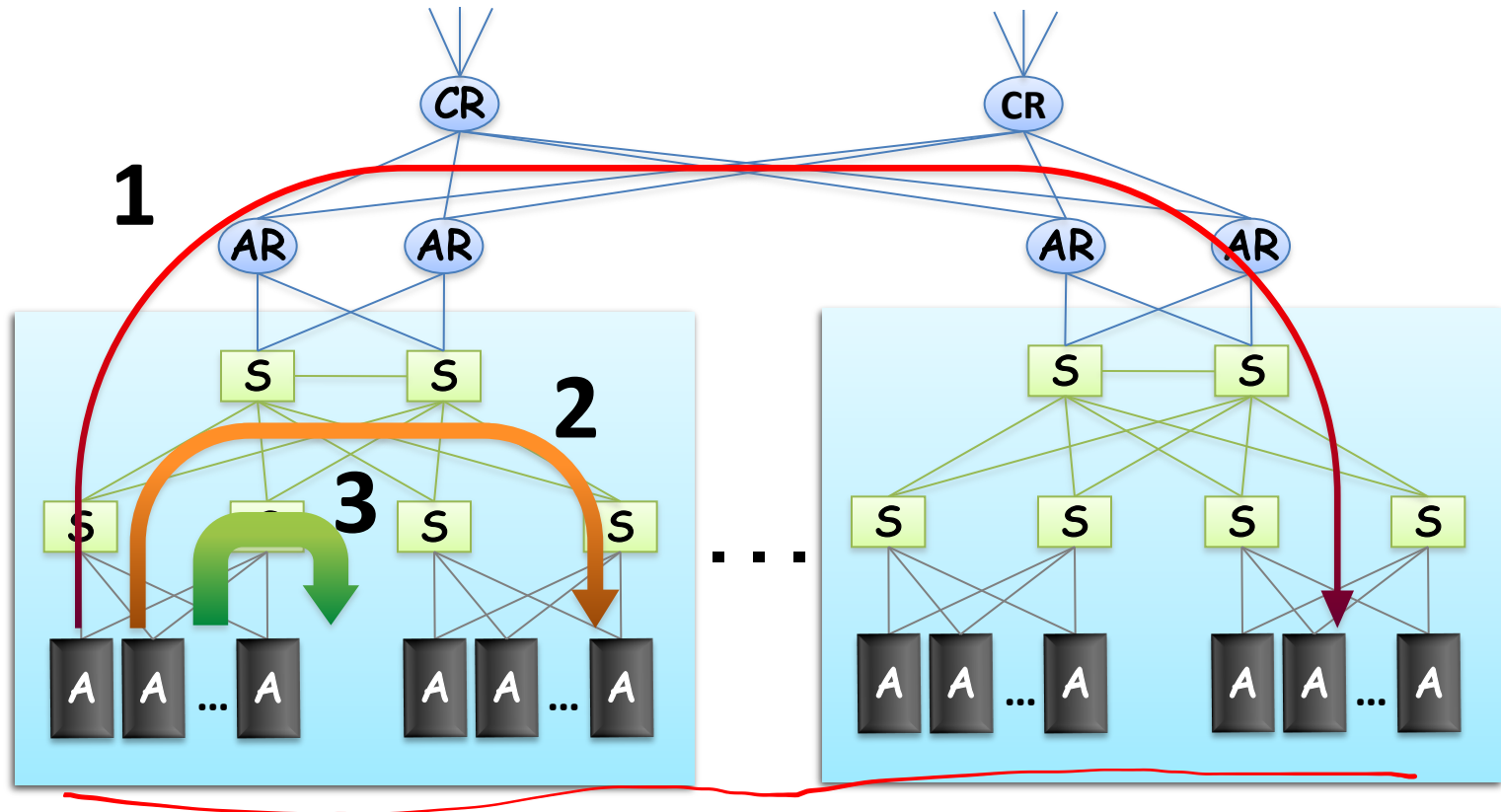
- **Rack of servers**
  - Commodity servers
  - And top-of-rack switch
- **Modular design**
  - Preconfigured racks
  - Power, network, and storage cabling



# Datacenter Network Topology

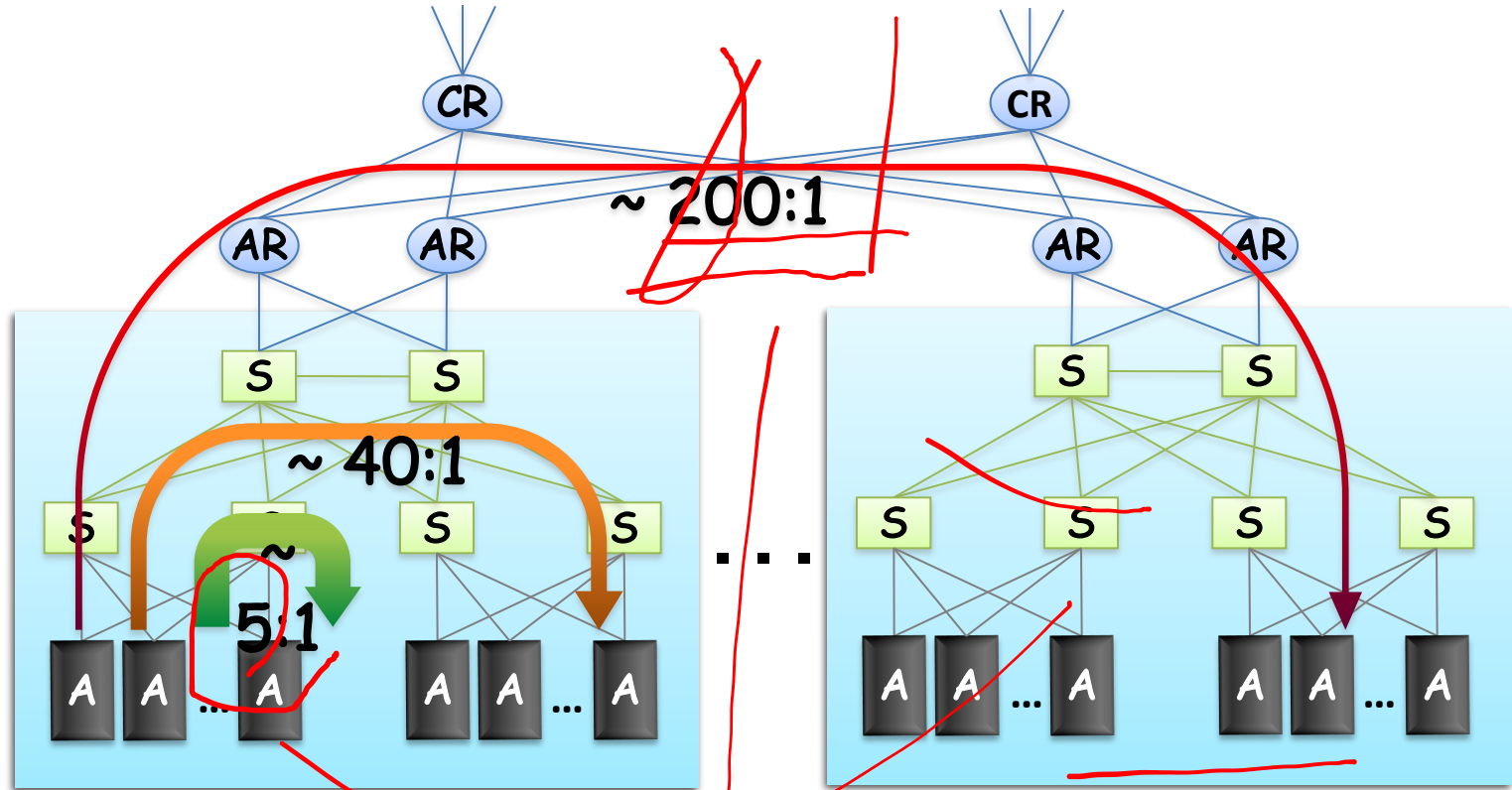


# Capacity Mismatch?

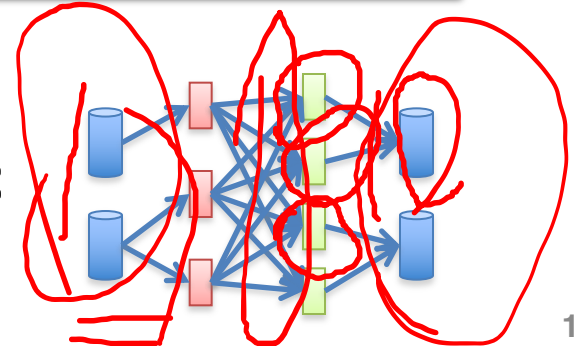


**“Oversubscription”:**  
Much more demand vs. supply for higher links

# Capacity Mismatch!



Particularly bad for east-west traffic

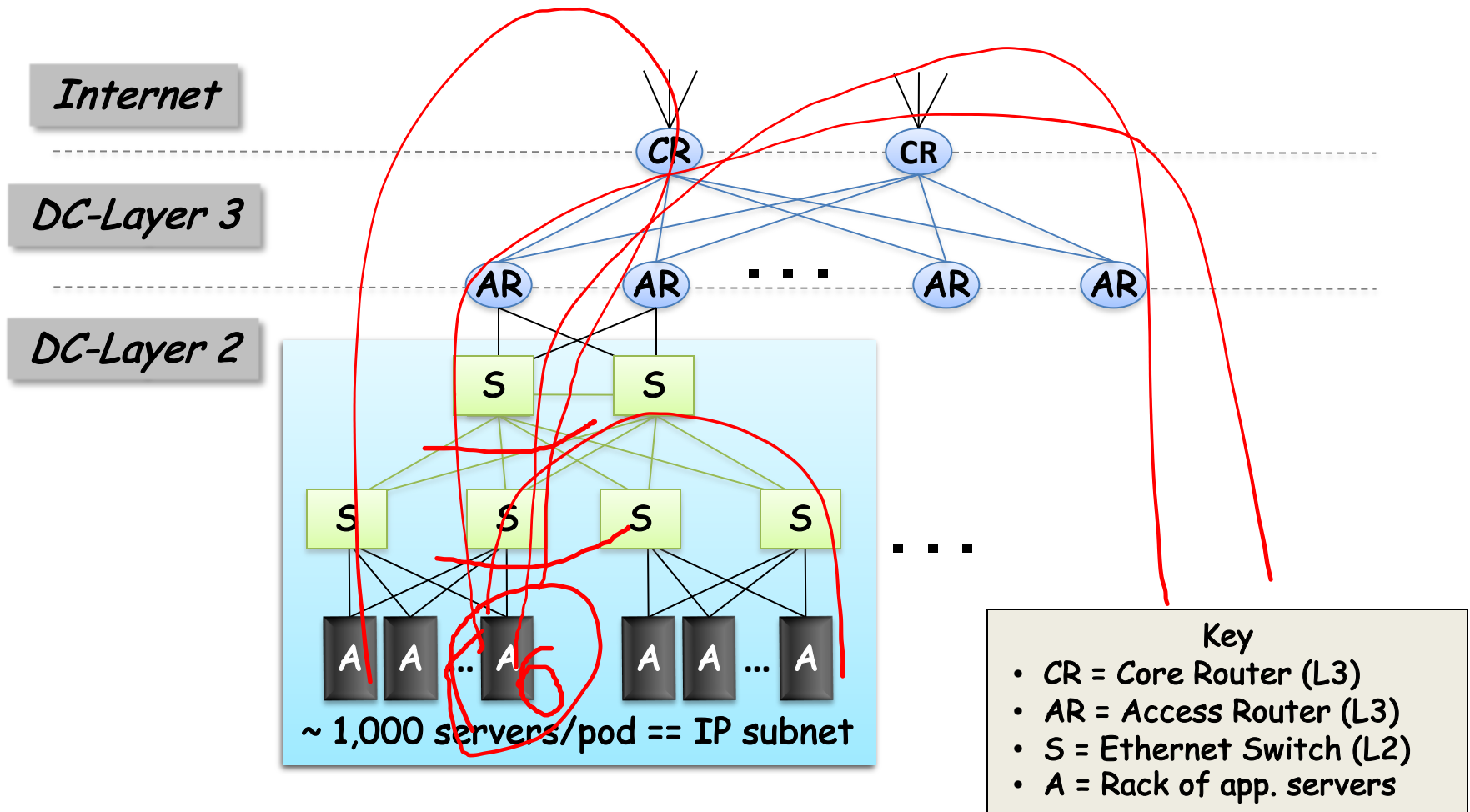


# Layer 2 vs. Layer 3?

- Ethernet switching (layer 2)
  - Cheaper switch equipment
  - Fixed addresses and auto-configuration
  - Seamless mobility, migration, and failover
- IP routing (layer 3)
  - Scalability through hierarchical addressing
  - Efficiency through shortest-path routing
  - Multipath routing through equal-cost multipath

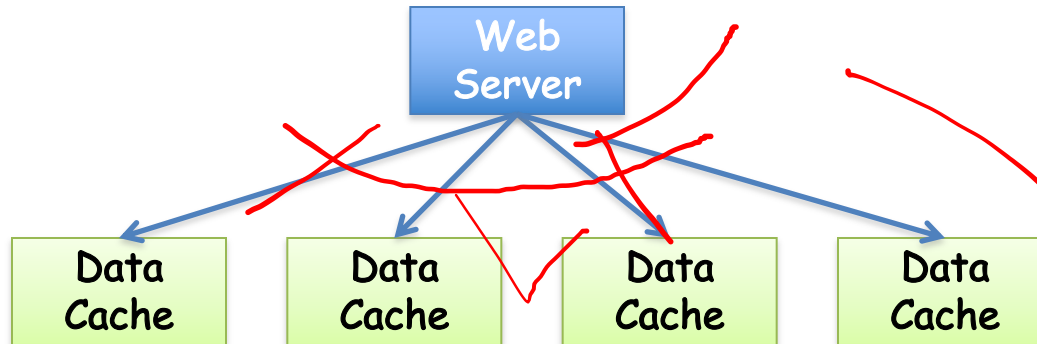


# Datacenter Routing



New datacenter networking  
problems have emerged...

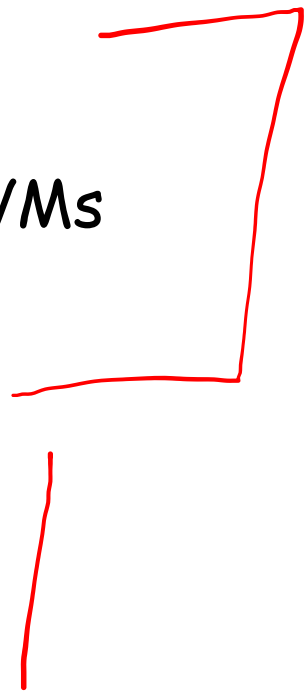
# Network *Incast* Problem



- **Incast arises from synchronized, parallel requests**
  - Web server sends out parallel request (“which friends of Johnny are online?”)
  - Nodes reply at same time, cause traffic burst
  - Replies potentially exceed switch’s buffer, causing drops

# Data Center Networks: Summary

- **Cloud computing**
  - Major trend in IT industry
  - Today's equivalent of factories
- **Datacenter networking**
  - Regular topologies interconnecting VMs
  - Mix of Ethernet and IP networking
- **Modular, multi-tier applications**
  - New ways of building applications
  - New performance challenges



# Today

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# Some Key Concepts

- Course was organized around protocols
  - But a small set of concepts recur in many protocols
- General CS concepts
  - Hierarchy, indirection, caching, randomization
- Networking-specific concepts
  - Soft state, layering, (de)multiplexing
  - End-to-end argument

# Hierarchy

- Scalability of large systems
  - Cannot store all information everywhere
  - Cannot centrally coordinate everything
- Hierarchy to manage scale
  - Divide system into smaller pieces
- Hierarchy to divide control
  - Decentralized management
- Examples from the Internet
  - IP addresses, routing protocols, DNS, P2P

# Indirection

- Referencing by name
  - Rather than the value itself
  - E.g., manipulating a variable through a pointer
- Benefits of indirection
  - Human convenience
  - Reducing overhead when things change
- Examples of indirection in the Internet
  - ~~– Names vs. addresses~~
  - Mobile IP



# Caching

- Duplicating data stored elsewhere
  - To reduce latency for accessing the data
  - To reduce resources consumed
- Caching is often quite effective
  - Speed difference between cache and primary copy
  - Locality of reference, and small set of popular data
- Examples from the Internet
  - DNS caching, Web caching, CDNs

# Randomization

- Distributed adaptive algorithms
  - Multiple distributed parties
  - Adapting independently
- Risk of synchronization
  - Many parties reacting at the same time
  - Leading to bad aggregate behavior
- Randomization can desynchronize
  - Ethernet back-off
- Rather than imposing centralized control

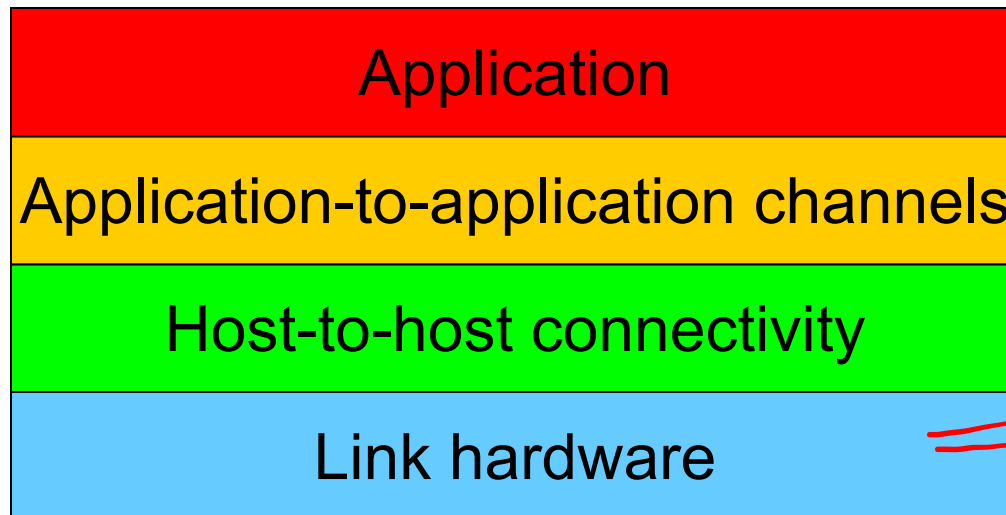
# Soft State

- **State: stored in nodes by network protocols**
  - Installed by receiver of a set-up message
  - Updated when conditions change
- **Hard state: valid unless told otherwise**
  - Removed by receiver of tear-down message
  - Requires error handling to deal with sender failure
- **Soft state: invalid if not told to refresh**
  - Periodically refreshed, removed by timeout
- **Soft state reduces complexity**
  - DNS caching, DHCP leases



# Layering: A Modular Approach

- **Sub-divide the problem**
  - Each layer relies on services from layer below
  - Each layer exports services to layer above
- **Interface between layers defines interaction**
  - Hides implementation details
  - Layers can change without disturbing other layers



# Power at the End Host

## End-to-End Principle

Whenever possible, communications protocol operations should be defined to occur at the **end-points** of a communications system.

## Programmability

With programmable end hosts, new network services can be added at **any time, by anyone**.

# The Internet of the Future

- Can we fix what ails the Internet
  - Security, performance, reliability
  - Upgradability, managability
- Without throwing out baby with bathwater
  - Ease of adding new hosts
  - Ease of adding new services
  - Ease of adding new link technologies
- An open technical and policy question...



# Final Exam

- Begins 9:00 AM on Wednesday, December 15.
- The exam will be due at 5:00 PM on Monday, December 20
  - Six hour Gradescope completion time limit
- online, open-book, open-461 material, calculators-allowed