





# Class Meeting: Lectures 17 and 18: Wireless, SDN

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[Various parts adapted from B. Karp, N. McKeown, J. Rexford]

# Today

- Wireless Networks
	- What makes wireless networks hard?
	- ALOHA: taking turns
	- MACA: sensing other transmissions

• Programmable Networks



## Wireless Links

- Interference / bit errors – More sources of corruption vs wired
- Multipath propagation – Signal does not travel in a straight line
- (Often) a *broadcast* medium – All traffic to everyone nearby
- Power trade-offs
	- Important for mobile, battery-powered devices

# Dealing With Bit Errors

- Wireless vs. wired links
	- Wired: most loss is due to queuing congestion
	- Wireless: higher, time-varying bit-error rate
- Dealing with high bit-error rates
	- Sender could increase transmission power
		- More interference with other senders
	- Stronger error detection and recovery
		- More powerful error detection/correction codes
		- Link-layer retransmission of corrupted frames

#### Wireless Broadcast and Interference: Interference matters at the receiver



A and B hear each other... B and C hear each other But, A and C do not So, A and C are unaware of their interference at B

#### Wireless LANs: a Timeline



#### ALOHAnet: Context

- Norm Abramson, 1970 at the University of Hawaii
	- Seven campuses, on four islands
	- Wanted to connect campus terminals and mainframe
	- Telephone costs high, so built a packet radio network



#### An Unslotted ALOHA Network



- Suppose: Chance new packet in time  $\Delta t$ :  $\Lambda \times \Delta t$ – N senders in total, send frames of time duration 1
- Then: *A* frames/sec *aggregate rate* from all N senders – Individual rate λ/N for each sender
- Collision and loss of data if the frames overlap (even a bit!)

#### Medium Access Control Refinement: "Slotted ALOHA"

- Divide time into slots of duration 1, synchronize so that nodes transmit only in a slot
	- Each of Nnodes transmits w/prob. p in each slot
	- So total transmission rate  $\lambda$  =  $\mathcal{N} \times p$
- As before, if exactly one transmission in slot, can receive; if two or more in slot, no one can receive (collision)



#### ALOHA Medium Access Control: Timeslots Double Throughput!



Just by forcing nodes to transmit on slot boundaries, we double peak medium utilization!



#### Assumptions

- Uniform, circular radio propagation
	- Fixed transmit power, all same ranges
	- Equal interference and communication ranges

**Radios modeled as "conditionally connected" wires based on circular radio ranges**

• Def'n: Node is connected to other node iff other located within circular radio range:



#### MACA: Goals

#### • Goals

- Fairness in sharing of medium
- Efficiency (total bandwidth achieved)
- Reliability of data transfer at MAC layer

#### When Does Listen-Before-Talk Carrier Sense (CS) Work Well?

• Two pairs far away from each other – Neither sender carrier-senses the other



#### B transmits to A, while D transmits to C.

## When Does CS Work Well?

• Both transmitters can carrier sense each other

#### But what about cases in between these extremes?



B transmits to A, D transmits to C, taking turns.

# Hidden Terminal Problem  $A \longrightarrow X \longleftarrow c$

- C can't hear A, so C will transmit while A transmits – Result: Collision at B
- Carrier Sense insufficient to detect all transmissions on wireless networks!
- Key insight: Collisions are spatially located at receiver

# Exposed Terminal Problem  $A \leftarrow B$

- If C transmits, does it cause a collision at A? – Yet C cannot transmit while B transmits to A!
- Same insight: Collisions spatially located at receiver
- One possibility: directional antennas rather than omnidirectional. Why does this help? Why is it hard?

#### MACA: Multiple Access with Collision Avoidance

• Carrier sense became adopted in packet radio

• But distances (cell size) remained large

• Hidden and Exposed terminals abounded

• Simple solution: use receiver's medium state to determine transmitter behavior

## RTS/CTS

- Exchange of two short messages: Request to<br>Send (RTS) and Clear to Send (CTS)
- Algorithm
	- 1. A sends an RTS (tells B to prepare)
	- 2. B replies an CTS (echoes message length)
	- 3. A sends its Data

1. "RTS, k bits"  
\n
$$
A \overline{3. "Data"}
$$
\n2. "CTS, k bits"

#### Deference to CTS

- Hear  $crs \rightarrow$  Defer for length of expected data transmission time
	- Solves hidden terminal problem

A B C 1. "RTS, k bits" 2. "CTS, k bits" defers 3. "Data"

## Deference to RTS, but not CS

- Hear RTS  $\rightarrow$  Defer one CTS-time (why.?)
- MACA: No carrier sense before sending!
	- Karn concluded useless because of hidden terminals
- So **exposed** terminals **B, C** can transmit concurrently:

A B C 1. "RTS, k bits" 2. "CTS, k bits" 3. "Data" D (No deference after Step 2)

# Today

• Wireless Networks

- Programmable Networks
	- Division of labor: control, data planes
	- Software-defined networking
	- Examples

#### The Internet: A Remarkable Story

• Tremendous success

…

– From research experiment to global infrastructure



- Brilliance of under-specifying
	- Network: best-effort packet delivery
	- Hosts: arbitrary applications
- Enables innovation in applications
	- Web, P2P, VoIP, social networks, smart cars,
- But, change is easy only at the edge...  $\odot$  24

Inside the 'Net: A Different Story…

- Closed equipment
	- Software bundled with hardware
	- Vendor-specific interfaces
- Over specified
	- Slow protocol standardization
- Few people can innovate
	- Equipment vendors write the code
	- Long delays to introduce new features

**Impacts performance, security, reliability, cost…**

## Networks are Hard to Manage

- Operating a network is expensive – More than half the cost of a network – Yet, operator error causes most outages
- Buggy software in the equipment – Routers with 20+ million lines of code – Cascading failures, vulnerabilities, etc.
- The network is "in the way" – Especially in data centers and the hon-

Rethinking the "Division of Labor"

### Traditional Computer Networks



#### **Forward, filter, buffer, mark, rate-limit, and measure packets**

### Traditional Computer Networks



#### **Track topology changes, compute routes, install forwarding rules**

### Traditional Computer Networks

#### **Management plane: Human time scale**

**Collect measurements and configure the equipment**

## Remove that Control Plane!

- Simpler management
	- No need to "invert" control-plane operations
- Faster pace of innovation – Less dependence on vendors and standards
- Easier interoperability – Compatibility only in "wire" protocols
- Simpler, cheaper equipment
	- Minimal software

#### Software Defined Networking (SDN)



### Data Plane: Simple Packet Handling

- Simple packet-handling rules
	- Pattern: match packet header bits



- Actions: drop, forward, modify, send to controller
- Priority: disambiguate overlapping patterns
- Counters: #bytes and #packets



- **1.**  $src=1.2.*.*$ ,  $dest=3.4.5.*$   $\rightarrow$  drop
- **2.**  $src = *.*.*.*$ ,  $dest = 3.4.*.*$   $\rightarrow$   $forward(2)$
- **3.** src=10.1.2.3, dest= $*,*,*,* \rightarrow$  send to controller

# E.g.: Dynamic Access Control

traffic

- Inspect first packet of a connection
- Consult the access control policy
- Install rules to block or route

# E.g.: Seamless Mobility/Migration • See host send traffic at new location • Modify rules to reroute the traffic

## Summary

- Wireless networks: de facto means of accessing the Internet
	- Evolution from ALOHAnet, Ethernet, MACA, toward IEEE 802.11 Wi-Fi

• Software-Defined Networks: new ways of managing networks – New API, OpenFlow, enables new applications