





#### Class Meeting: Lectures 17 and 18: Wireless, SDN

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[Various parts adapted from B. Karp, N. McKeown, J. Rexford]

# Today

- Wireless Networks
  - What makes wireless networks hard?
  - ALOHA: taking turns
  - MACA: sensing other transmissions
- Programmable Networks



## Wireless Links

- Interference / bit errors
  - More sources of corruption vs wired
- Multipath propagation

   Signal does not travel in a straight line
- (Often) a broadcast medium
   All traffic to everyone nearby
- Power trade-offs
  - Important for mobile, battery-powered devices

## Dealing With Bit Errors

- Wireless vs. wired links
  - Wired: most loss is due to queuing congestion
  - Wireless: higher, time-varying bit-error rate
- Dealing with high bit-error rates
  - Sender could increase transmission power
    - More interference with other senders
  - Stronger error detection and recovery
    - More powerful error detection/correction codes
    - Link-layer retransmission of corrupted frames

#### Wireless Broadcast and Interference: Interference matters <u>at the receiver</u>



A and B hear each other... B and C hear each other But, A and C do not So, A and C are unaware of their interference <u>at B</u>

#### Wireless LANs: a Timeline



#### ALOHAnet: Context

- Norm Abramson, 1970 at the University of Hawaii
  - Seven campuses, on four islands
  - Wanted to connect campus terminals and mainframe
  - Telephone costs high, so built a packet radio network



#### An Unslotted ALOHA Network



- Suppose: Chance new packet in time  $\Delta t$ :  $\Lambda \times \Delta t$ – Nsenders in total, send frames of time duration 1
- Then: A frames/sec aggregate rate from all Nsenders
  - Individual rate N/N for each sender
- Collision and loss of data if the frames overlap (even a bit!)

#### Medium Access Control Refinement: "Slotted ALOHA"

- Divide time into slots of duration 1, synchronize so that nodes transmit only in a slot
  - Each of Nnodes transmits w/prob. p in each slot
  - So total transmission rate  $\Lambda = N \times p$
- As before, if exactly one transmission in slot, can receive; if two or more in slot, no one can receive (collision)



#### ALOHA Medium Access Control: <u>Timeslots</u> Double Throughput!



Just by forcing nodes to transmit on slot boundaries, we double peak medium utilization!



#### Assumptions

- Uniform, circular radio propagation
  - Fixed transmit power, all same ranges
  - Equal interference and communication ranges

Radios modeled as "conditionally connected" wires based on circular radio ranges

<u>Def'n</u>: Node is connected to other node *iff other located within* circular radio range:



#### MACA: Goals

#### • Goals

- Fairness in sharing of medium
- Efficiency (total bandwidth achieved)
- Reliability of data transfer at MAC layer

#### When Does Listen-Before-Talk *Carrier Sense* (CS) Work Well?

Two pairs far away from each other
 Neither sender carrier-senses the other



#### B transmits to A, while D transmits to C.

## When Does CS Work Well?

• Both transmitters can carrier sense each other

# But what about cases in between these extremes?



B transmits to A, D transmits to C, taking turns.

# Hidden Terminal Problem

- C can't hear A, so C will transmit while A transmits
   Result: Collision at B
- Carrier Sense insufficient to detect all transmissions on wireless networks!
- Key insight: Collisions are spatially located at receiver

# Exposed Terminal Problem

- If C transmits, does it cause a collision at A?
   Yet C cannot transmit while B transmits to A!
- Same insight: Collisions spatially located at receiver
- One possibility: directional antennas rather than omnidirectional. Why does this help? Why is it hard?

#### MACA: Multiple Access with Collision Avoidance

• Carrier sense became adopted in packet radio

• But distances (cell size) remained large

Hidden and Exposed terminals abounded

• Simple solution: use *receiver's* medium state to determine transmitter behavior

## RTS/CTS

- Exchange of two short messages: Request to Send (RTS) and Clear to Send (CTS)
- Algorithm
  - 1. A sends an RTS (tells B to prepare)
  - 2. B replies an CTS (echoes message length)
  - 3. A sends its Data

#### Deference to CTS

- Hear CTS → Defer for length of expected data transmission time
  - Solves hidden terminal problem

## Deference to RTS, but not CS

- Hear RTS → Defer one CTS-time (why?)
- MACA: No carrier sense before sending!
  - Karn concluded useless because of hidden terminals
- So exposed terminals B, C can transmit concurrently:

## Today

Wireless Networks

- Programmable Networks
  - Division of labor: control, data planes
  - Software-defined networking
  - Examples

#### The Internet: A Remarkable Story

- Tremendous success
  - From research experiment
     to global infrastructure



- Brilliance of under-specifying
  - Network: best-effort packet delivery
  - Hosts: arbitrary applications
- Enables innovation in applications
  - Web, P2P, VoIP, social networks, smart cars,
- But, change is easy only at the edge...  $\otimes$  <sup>24</sup>

Inside the 'Net: A Different Story ...

- Closed equipment
  - Software bundled with hardware
  - Vendor-specific interfaces
- Over specified
  - Slow protocol standardization-
- Few people can innovate
  - Equipment vendors write the code
  - -Long delays to introduce new features

Impacts performance, security, reliability, cost...

## Networks are Hard to Manage

- Operating a network is expensive

   More than half the cost of a network
   Yet, operator error causes most outages
- Buggy software in the equipment B – Routers with 20+ million lines of code – Cascading failures, vulnerabilities, etc.
- The network is "in the way"
   Especially in data centers and the hon

Rethinking the "Division of Labor"

### Traditional Computer Networks



# Forward, filter, buffer, mark, rate-limit, and measure packets

#### Traditional Computer Networks



# Track topology changes, compute routes, install forwarding rules

## Traditional Computer Networks

#### Management plane: Human time scale

Collect measurements and configure the equipment

## Remove that Control Plane!

- Simpler management
  - No need to "invert" control-plane operations
- Faster pace of innovation

   Less dependence on vendors and standards
- Easier interoperability
  - Compatibility only in "wire" protocols
- Simpler, cheaper equipment
  - Minimal software

#### Software Defined Networking (SDN)



### Data Plane: Simple Packet Handling

- Simple packet-handling rules
  - Pattern: match packet header bits



- Actions: drop, forward, modify, send to controller
- Priority: disambiguate overlapping patterns
- Counters: #bytes and #packets



- 1. src=1.2.\*.\*, dest=3.4.5.\* → drop
- 2. src = \*.\*.\*, dest=3.4.\*.\*  $\rightarrow$  forward(2)
- 3. src=10.1.2.3, dest=\*.\*.\*  $\rightarrow$  send to controller

## E.g.: Dynamic Access Control

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- Inspect first packet of a connection
- Consult the access control policy
- Install rules to block or route



## Summary

- Wireless networks: de facto means of accessing the Internet
  - Evolution from ALOHAnet, Ethernet, MACA, toward IEEE 802.11 Wi-Fi

 Software-Defined Networks: new ways of managing networks

 New API, OpenFlow, enables new applications