COS 461 Computer Networks

Class Meeting, Lectures 3 & 4

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Need to miss 461? Just reach out

Short-Term Remote Teaching and Other Contingency Measures

Under what circumstances might students or faculty need to miss class due to COVID?

Three COVID-related situations would require a student or faculty member to stay in their dorm or at home (as applicable) during the semester:

- 1. Anyone experiencing **COVID-like symptoms** should stay home/in their dorm room until they are tested and receive a negative test result.
- Anyone who receives a positive COVID test result will be required to self-isolate.
 Isolation will be a minimum of 10 days; isolation longer than 10 days is rare, but possible. Isolation means restricting activities outside of one's home or isolation dorm, except for obtaining medical care when necessary.

IP Protocol Stack: Key Abstractions

| Application | Applications | |
|-------------|--|----------|
| Transport | Reliable streams | Messages |
| Network | Best-effort global packet delivery | |
| Link | Best-effort <i>local</i> packet delivery | |

Best-Effort Global Packet Delivery

Circuit Switching (e.g., Phone Network)

- Source establishes connection
 - Reserve resources along hops in the path
- Source sends data

- Transmit data over the established connection

- Source tears down connection
 - Free the resources for future connections



Circuit Switching: Static Allocation

- Time-division
 - Each circuit allocated certain time slots
- Frequency-division
 - Each circuit allocated certain frequencies





Packet Switching

- Message divided into packets
 - Header identifies the destination address
- Packets travel separately through the network
 - Forwarding based on the destination address
 - Packets may be buffered temporarily
- Destination reconstructs the message



Is Best Effort Good Enough?

- Packet loss and delay
 Sender can resend
- Packet corruption
 - Receiver can detect, and sender can resend
- Out-of-order delivery
 - Receiver can put the data back in order

- Packets follow different paths
 - Doesn't matter
- Network failure
 - Drop the packet
- Network congestion

Drop the packet

Packet (Y) vs. Circuit Switching (A)?

- Predictable performance
- Network never blocks senders
- Reliable, in-order delivery
- Low delay to send data
- Simple forwarding
- No overhead for packet headers
- High utilization under most workloads
- No per-connection network state

Circuit

Packet

Circuit

Packet

Circuit

Circuit

Packet

Packet

Network Addresses

Grouping Related Hosts

- The Internet is an "inter-network"
 - Used to connect networks together, not hosts
 - Need to address a network (i.e., group of hosts)



- LAN = Local Area Network
- WAN = Wide Area Network

Scalability Challenge

- Suppose hosts had arbitrary addresses
 - Then every router would need a lot of information
 - ...to know how to direct packets toward every host



forwarding table

Hierarchical Addressing: IP Prefixes

- Network and host portions (left and right)
- 12.34.158.0/24 is a 24-bit prefix with 2⁸ addresses



Scalability Improved

- Number related hosts from a common subnet
 - 1.2.3.0/24 on the left LAN
 - 5.6.7.0/24 on the right LAN



Easy to Add New Hosts

- No need to update the routers
 - E.g., adding a new host 5.6.7.213 on the right
 - Doesn't require adding a new forwarding-table entry



Classless Inter-Domain Routing (CIDR)

• Use two 32-bit numbers to represent network:

Network number = IP address + Mask

IP Address : 12.4.0.0 IP Mask: 255.254.0.0



Hierarchical Address Allocation

- Hierarchy is key to scalability
 - Address allocated in contiguous chunks (prefixes)
 - Today, the Internet has about 600-800,000 prefixes



Packet Forwarding

Hop-by-Hop Packet Forwarding

- Each router has a forwarding table
 - Maps destination address to outgoing interface
- Upon receiving a packet
 - Inspect the destination address in the header
 - Index into the table
 - Determine the outgoing interface
 - Forward the packet out that interface
- Then, the next router in the path repeats

Separate Forwarding Entry Per Prefix

- Prefix-based forwarding
 - Map the destination address to matching prefix
 - Forward to the outgoing interface



CIDR Makes Packet Forwarding Harder

- Forwarding table may have many matches
 - E.g., entries for 201.10.0.0/21 and 201.10.6.0/23
 - The IP address 201.10.6.17 would match both!



Longest Prefix Match Forwarding

- Destination-based forwarding
 - Packet has a destination address
 - Router identifies longest-matching prefix
 - Cute algorithmic problem: very fast lookups



"Layer 2" Hubs and Switches

Bridges/Switches: Traffic Isolation

- Switch filters packets
 - Frame only forwarded to the necessary segments
 - Segments can support separate transmissions



Self Learning: Building the Table

- When a frame arrives
 - Inspect the *source* MAC address
 - Associate the address with the *incoming* interface
 - Store the mapping in the switch table
 - Use a timer to eventually forget the mapping



Self Learning: Handling Misses

- When frame arrives with unfamiliar destination
 - Forward the frame out all of the interfaces
 - ... except for the one where the frame arrived
 - Hopefully, this case won't happen very often!



Link Layer: Switches

- Typically connects individual computers

 A switch is essentially the same as a bridge
 ... though typically used to connect hosts
- Supports concurrent communication
 Host A can talk to C, while B talks to D



L3 Routers: Looking closer...

Basic Router Architecture

- Each switch/router has a forwarding table
 Maps destination address to outgoing interface
 - Maps destination address to outgoing interface

- Basic operation T/s
 - 1. Receive packet
 - 2. Look at header to determine destination address
 - 3. Look in forwarding table to determine output interface
 - 4. Modify packet header (e.g., decr TTL, update chksum)
 - 5. Send packet to output interface

Basic Router Architecture



- Basic operation
 - 1. Receive packet
 - 2. Look at header to determine destination address
 - 3. Look in forwarding table to determine output interface

Line Card (I/O)

- 4. Modify packet header (e.g., decr TTL, update chksum)
- 5. Send packet to output interface

Router



LPM: Motivation

- Each packet has destination IP address
- Router looks up table entry that matches address
- Benefits of CIDR allocation and LPM

Efficiency: Prefixes can be allocated at much finer granularity

Hierarchical aggregation: Upstream ISP can aggregate
 2 contiguous prefixes from downstream ISPs to
 shorter prefix

Decision: Crossbar switch



Decision: Crossbar switch Shared bus Only one input can speak to one output at a time

- Crossbar switch / switched backplane
 - Input / output pairs that don't compete can send in same timeslot

Crossbar switching

- Every input port has connection to every output port
- In each timeslot, each input connected to zero or more outputs



Crossbar switching

- Every input port has connection to every output port
- In each timeslot, each input connected to zero or more outputs



Problem: Head-of-line blocking

- Packet at front of queue blocks packets behind it from being processed
 - e.g.: 1st packet at Input 1 wants to go to Output 5;
 next packet at Input 1 that wants Output 4 is still blocked



Solution: *Virtual output queues*

- One queue at input, per output port (for all inputs)
- So avoids head-of-line blocking during crossbar scheduling



Processo Data, Control, & Management Planes Switching Fabric Data Control Management **Event** Time-Human Packet (ns) scale (10 ms to sec) (min to hours) Forwarding, buffering, Routing, Analysis, **Tasks** configuration filtering, signaling scheduling Line-card Router Humans or Location software hardware scripts

Coming Up in 461

Next Class Meeting

Lectures 5 (Transport Layer) and 6 (Congestion Control)

Precepts this Thursday and Friday