First Precept!

COS 318: Operating Systems

Precept Objectives:

In precepts we will usually cover two components:

- Outline of assignment, and what is necessary to know/do
- Understanding why the assignment is necessary & useful for Operating Systems

Today we will also be going over some of the high-level details of the course, and covering x86 asm, which is necessary for the 1st assignment.

Icebreaker

Because this may be your first in-person precepts in a while, it would be good to meet the other people interested in Operating Systems!

High Level Objectives of OS

Below, sorted by amount of time to learn & difficulty to achieve:

- 1. Be able to understand the various functions and structure of an OS
 - a. Better understand the interface w/ the system in user programs
- 2. Can read & understand existing OS source, given enough time
- 3. Can modify existing OSs (Linux Patches!)
- 4. Can develop new components/structure of an OS

- Google has Fuchsia & Android, Microsoft has Windows, Apple has MacOS/iOS
- Plenty of others, like VMWare which have OS adjacent services

Today's Material

- Timeline for Project 1
 - Also covering what is expected in the design review
- Review of x86 assembly (or introduction for those who used ARM) which should be sufficient for doing the 1st assignment



Project 1 Schedule

- Design Review: Mon(9/13) & Wed(9/15)
 - Sign up for 10-min slot from Mon (8:30pm-10:30pm) or Wed (3pm-7pm)
 - Complete set up and answer posted questions
- Project 1 Precept: Mon (9/13) & Tue (9/14),
 7:30pm 8:20pm. Due: Sun (9/26), 11:55pm



Design Review

- USBs will be given at the Design Review
- Write print_char and print_string assembly functions
- Be ready to describe:
 - How to move the kernel from disk to memory
 - How to create disk image
 - (More specific guidelines are provided on the project page)



x86 Assembly Tutorial

• x86/IA-32/i386 Assembly Overview

Components:

- Registers, Flags, Memory Addressing, Instructions, Stack / Calling Convention, Directives, Segments
- BIOS
- GDB

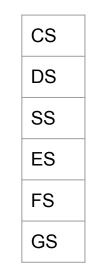
Registers

General Purpose Registers: 8,16,32 bits

Size (in Bits)

32	16	8
EAX	AX	AH/AL
EBX	BX	BH/BL
ECX	CX	CH/CL
EDX	DX	DH/DL
EDI	DI	DIL
ESI	SI	SIL
EBP	BP	BPL
ESP	SP	SPL

Segment Registers: 16 bits



Instruction Pointer (EIP): 32 bits Flags (EFLAGS): 32 bits







Flags are single bits in the EFLAGS register, where each bit has a different purpose such as:

- Controlling the behavior of CPU
- Saving the status of the last instruction
- More details at: <u>https://en.wikipedia.org/wiki/FLAGS_register</u>



Flags

• Status:

- CF: Carry flag
- ZF: Zero flag (op = 0)
- SF: Sign flag (op < 0)
- Control Flags:
 - IF: Interrupt flag (sti, cli), En/Disable interrupts



AT&T Syntax

- Instruction format: `instr src, dest`
 - `movw %ax,%bx`
- Prefix register names with %: %ax
- Prefix constants (immediate values) with \$
 - `movw \$1,%ax`
- Suffix instructions with size of data
 - b for byte (8 bit), w for word (16 bit), I for long (32 bit)
 - i.e. movb, movw, movl



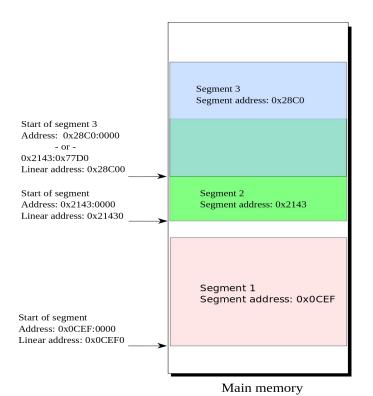
Memory Addressing in Real Mode

- 1MB memory
 - Valid address range: 0x00000-0xFFFF= 20 bits = 1MB
- 16-bit segments and 16-bit offsets into each segment: (segment << 4) + offset



Memory Addressing (Real Mode)

- Format (AT&T syntax):
 - segment:displacement(base,index,scale)
- Offset = Base + Index * Scale + Displacement
- Address = (Segment << 4) + Offset
- Displacement/Scale must be a constant w/o \$
- register that could be base: i.e. %bx, %bp
- Register that could be index: i.e. %si, %di
- Segment registers: %cs, %ds, %ss, %es, %fs, %gs





Instructions: Arithmetic & Logic

- add/sub{l,w,b} source,dest (dest source)
- inc/dec/neg{l,w,b} dest
- ource,dest (dest > source)
- and/or/xor{l,w,b} source,dest ...
- Restrictions
 - No more than one memory operand, other must be register



Instructions: Data Transfer

- mov{l,w,b} source, dest
- xchg{l,w,b} source, dest
- movsb/movsw
 - %es:(%di) \leftarrow %ds:(%si), must be these registers
 - Often used with %cx to move a number of bytes
 - movw \$0x10,%cx
 - rep movsw

• Segment registers can only address mem w/ registers



Stack Layout

- Grows from high to low
 - Lowest value address = "top" of stack
- %sp points to top of the stack
 - Used to reference temporary variables
- %bp points to bottom of stack frame
 - Used for local vars + function args.

(function arg. n)	
function arg. 1	06bp + 9
return address	• %bp + 8
old %ebp	• %bp + 4
local var. 1	• %bp
(local var. n)	%bp - 4
(callee-save regs)	
callee-save reg 1	
(temp var. n)	
temp var. 1	←──── %sp + 4 ←─── %sp



Calling Convention

- When calling a function:
 - 1. Push caller-save regs onto stack
 - 2. Push function args onto stack
 - 3. Push return address then jump
- In function after call:
 - 1. Push old %bp + set %bp = %sp
 - 2. Allocate space for local variables
 - 3. Push callee-save regs if necessary

(function arg. n)	
function arg. 1	s + 8≪
return address	•
old %ebp	• %bp + 4
local var. 1	←─── %bp ←─── %bp - 4
(local var. n)	- 900p - 4
(callee-save regs)	
callee-save reg 1	
(temp var. n)	
temp var. 1	←──── %sp + 4 ←─── %sp



Instructions: Stack Access

• pushl source

- %sp ← %sp 4
- \circ %ss:(%sp) ← source
- popl dest
 - dest ← %ss:(%sp)
 - \circ %sp \leftarrow %sp + 4

(function arg. n …)	
function arg. 1	0%bp ± 8
return address	%bp + 8
old %ebp	%bp + 4
local var. 1	← %bp
(local var. n)	• %bp - 4
(callee-save regs)	
callee-save reg 1	
(temp var. n)	• • •
temp var. 1	← %sp + 4 ← %sp



Instructions: Control Flow

- jmp label
 - \circ %eip \leftarrow label
- ljmp NEW_CS, offset
 - \circ %cs \leftarrow NEW_CS
 - \circ %eip \leftarrow offset

- call label
 - push %eip
 - o %eip ← label
- ret
 - pop %eip



Instructions: Conditional Jump

- Relies on %eflags bits
 - Most arithmetic operations change %eflags
- j* label
 - Jump to label if * flag is 1, can be z(ero), e(qual), etc
- jn* label
 - \circ Jump to label if * flag is 0



Assembler Directives

- Commands that "direct" the assembler
 - \circ Are not instructions
- Examples:
 - .globl defines a list of symbols as global
 - .equ defines a constant (like #define)
 - .bytes, .word, .asciz reserve space in RO memory

https://docs.oracle.com/cd/E26502_01/html/E28388/eoiyg.html



Assembler Segments

- Organize memory by data properties
 - .text holds executable instructions
 - .bss holds zero-initialized data (e.g. static int i;)
 - .data holds initialized data (e.g. char c = 'a';)
 - .rodata holds read-only data
- Stack Initialized by linker & loader
- Heap Start defined by compiler, contents by programmer (usually thru an included library)

BIOS = Basic Input/Output System

- Firmware (compiled with the device itself)
- Configures buses/connections to and from various hardware devices
- Setups RAM for usage, puts CPU in 20-bit Real Mode
- Determines which device (CD, USB, etc) has a loadable boot sequence.
 - Loads boot sequence into RAM, then transfers control to top of boot sequence
- Afterwards, BIOS provides abstractions over low-level interfaces, such as reading/writing from external disks

- Source:
 - <u>https://wiki.osdev.org/System_Initialization_(x86)</u>
 - https://wiki.osdev.org/BIOS



BIOS Services

- Use BIOS services through `int` instruction
 - Store parameters in specified registers, such as AH/AX
 - Software Interrupt: triggers a func in firmware
- asm: `int SERVICE_NUM`
 - i.e. int \$0x10: Video services, int \$0x13: Disk services



Useful GDB Commands

- r show register values
- sreg show segment registers
- s step into instruction
- n next instruction
- c continue
- u <start> <stop> disassembles
 C code into assembly

- b set a breakpoint
- d <n> delete a breakpoint
- bpd / bpe <n> disable / enable a breakpoint
- x/Nx addr display hex dump of N words, starting at addr
- x/Ni addr display N instructions, starting at addr

Assembly can be useful

In order to understand **efficiency** of implementation, sometimes looking at the assembly of generated code can be useful:

https://godbolt.org/z/5rxdz8s8K