

# Polymorphism

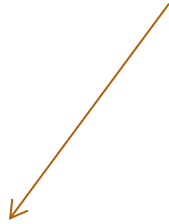
COS 326

Speaker: Andrew Appel

Princeton University



polymorphic,  
higher-order  
programming



**POLY-HO!**



## Here's an annoying thing

```
let rec map (f:int->int) (xs:int list) : int list =  
  match xs with  
  | [] -> []  
  | hd::tl -> (f hd)::(map f tl);;
```

What if I want to increment a list of floats?

Alas, I can't just call this map. It works on ints!



# Here's an annoying thing

```
let rec map (f:int->int) (xs:int list) : int list =  
  match xs with  
  | [] -> []  
  | hd::tl -> (f hd)::(map f tl);;
```

What if I want to increment a list of floats?

Alas, I can't just call this map. It works on ints!

```
let rec mapfloat (f:float->float) (xs:float list) :  
  float list =  
  match xs with  
  | [] -> []  
  | hd::tl -> (f hd)::(mapfloat f tl);;
```



# Turns out

```
let rec map f xs =  
  match xs with  
  | [] -> []  
  | hd::tl -> (f hd)::(map f tl)
```

```
let ints = map (fun x -> x + 1) [1; 2; 3; 4]
```

```
let floats = map (fun x -> x +. 2.0) [3.1415; 2.718]
```

```
let strings = map String.uppercase ["sarah"; "joe"]
```



# Type of the undecorated map?

```
let rec map f xs =  
  match xs with  
  | [] -> []  
  | hd::tl -> (f hd)::(map f tl)  
  
map : ('a -> 'b) -> 'a list -> 'b list
```



# Type of the undecorated map?

```
let rec map f xs =  
  match xs with  
  | [] -> []  
  | hd::tl -> (f hd)::(map f tl)  
  
map : ('a -> 'b) -> 'a list -> 'b list
```

We often use greek letters like  $\alpha$  or  $\beta$  to represent type variables.

Read as:

- for any types 'a and 'b,
- if you give map a function from 'a to 'b,
- it will return a function
  - which when given a list of 'a values
  - returns a list of 'b values.



# We can say this explicitly

```
let rec map (f:'a -> 'b) (xs:'a list) : 'b list =  
  match xs with  
  | [] -> []  
  | hd::tl -> (f hd)::(map f tl)  
  
map : ('a -> 'b) -> 'a list -> 'b list
```

The OCaml compiler is smart enough to figure out that this is the *most general* type that you can assign to the code.  
(technical term: *principal type*)

We say map is *polymorphic* in the types 'a and 'b – just a fancy way to say map can be used on any types 'a and 'b.

Java generics derived from ML-style polymorphism (but added after the fact and more complicated due to subtyping)





# More realistic polymorphic functions

```
let rec merge (lt:'a->'a->bool) (xs:'a list) (ys:'a list) : 'a list =  
  match (xs,ys) with  
  | ([],_) -> ys  
  | (_,[]) -> xs  
  | (x::xst, y::yst) ->  
    if lt x y then x::(merge lt xst ys)  
    else y::(merge lt xs yst)
```

```
let rec split (xs:'a list) (ys:'a list) (zs:'a list) : 'a list * 'a list =  
  match xs with  
  | [] -> (ys, zs)  
  | x::rest -> split rest zs (x::ys)
```

```
let rec mergesort (lt:'a->'a->bool) (xs:'a list) : 'a list =  
  match xs with  
  | ([] | _::[]) -> xs  
  | _ -> let (first,second) = split xs [] [] in  
    merge lt (mergesort lt first) (mergesort lt second)
```



# More realistic polymorphic functions

```
mergesort : ('a->'a->bool) -> 'a list -> 'a list
```

```
mergesort (<) [3;2;7;1]  
  == [1;2;3;7]
```

```
mergesort (>) [2; 3; 42]  
  == [42 ; 3; 2]
```

```
mergesort (fun x y -> String.compare x y < 0) ["Hi"; "Bi"]  
  == ["Bi"; "Hi"]
```

```
let int_sort = mergesort (<)
```

```
let int_sort_down = mergesort (>)
```

```
let str_sort = mergesort (fun x y -> String.compare x y < 0)
```



# Another Interesting Function

```
let comp f g x = f (g x)
```

```
let mystery = comp (add 1) square
```



```
let comp = fun f -> (fun g -> (fun x -> f (g x)))
```

```
let mystery = comp (add 1) square
```



```
let mystery =  
  (fun f -> (fun g -> (fun x -> f (g x)))) (add 1) square
```

A diagram with orange arrows and brackets. One arrow points from the 'f' parameter of the lambda function to the 'add 1' argument. Another arrow points from the 'g' parameter to the 'square' argument. A third arrow points from the 'x' parameter to the 'f (g x)' body. Brackets are placed under 'add 1' and 'square' to indicate they are arguments.

```
let mystery = fun x -> (add 1) (square x)
```



```
let mystery x = add 1 (square x)
```



# Function composition!

```
let comp f g x = f (g x)
```

```
let mystery = comp (add 1) square
```

$$(f \circ g)(x) = f(g(x))$$

$$\text{mystery} = (\text{add } 1) \circ \text{square}$$

$$\text{mystery}(x) = (\text{add } 1) (\text{square } (x))$$



# What is the type of comp?

```
let comp f g x = f (g x)
```

```
let comp (f: 'b->'c) (g: 'a->'b) (x: 'a) : 'c  
= f (g x)
```

```
comp : ('b -> 'c) ->  
        ('a -> 'b) ->  
        ('a -> 'c)
```

# Optimization

What does this program do?

```
map f (map g [x1; x2; ...; xn])
```

For each element of the list  $x_1, x_2, x_3 \dots x_n$ , it executes  $g$ , creating:

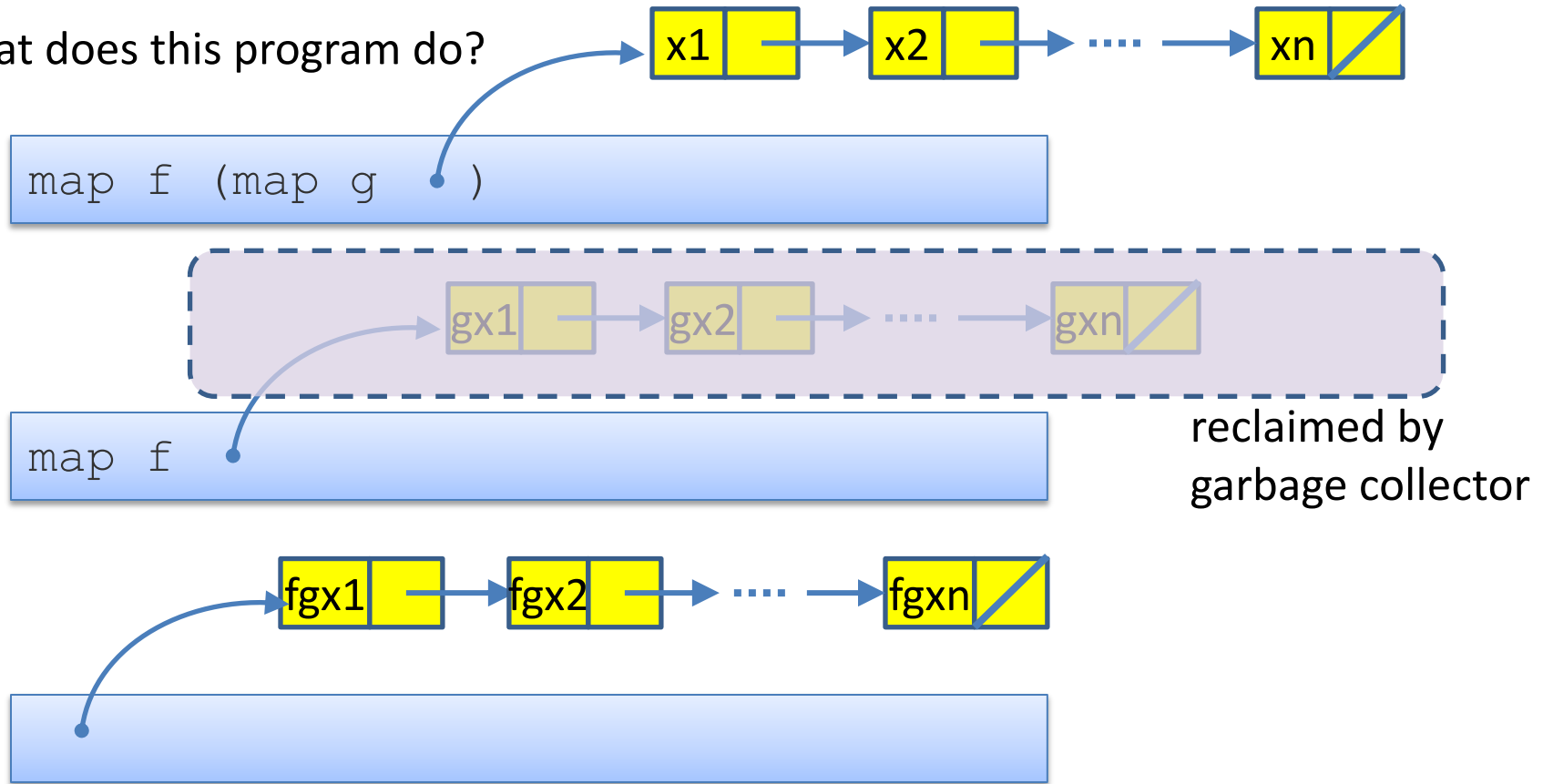
```
map f ([g x1; g x2; ...; g xn])
```

Then for each element of the list  $[g x_1, g x_2, g x_3 \dots g x_n]$ , it executes  $f$ , creating:

```
[f (g x1); f (g x2); ...; f (g xn)]
```

# Optimization

What does this program do?



# Optimization

What does this program do?

```
map f (map g [x1; x2; ...; xn])
```

For each element of the list  $x_1, x_2, x_3 \dots x_n$ , it executes  $g$ , creating:

```
map f ([g x1; g x2; ...; g xn])
```

Then for each element of the list  $[g x_1, g x_2, g x_3 \dots g x_n]$ , it executes  $f$ , creating:

```
[f (g x1); f (g x2); ...; f (g xn)]
```

Is there a faster way? Yes! (And query optimizers for SQL do it for you.)

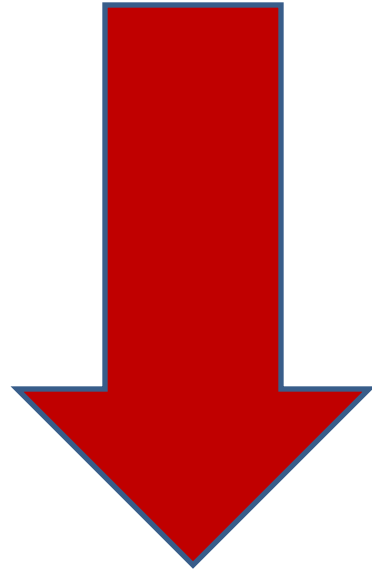
```
map (comp f g) [x1; x2; ...; xn]
```





# Deforestation

```
map f (map g [x1; x2; ...; xn])
```



This kind of optimization has a name:

**deforestation**

(because it eliminates intermediate lists and, um, trees...)

```
map (comp f g) [x1; x2; ...; xn]
```

# How about reduce?

```
let rec reduce f u xs =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)
```

What's the most general type of reduce?

# How about reduce?

```
let rec reduce f u xs =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)
```

What's the most general type for `reduce`?

Based on the patterns, we know `xs` must be a ('a list) for some type 'a.

# How about reduce?

```
let rec reduce f u (xs: 'a list) =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)
```

What's the most general type of reduce?

# How about reduce?

```
let rec reduce f u (xs: 'a list) =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)
```

What's the most general type of reduce?

f is called so it  
must be a  
function of two  
arguments.

# How about reduce?

```
let rec reduce (f:? -> ? -> ?) u (xs: 'a list) =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)
```

What's the most general type of reduce?

# How about reduce?

```
let rec reduce (f:? -> ? -> ?) u (xs: 'a list) =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)
```

What's the most general type of reduce?

Furthermore, `hd` came from `xs`, so `f` must take an `'a` value as its first argument.

# How about reduce?

```
let rec reduce (f:'a -> ? -> ?) u (xs: 'a list) =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)
```

What's the most general type of reduce?



# How about reduce?

```
let rec reduce (f:'a -> ? -> ?) u (xs: 'a list) =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)
```

What's the most general type of reduce?

The second argument to f must have the same type as the result of reduce. Let's call it 'b.'

# How about reduce?

```
let rec reduce (f:'a -> 'b -> ?) u (xs: 'a list) : 'b =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)
```

What's the most general type of reduce?

The result of f  
must have the  
same type as the  
result of reduce  
overall: 'b.

# How about reduce?

```
let rec reduce (f:'a -> 'b -> 'b) u (xs: 'a list) : 'b =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)
```

What's the most general type of reduce?



# How about reduce?

```
let rec reduce (f:'a -> 'b -> ?) u (xs: 'a list) : 'b =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)
```

What's the most general type of reduce?

If xs is empty,  
then reduce  
returns u. So u's  
type must be 'b.

# How about reduce?

```
let rec reduce (f:'a -> 'b -> ?) (u:'b) (xs: 'a list) : 'b =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)
```

What's the most general type of reduce?

# How about reduce?

```
let rec reduce (f:'a -> 'b -> ?) (u:'b) (xs: 'a list) : 'b =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)
```

What's the most general type of reduce?

reduce returns  
the result of f. So  
f's result type  
must be 'b.

# How about reduce?

```
let rec reduce (f:'a -> 'b -> 'b) (u:'b) (xs: 'a list) : 'b =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)
```

What's the most general type of reduce?

# How about reduce?

```
let rec reduce (f:'a -> 'b -> 'b) (u:'b) (xs: 'a list) : 'b =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)
```

What's the most general type of reduce?

```
('a -> 'b -> 'b) -> 'b -> 'a list -> 'b
```





# What does this do?


```
let rec reduce f u xs =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)  
  
let mystery0 = reduce (fun x y -> 1+y) 0
```

# What does this do?

```
let rec reduce f u xs =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl);;  
  
let mystery0 = reduce (fun x y -> 1+y) 0;;  
  
let rec mystery0 xs =  
  match xs with  
  | [] -> 0  
  | hd::tl ->  
    (fun x y -> 1+y) hd (reduce (fun ...) 0 tl)
```

# What does this do?

```
let rec reduce f u xs =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl);;  
  
let mystery0 = reduce (fun x y -> 1+y) 0;;  
  
let rec mystery0 xs =  
  match xs with  
  | [] -> 0  
  | hd::tl ->  
    (fun x y -> 1+y) hd (reduce (fun ... ) 0 tl)
```



# What does this do?

```
let rec reduce f u xs =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl);;  
  
let mystery0 = reduce (fun x y -> 1+y) 0;;  
  
let rec mystery0 xs =  
  match xs with  
  | [] -> 0  
  | hd::tl ->  
    (fun y -> 1+y) (reduce (fun ...) 0 tl)
```

# What does this do?

```
let rec reduce f u xs =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)  
  
let mystery0 = reduce (fun x y -> 1+y) 0  
  
let rec mystery0 xs =  
  match xs with  
  | [] -> 0  
  | hd::tl -> 1 + reduce (fun ...) 0 tl
```

# What does this do?

```
let rec reduce f u xs =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)  
  
let mystery0 = reduce (fun x y -> 1+y) 0  
  
let rec mystery0 xs =  
  match xs with  
  | [] -> 0  
  | hd::tl -> 1 + mystery0 tl
```

# What does this do?

```
let rec reduce f u xs =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)  
  
let mystery0 = reduce (fun x y -> 1+y) 0  
  
let rec mystery0 xs =  
  match xs with  
  | [] -> 0  
  | hd::tl -> 1 + mystery0 tl  List Length!
```

# What does this do?

```
let rec reduce f u xs =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl);;  
  
let mystery1 = reduce (fun x y -> x::y) []
```



# What does this do?

```
let rec reduce f u xs =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)  
  
let mystery1 = reduce (fun x y -> x::y) []  
  
let rec mystery1 xs =  
  match xs with  
  | [] -> []  
  | hd::tl -> hd::(mystery1 tl)  Copy!
```

## And this one?

```
let rec reduce f u xs =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)  
  
let mystery2 g =  
  reduce (fun a b -> (g a)::b) []
```

# And this one?

```
let rec reduce f u xs =  
  match xs with  
  | [] -> u  
  | hd::tl -> f hd (reduce f u tl)
```

```
let mystery2 g =  
  reduce (fun a b -> (g a)::b) []
```

```
let rec mystery2 g xs =  
  match xs with  
  | [] -> []  
  | hd::tl -> (g hd)::(mystery2 g tl) map!
```

# Map and Reduce

```
val map : ('a -> 'b) -> 'a list -> 'b list
```

```
val reduce : ('a -> 'b -> 'b) -> 'b -> 'a list -> 'b
```

We coded **map** in terms of **reduce**:

- ie: we showed we can compute **map f xs** using a call to **reduce** ??? just by passing the right arguments in place of ???

Can we code **reduce** in terms of **map**?



# Map and Reduce

```
val map : ('a -> 'b) -> 'a list -> 'b list
```

```
val reduce : ('a -> 'b -> 'b) -> 'b -> 'a list -> 'b
```

let reduce f u xs = ... map (...) (...) ...

*(use only: map, f, u, xs; don't use rec )*

reduce (+) 0 [1;2;3] = ... map (...) (...) ...



# Some Other Combinators: List Module

<https://caml.inria.fr/pub/docs/manual-ocaml/libref/List.html>

```
val fold_left : ('a -> 'b -> 'a) -> 'a -> 'b list -> 'a  
val fold_right : ('a -> 'b -> 'b) -> 'a list -> 'b -> 'b
```

```
val mapi : (int -> 'a -> 'b) -> 'a list -> 'b list  
List.mapi f [a0; ...; an] == [f 0 a0; ... ; f n an]
```

```
val map2 : ('a -> 'b -> 'c) -> 'a list -> 'b list -> 'c list  
List.map2 f [a0; ...; an] [b0; ...; bn] == [f a0 b0 ; ... ; f an bn]
```

```
val iter : ('a -> unit) -> 'a list -> unit  
List.iter f [a0; ...; an] == f a0; ... ; f an
```



# Summary

- Map and reduce are two *higher-order functions* that capture very, very common *recursion patterns*
- Reduce is especially powerful:
  - related to the “visitor pattern” of OO languages like Java.
  - can implement most list-processing functions using it, including things like copy, append, filter, reverse, map, etc.
- We can write clear, terse, reusable code by exploiting:
  - higher-order functions
  - anonymous functions
  - first-class functions
  - polymorphism



# Practice Problems

Using map, write a function that takes a list of pairs of integers, and produces a list of the sums of the pairs.

- e.g., `list_add [(1,3); (4,2); (3,0)] = [4; 6; 3]`
- Write `list_add` directly using `reduce`.

Using map, write a function that takes a list of pairs of integers, and produces their quotient if it exists.

- e.g., `list_div [(1,3); (4,2); (3,0)] = [Some 0; Some 2; None]`
- Write `list_div` directly using `reduce`.

Using reduce, write a function that takes a list of optional integers, and filters out all of the `None`'s.

- e.g., `filter_none [Some 0; Some 2; None; Some 1] = [0;2;1]`
- Why can't we directly use `filter`? How would you generalize `filter` so that you can compute `filter_none`? Alternatively, rig up a solution using `filter` + `map`.

Using reduce, write a function to compute the sum of squares of a list of numbers.

- e.g., `sum_squares = [3,5,2] = 38`

