

OCaml Datatypes

COS 326

David Walker

Princeton University

OCaml So Far

- We have seen a number of basic types:
 - int
 - float
 - char
 - string
 - bool
- We have seen a few structured types:
 - pairs
 - tuples
 - options
 - lists
- In this lecture, we will see some more general ways to define our own new types and data structures

Type Abbreviations

- We have already seen some type abbreviations:

```
type point = float * float
```

Type Abbreviations

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```
type point = float * float
```

- These abbreviations can be helpful documentation:

```
let distance (p1:point) (p2:point) : float =  
  let square x = x *. x in  
  let (x1,y1) = p1 in  
  let (x2,y2) = p2 in  
  sqrt (square (x2 -. x1) +. square (y2 -. y1))
```

- But they add nothing of *substance* to the language
 - they are *equal* in every way to an existing type

Type Abbreviations

- We have already seen some type abbreviations:

```
type point = float * float
```

- As far as OCaml is concerned, you could have written:

```
let distance (p1:float*float)
             (p2:float*float) : float =
  let square x = x *. x in
  let (x1,y1) = p1 in
  let (x2,y2) = p2 in
  sqrt (square (x2 -. x1) +. square (y2 -. y1))
```

- Since the types are equal, you can *substitute* the definition for the name wherever you want
 - we have not added any new data structures

DATA TYPES

Data types

- OCaml provides a general mechanism called a **data type** for defining new data structures that consist of many alternatives

```
type my_bool = Tru | Fal
```

a **value** with type **my_bool**
is one of two things:

- **Tru**, or
- **Fal**

read the "|" as "or"

Data types

- OCaml provides a general mechanism called a **data type** for defining new data structures that consist of many alternatives

```
type my_bool = Tru | Fal
```

Tru and Fal are called
"constructors"

a **value** with type **my_bool**
is one of two things:

- **Tru**, or
- **Fal**

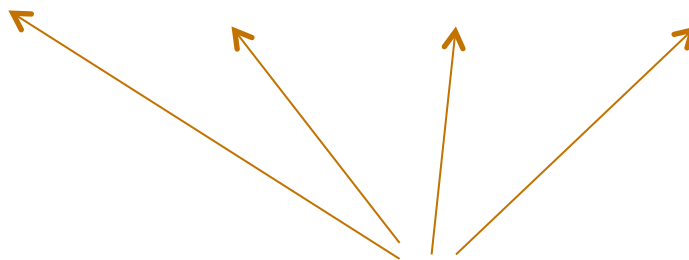
read the "|" as "or"

Data types

- OCaml provides a general mechanism called a **data type** for defining new data structures that consist of many alternatives

```
type my_bool = Tru | Fal
```

```
type color = Blue | Yellow | Green | Red
```



there's no need to stop
at 2 cases; define as many
alternatives as you want

Data types

- OCaml provides a general mechanism called a **data type** for defining new data structures that consist of many alternatives

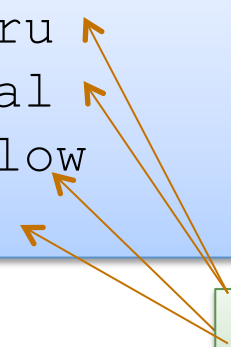
```
type my_bool = Tru | Fal
```

```
type color = Blue | Yellow | Green | Red
```

- Creating values:

```
let b1 : my_bool = Tru
let b2 : my_bool = Fal
let c1 : color = Yellow
let c2 : color = Red
```

use constructors to create values



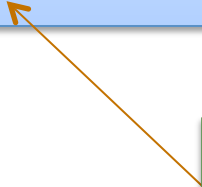
Data types

```
type color = Blue | Yellow | Green | Red

let c1 : color = Yellow
let c2 : color = Red
```

- Using data type values:

```
let print_color (c:color) : unit =
  match c with
  | Blue ->
  | Yellow ->
  | Green ->
  | Red ->
```



use pattern matching to determine which color you have; act accordingly

Data types

```
type color = Blue | Yellow | Green | Red
```

```
let c1 : color = Yellow
```

```
let c2 : color = Red
```

- Using data type values:

```
let print_color (c:color) : unit =  
  match c with  
  | Blue -> print_string "blue"  
  | Yellow -> print_string "yellow"  
  | Green -> print_string "green"  
  | Red -> print_string "red"
```

Data types

```
type color = Blue | Yellow | Green | Red

let c1 : color = Yellow
let c2 : color = Red
```

- Using data type values:

```
let print_color (c:color) : unit =
  match c with
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  | Red -> print_string "red"
```

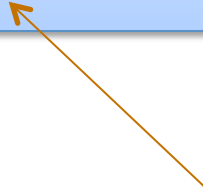
Why not just use strings to represent colors instead of defining a new type?

Data types

```
type color = Blue | Yellow | Green | Red
```

oops!:

```
let print_color (c:color) : unit =  
  match c with  
  | Blue -> print_string "blue"  
  | Yellow -> print_string "yellow"  
  | Red -> print_string "red"
```



Warning 8: this pattern-matching is not exhaustive.
Here is an example of a value that is not matched:
Green

Data types

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type color = Blue | Yellow | Green | Red
```

oops!:

```
let print_color (c:color) : unit =  
  match c with  
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```

Warning 8: this pattern-matching is not exhaustive.
Here is an example of a value that is not matched:
Green

OCaml's datatype mechanism allow you to create types
that contain *precisely* the values you want!

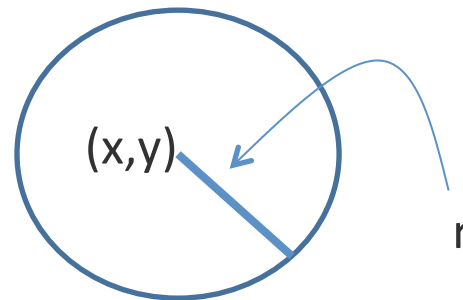
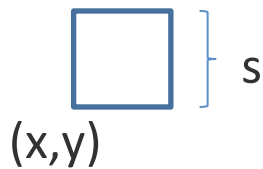
Data Types Can Carry Additional Values

- Data types are more than just enumerations of constants:

```
type point = float * float

type simple_shape =
  Circle of point * float
| Square of point * float
```

- Read as: a **simple_shape** is either:
 - a **Circle**, which contains a **pair** of a **point** and **float**, or
 - a **Square**, which contains a **pair** of a **point** and **float**



Data Types Can Carry Additional Values

- Data types are more than just enumerations of constants:

```
type point = float * float

type simple_shape =
  Circle of point * float
| Square of point * float

let origin : point = (0.0, 0.0)

let circ1   : simple_shape = Circle (origin, 1.0)
let circ2   : simple_shape = Circle ((1.0, 1.0), 5.0)
let square  : simple_shape = Square (origin, 2.3)
```

Data Types Can Carry Additional Values

- Data types are more than just enumerations of constants:

```
type point = float * float

type simple_shape =
  | Circle of point * float
  | Square of point * float

let simple_area (s:simple_shape) : float =
  match s with
  | Circle (_, radius) -> 3.14 *. radius *. radius
  | Square (_, side) -> side *. side
```

Compare

- Data types are more than just enumerations of constants:

```
type point = float * float

type simple_shape =
  | Circle of point * float
  | Square of point * float

let simple_area (s:simple_shape) : float =
  match s with
  | Circle (_, radius) -> 3.14 *. radius *. radius
  | Square (_, side) -> side *. side
```

```
type my_shape = point * float

let simple_area (s:my_shape) : float =
  (3.14 *. radius *. radius) ?? or ?? (side *. side)
```

More General Shapes

```

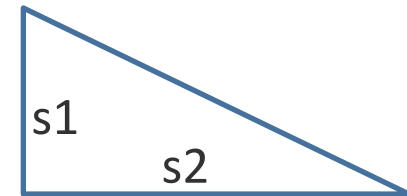
type point = float * float

type shape =
  Square of float
| Ellipse of float * float
| RtTriangle of float * float
| Polygon of point list
  
```

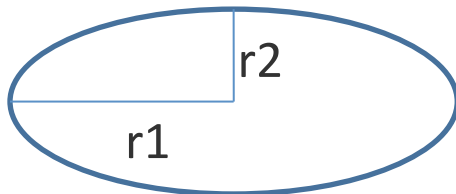
Square $s =$



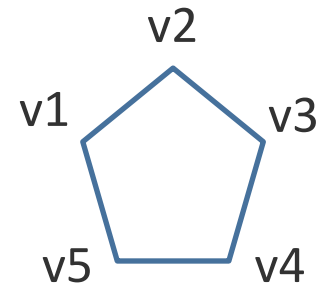
RtTriangle $(s1, s2) =$



Ellipse $(r1, r2) =$



Polygon $[v1; \dots; v5] =$



More General Shapes

```

type point = float * float
type radius = float
type side = float

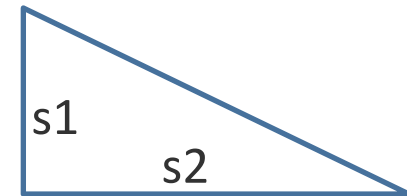
type shape =
  | Square of side
  | Ellipse of radius * radius
  | RtTriangle of side * side
  | Polygon of point list
  
```

Type abbreviations can aid readability

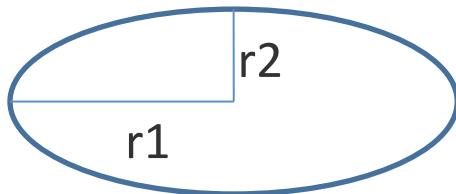
Square $s =$



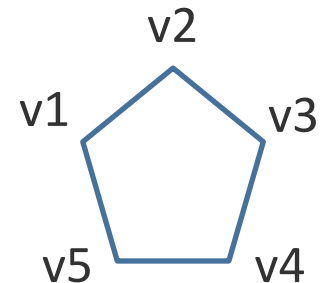
RtTriangle $(s1, s2) =$



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RtTriangle $[v1; \dots; v5] =$



More General Shapes

```
type point = float * float
type radius = float
type side = float

type shape =
  | Square of side
  | Ellipse of radius * radius
  | RtTriangle of side * side
  | Polygon of point list
```

Square builds a shape
from a single side

RtTriangle builds a shape
from a pair of sides

```
let sq    : shape = Square 17.0
let ell   : shape = Ellipse (1.0, 2.0)
let rt    : shape = RtTriangle (1.0, 1.0)
let poly  : shape = Polygon [(0., 0.); (1., 0.); (0.; 1.)]
```

they are all shapes;
they are constructed in
different ways

Polygon builds a shape
from a list of points
(where each point is itself a pair)

More General Shapes

```
type point = float * float
type radius = float
type side = float

type shape =
  Square of side
| Ellipse of radius * radius
| RtTriangle of side * side
| Polygon of point list
```

```
let area (s : shape) : float =
  match s with
  | Square s ->
  | Ellipse (r1, r2)->
  | RtTriangle (s1, s2) ->
  | Polygon ps ->
```

a data type also defines
a pattern for matching

More General Shapes

```
type point = float * float
type radius = float
type side = float

type shape =
  Square of side
  | Ellipse of radius * radius
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a data type also defines
a pattern for matching

```
let area (s : shape) : float =
  match s with
  | Square s ->
  | Ellipse (r1, r2) ->
  | RtTriangle (s1, s2) ->
  | Polygon ps ->
```

Square carries a value
with type **float** so **s** is
a pattern for float values

RtTriangle carries a value
with type **float * float**
so **(s1, s2)** is a pattern
for that type

More General Shapes

```
type point = float * float
type radius = float
type side = float

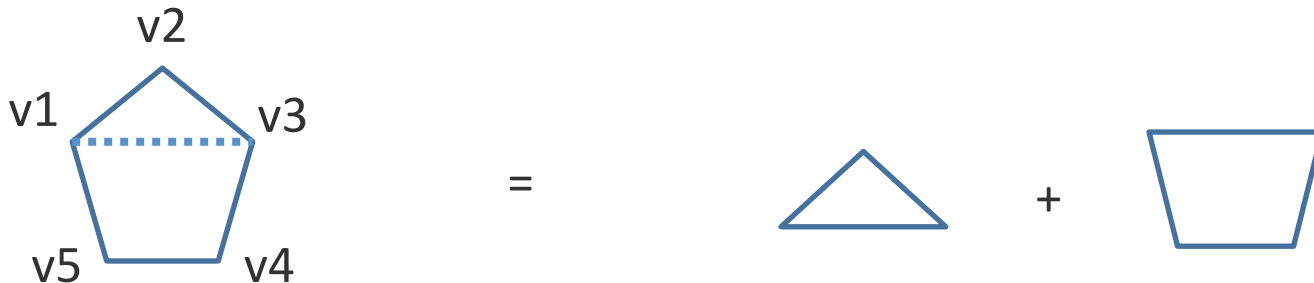
type shape =
  Square of side
  | Ellipse of radius * radius
  | RtTriangle of side * side
  | Polygon of point list
```

a data type also defines
a pattern for matching

```
let area (s : shape) : float =
  match s with
  | Square s -> s *. s
  | Ellipse (r1, r2) -> pi *. r1 *. r2
  | RtTriangle (s1, s2) -> s1 *. s2 /. 2.
  | Polygon ps -> ???
```

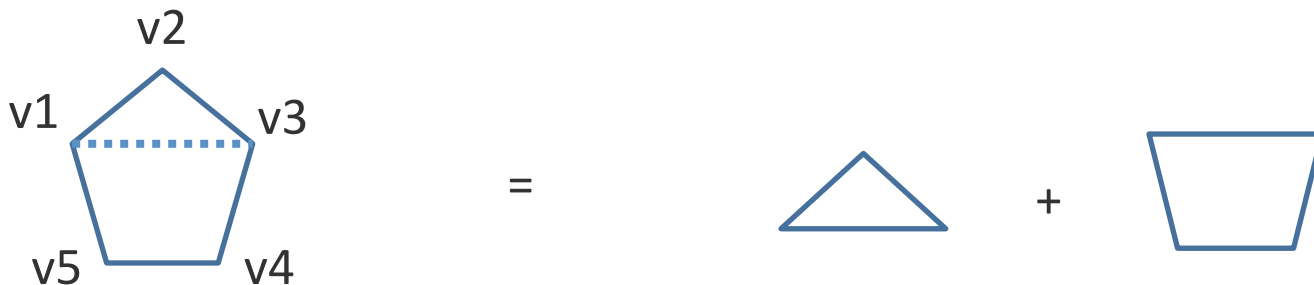
Computing Area

- How do we compute polygon area?
- For convex polygons:
 - **Case: the polygon has fewer than 3 points:**
 - it has 0 area! (it is a line or a point or nothing at all)
 - **Case: the polygon has 3 or more points:**
 - Compute the area of the triangle formed by the first 3 vertices
 - Delete the second vertex to form a new polygon
 - Sum the area of the triangle and the new polygon



Computing Area

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 - Compute the area of the triangle formed by the first 3 vertices
 - Delete the second vertex to form a new polygon
 - Sum the area of the triangle and the new polygon
- Note: This is a beautiful **inductive algorithm**:
 - the area of a polygon with n points is computed in terms of a smaller polygon with only $n-1$ points!

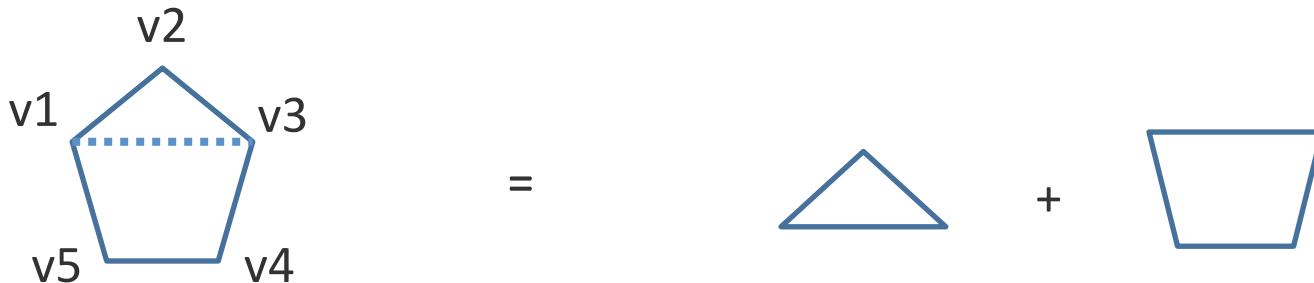


Computing Area

```
let area (s : shape) : float =
  match s with
  | Square s -> s *. s
  | Ellipse (r1, r2) -> r1 *. r2
  | RtTriangle (s1, s2) -> s1 *. s2 /. 2.
  | Polygon ps -> poly_area ps
```

This pattern says the list has at least 3 items

```
let poly_area (ps : point list) : float =
  match ps with
  | p1 :: p2 :: p3 :: tail ->
    tri_area p1 p2 p3 +. poly_area (p1::p3::tail)
  | _ -> 0.
```



Computing Area

```
let tri_area (p1:point) (p2:point) (p3:point) : float =
  let a = distance p1 p2 in
  let b = distance p2 p3 in
  let c = distance p3 p1 in
  let s = 0.5 *. (a +. b +. c) in
  sqrt (s *. (s -. a) *. (s -. b) *. (s -. c))
```

```
let rec poly_area (ps : point list) : float =
  match ps with
  | p1 :: p2 :: p3 :: tail ->
    tri_area p1 p2 p3 +. poly_area (p1::p3::tail)
  | _ -> 0.
```

```
let area (s : shape) : float =
  match s with
  | Square s -> s *. s
  | Ellipse (r1, r2)-> pi *. r1 *. r2
  | RtTriangle (s1, s2) -> s1 *. s2 /. 2.
  | Polygon ps -> poly_area ps
```

INDUCTIVE DATA TYPES

Inductive data types

- We can use data types to define inductive data
- A binary tree is:
 - a **Leaf** containing no data
 - a **Node** containing a **key**, a **value**, a left subtree and a right subtree

Inductive data types

- We can use data types to define inductive data
- A binary tree is:
 - a **Leaf** containing no data
 - a **Node** containing a **key**, a **value**, a left subtree and a right subtree

```
type key = string
type value = int

type tree =
  Leaf
| Node of key * value * tree * tree
```


Inductive data types

```
type key = int
type value = string

type tree =
  | Leaf
  | Node of key * value * tree * tree
```

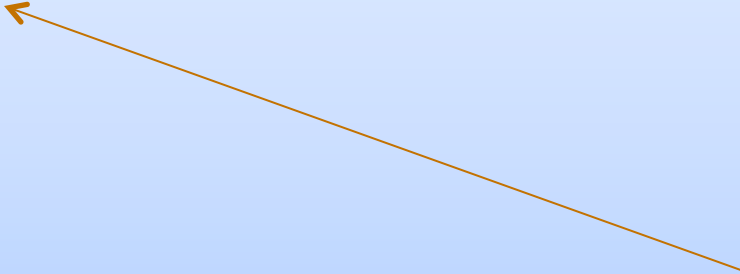
```
let rec insert (t:tree) (k:key) (v:value) : tree =
```

Inductive data types

```
type key = int
type value = string

type tree =
  Leaf
| Node of key * value * tree * tree
```

```
let rec insert (t:tree) (k:key) (v:value) : tree =
  match t with
  | Leaf ->
  | Node (k', v', left, right) ->
```



Again, the type definition specifies the cases you must consider

Inductive data types

```
type key = int
type value = string

type tree =
  Leaf
| Node of key * value * tree * tree
```

```
let rec insert (t:tree) (k:key) (v:value) : tree =
  match t with
  | Leaf -> Node (k, v, Leaf, Leaf)
  | Node (k', v', left, right) ->
```

Inductive data types

```
type key = int
type value = string

type tree =
  Leaf
| Node of key * value * tree * tree
```

```
let rec insert (t:tree) (k:key) (v:value) : tree =
  match t with
  | Leaf -> Node (k, v, Leaf, Leaf)
  | Node (k', v', left, right) ->
    if k < k' then
      Node (k', v', insert left k v, right)
    else if k > k' then
      Node (k', v', left, insert right k v)
    else
      Node (k, v, left, right)
```

Inductive data types

```
type key = int
type value = string

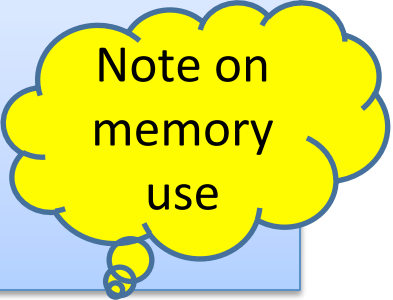
type tree =
  Leaf
| Node of key * value * tree * tree
```

```
let rec insert (t:tree) (k:key) (v:value) : tree =
  match t with
  | Leaf -> Node (k, v, Leaf, Leaf)
  | Node (k', v', left, right) ->
    if k < k' then
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Inductive data types

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type key = int
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```



Note on
memory
use

```
let rec insert (t:tree) (k:key) (v:value) : tree =
  match t with
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      Node (k', v', insert left k v, right)
    else if k > k' then
      Node (k', v', left, insert right k v)
    else
      Node (k, v, left, right)
```

Inductive data types: Another Example

- Recall, we used the type "int" to represent natural numbers
 - but that was kind of broken: it also contained negative numbers
 - we had to use a dynamic test to guard entry to a function:

```
let double (n : int) : int =  
  if n < 0 then  
    raise (Failure "negative input!")  
  else  
    double_nat n
```

- it would be nice if there was a way to define the natural numbers **exactly**, and use OCaml's type system to guarantee no client ever attempts to double a negative number

Inductive data types

- Recall, a natural number n is either:
 - zero, or
 - $m + 1$
- We use a data type to represent this definition exactly:

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 - zero, or
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```
type nat = Zero | Succ of nat
```

Inductive data types

- Recall, a natural number n is either:
 - zero, or
 - $m + 1$
- We use a data type to represent this definition exactly:

```
type nat = Zero | Succ of nat

let rec nat_to_int (n : nat) : int =
  match n with
  | Zero -> 0
  | Succ n -> 1 + nat_to_int n
```

Inductive data types

- Recall, a natural number n is either:
 - zero, or
 - $m + 1$
- We use a data type to represent this definition exactly:

```
type nat = Zero | Succ of nat

let rec nat_to_int (n : nat) : int =
  match n with
  | Zero -> 0
  | Succ n -> 1 + nat_to_int n

let rec double_nat (n : nat) : nat =
  match n with
  | Zero -> Zero
  | Succ m -> Succ (Succ(double_nat m))
```

Summary

- OCaml data types: a powerful mechanism for defining complex data structures:
 - They are precise
 - contain exactly the elements you want, not more elements
 - They are general
 - recursive, non-recursive (mutually recursive and polymorphic)
 - The type checker helps you detect errors
 - missing cases in your functions