COS 318: Operating Systems

Storage Devices

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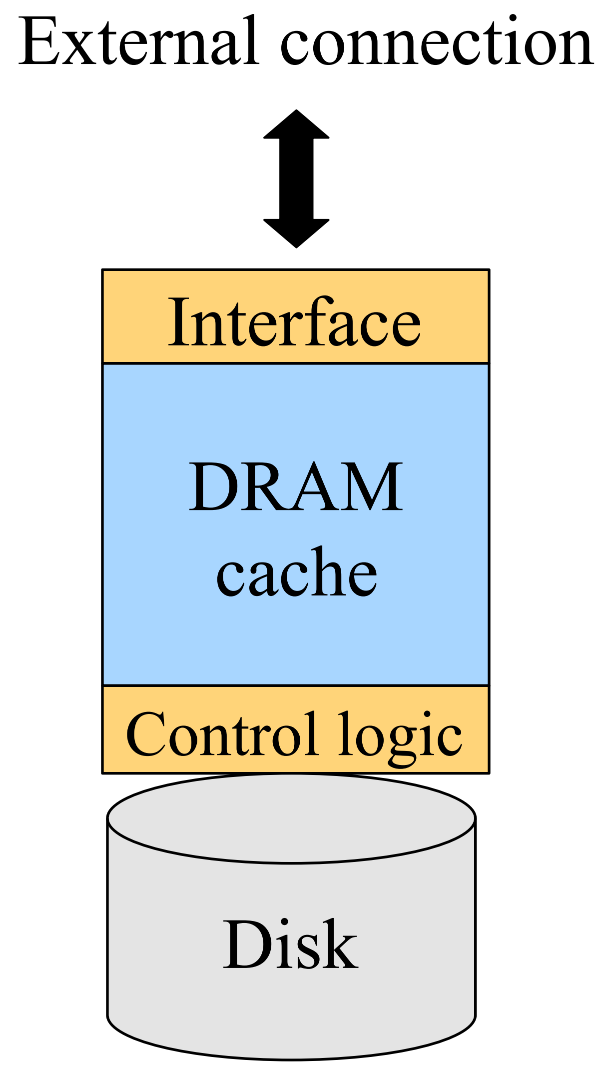
http://www.cs.princeton.edu/courses/archive/fall13/cos318/

Today’s Topics

- Magnetic disks
- Magnetic disk performance
- Disk arrays
- Flash memory

A Typical Magnetic Disk Controller

- External interfaces
  - IDE/ATA, **SATA(1.0, 2.0, 3.0)**
  - SCSI (1, 2, 3),
    Ultra-(160, 320, 640) SCSI
  - Fibre channel
- Cache
  - Buffer data between disk and interface
- Control logic
  - Read/write operations
  - Cache replacement
  - Failure detection and recovery

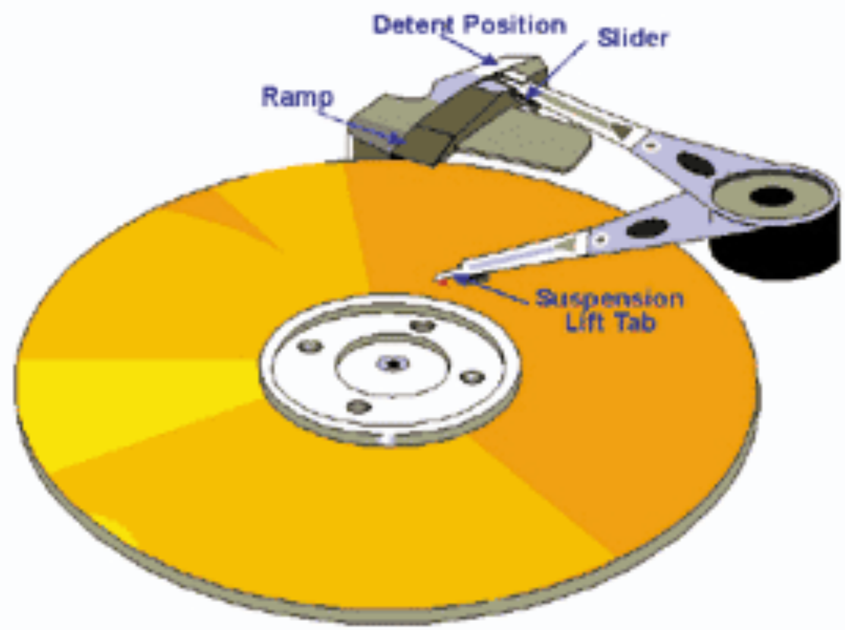


Disk Caching

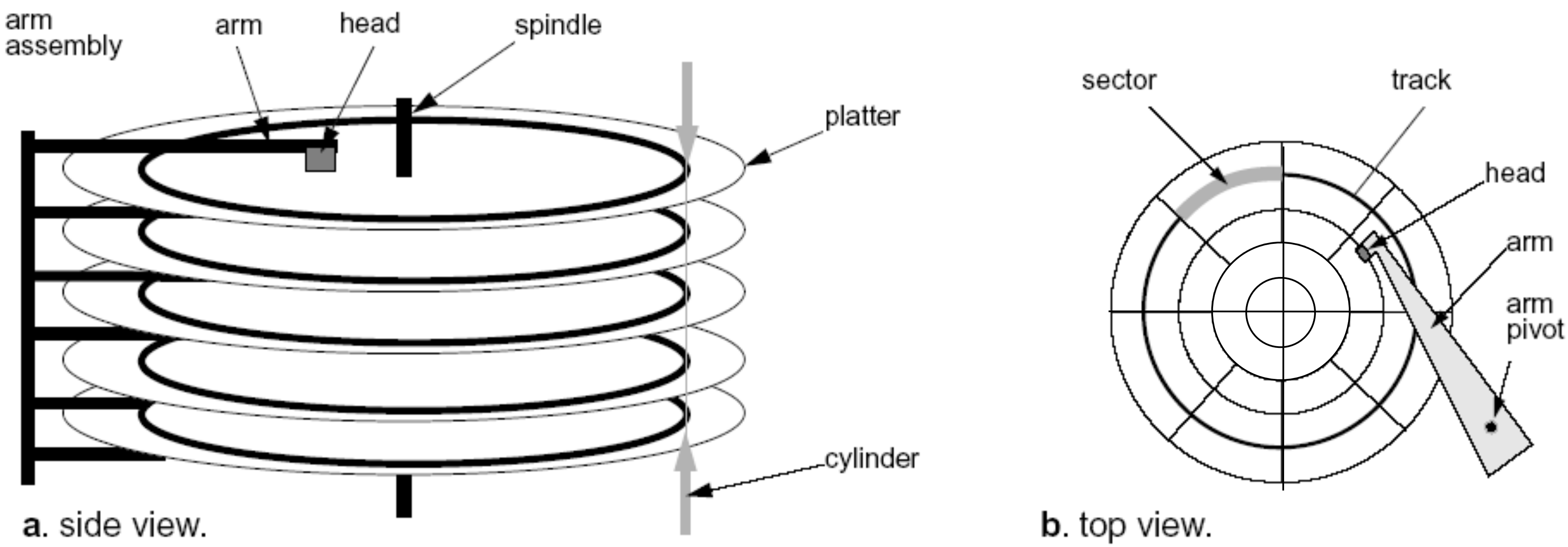
- Method
  - Use DRAM to cache recently accessed blocks
    - Typically a disk has 64-128 MB blocks
    - Some of the RAM space stores “firmware” (an embedded OS)
  - Blocks are replaced usually in an LRU order + “tracks”
- Pros
  - Good for reads if accesses have locality
- Cons
  - Doesn’t really help with writes since they need to be reliable

Disk Arm and Head

- Disk arm
  - A disk arm carries disk heads
- Disk head
  - Mounted on an actuator
  - Read/write on disk surface
- Read/write operation
  - Read/write with (track, sector)
  - Seek the right cylinder (tracks)
  - Wait until the sector comes
  - Perform read/write



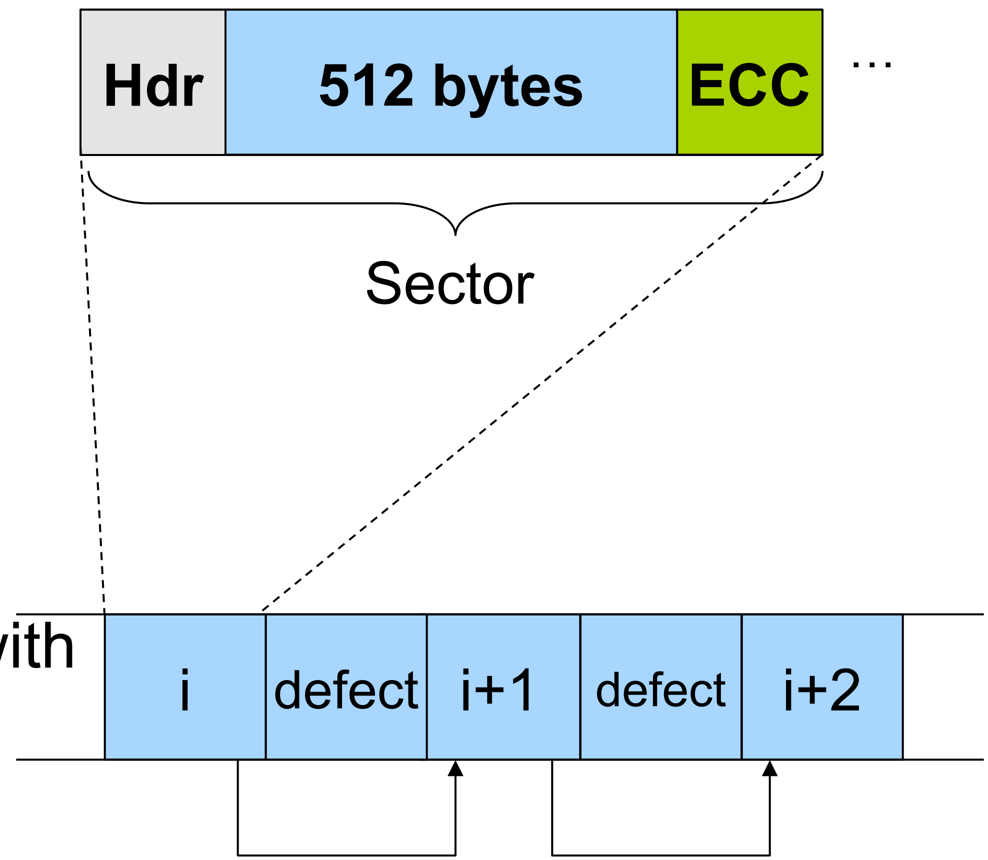
Mechanical Component of A Disk Drive



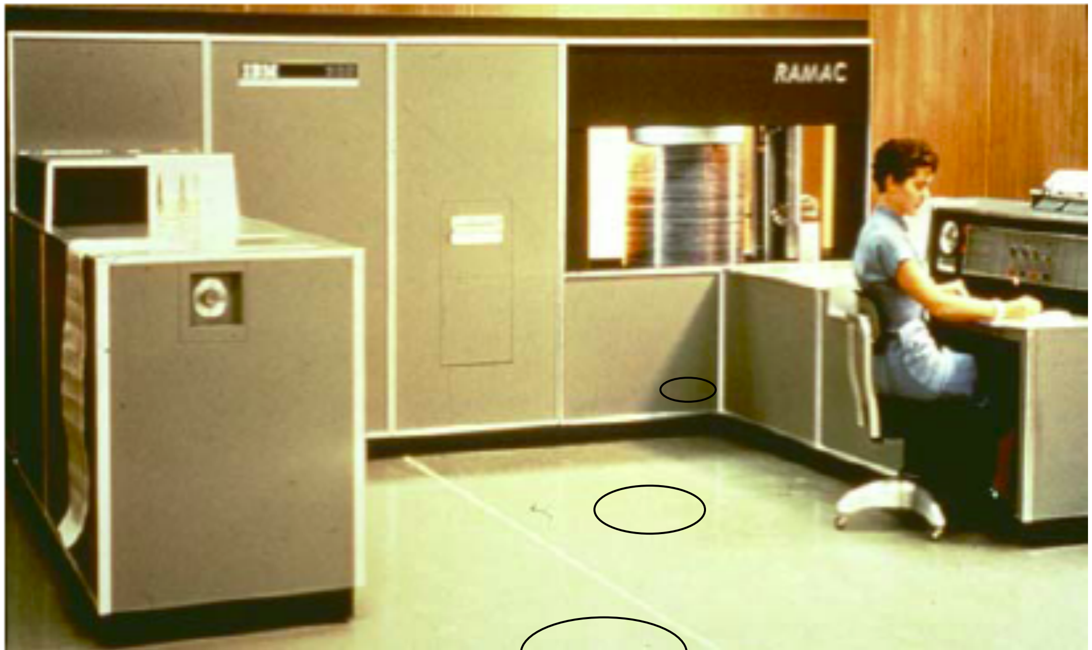
- Tracks
  - Concentric rings around disk surface, bits laid out serially along each track
- Cylinder
  - A track of the platter, 1000-5000 cylinders per zone, 1 spare per zone
- Sectors
  - Arc of track holding some min # of bytes, variable # sectors/track

Disk Sectors

- Where do they come from?
  - Formatting process
  - Logical maps to physical
- What is a sector?
  - Header (ID, defect flag, …)
  - Real space (e.g. 512 bytes)
  - Trailer (ECC code)
- What about errors?
  - Detect errors in a sector
  - Correct them with ECC
  - If not recoverable, replace it with a spare
  - Skip bad sectors in the future



Disks Were Large



First Disk:
IBM 305 RAMAC (1956)
5MB capacity
50 platters, each 24” diam

Storage Form Factors Are Changing

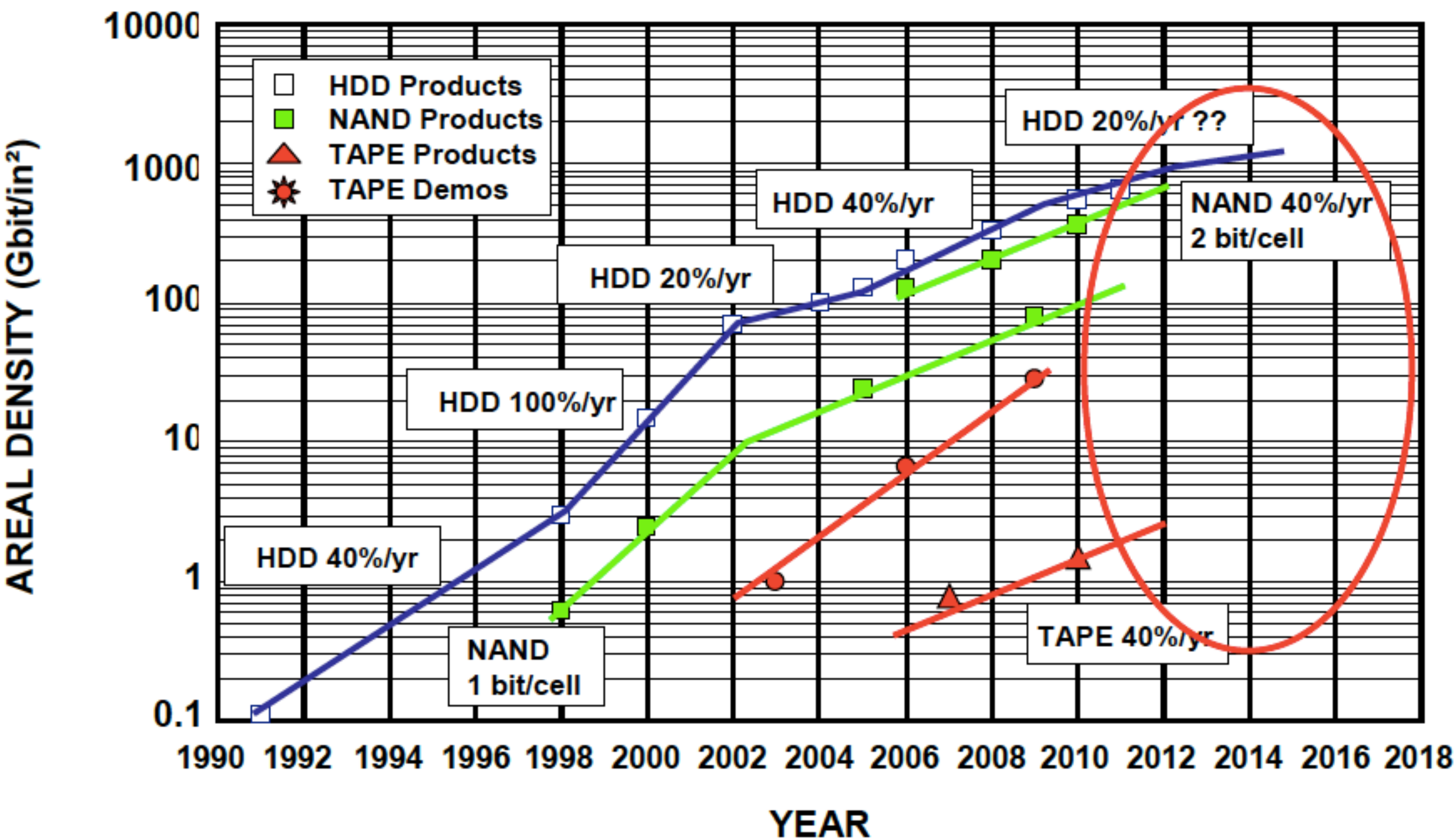
Form factor:
.5-1”· 4”· 5.7”
Storage:
0.5-6TB

Form factor:
.4-.7” · 2.7” · 3.9”
Storage:
0.5-2TB

Form factor: 24mm · 32mm · 2.1mm
Storage: 1-1TB

Form factor: PCI card
Storage: 0.5-10TB

Areal Density vs. Moore's Law



(Fontana, Decad, Hetzler, 2012)

50 Years (Mark Kryder at SNW 2006)

	IBM RAMAC (1956)	Seagate Momentus (2006)	Difference
Capacity	5MB	160GB	32,000
Areal Density	2K bits/in^2	130 Gbits/in^2	65,000,000
Disks	50 @ 24” diameter	2 @ 2.5” diameter	1 / 2,300
Price/MB	$1,000	$0.01	1 / 100,000
Spindle Speed	1,200 RPM	5,400 RPM	5
Seek Time	600 ms	10 ms	1 / 60
Data Rate	10 KB/s	44 MB/s	4,400
Power	5000 W	2 W	1 / 2,500
Weight	~ 1 ton	4 oz	1 / 9,000

Sample Disk Specs (from Seagate)

	Enterprise Performance	Desktop HDD
Capacity		
Formatted capacity (GB)	600	**4096**
Discs / heads	3 / 6	4 / 8
Sector size (bytes)	512	512
Performance		
External interface	STA	SATA
Spindle speed (RPM)	15,000	7,200
Average latency (msec)	**2.0**	4.16
Seek time, read/write (ms)	**3.5/3.9**	8.5/9.5
Track-to-track read/write (ms)	**0.2-0.4**	0.8/1.0
Transfer rate (MB/sec)	**138-258**	146
Cache size (MB)	128	64
Power		
Average / Idle / Sleep	8.5 / 6 / NA	7.5 / **5 / 0.75**
Reliability		
Recoverable read errors	1 per 10^{12} bits read	1 per 10^{10} bits read
Non-recoverable read errors	**1 per 10^{16} bits read**	1 per 10^{14} bits read

Disk Performance

- Seek
  - Position heads over cylinder, typically 3.5-9.5 ms
- Rotational delay
  - Wait for a sector to rotate underneath the heads
  - Typically 2 – 4 ms (7,200 – 15,000RPM, 4 - 8 ms per rotation)
- Transfer bytes
  - Transfer bandwidth is typically 70 -250 Mbytes/sec
- Example:
  - Performance of transfer 1 Kbytes of Desktop HDD, assuming BW = 100MB/sec, seek = 5ms, rotation = 4ms
  - Total time = 5ms + 4ms + 0.01ms = 9.01ms
  - What is the effective bandwidth?

More on Performance

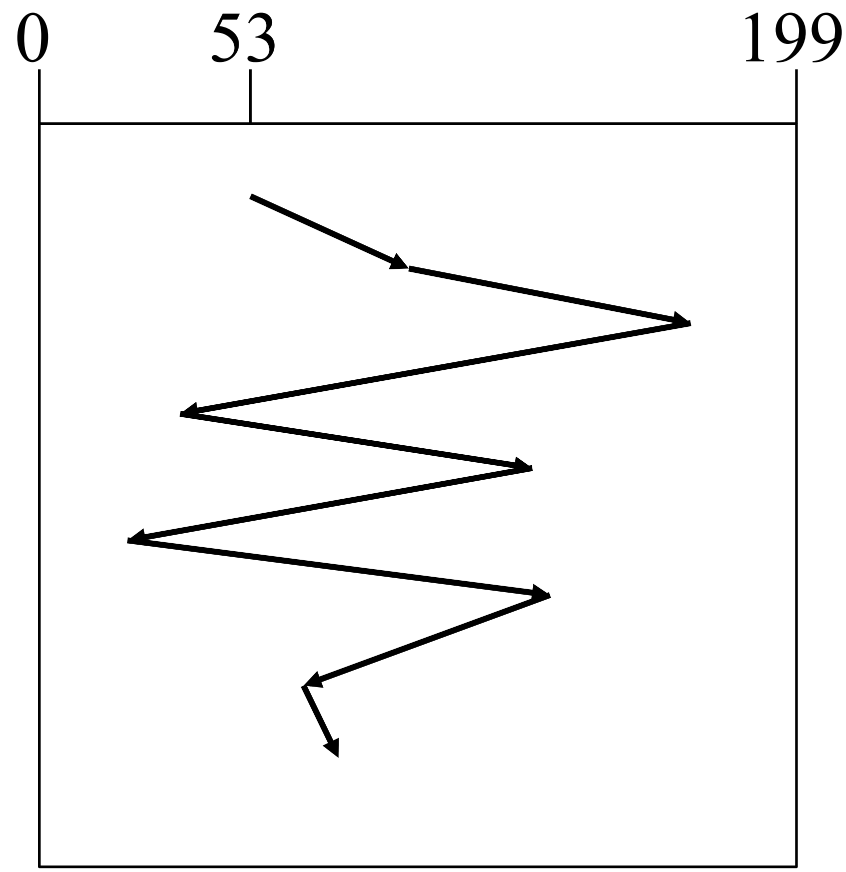
- What transfer size can get 90% of the disk bandwidth?
  - Assume Disk BW = 100MB/sec, avg rotation = 4ms, avg seek = 5ms
  - size / (size/BW + rotation + seek) = BW * 90%
  - size = BW * (rotation + seek) * 0.9 / (1 – 0.9)
    = 100MB * 0.009 * 0.9 / 0.1 = 8.1MB

Block Size (Kbytes)	% of Disk Transfer Bandwidth
9Kbytes	1%
100Kbytes	10%
0.9Mbytes	50%
8.1Mbytes	90%

- Seek and rotational times dominate the cost of small accesses
  - Disk transfer bandwidth are wasted
  - Need algorithms to reduce seek time

FIFO (FCFS) order

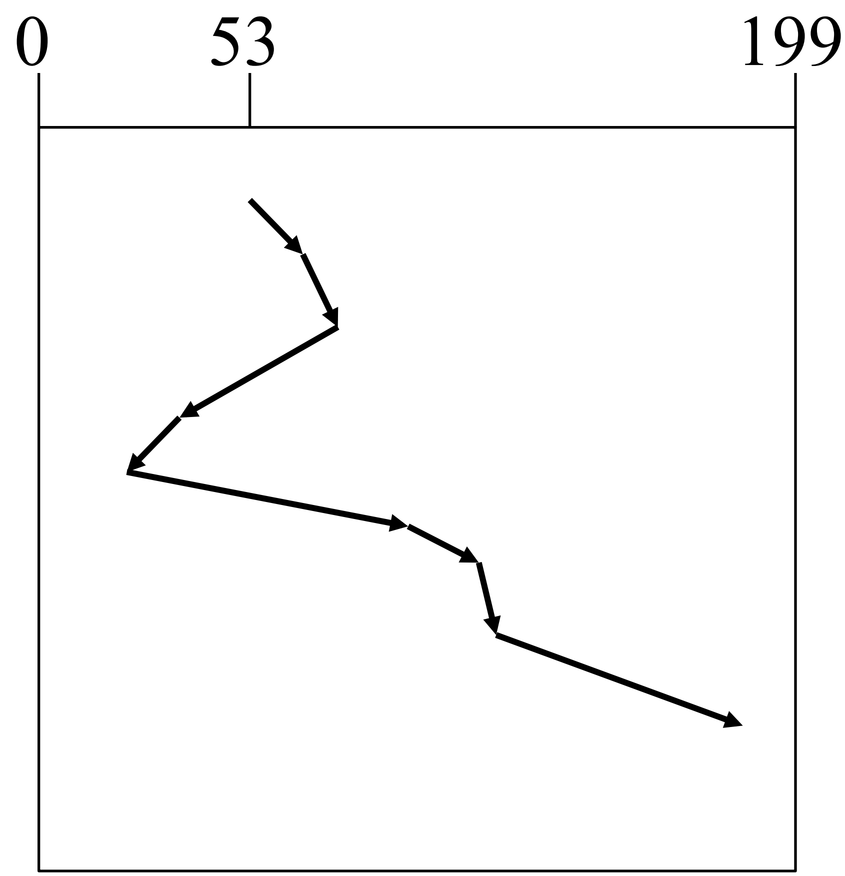
- Method
  - First come first serve
- Pros
  - Fairness among requests
  - In the order applications expect
- Cons
  - Arrival may be on random spots on the disk (long seeks)
  - Wild swing can happen



98, 183, 37, 122, 14, 124, 65, 67

SSTF (Shortest Seek Time First)

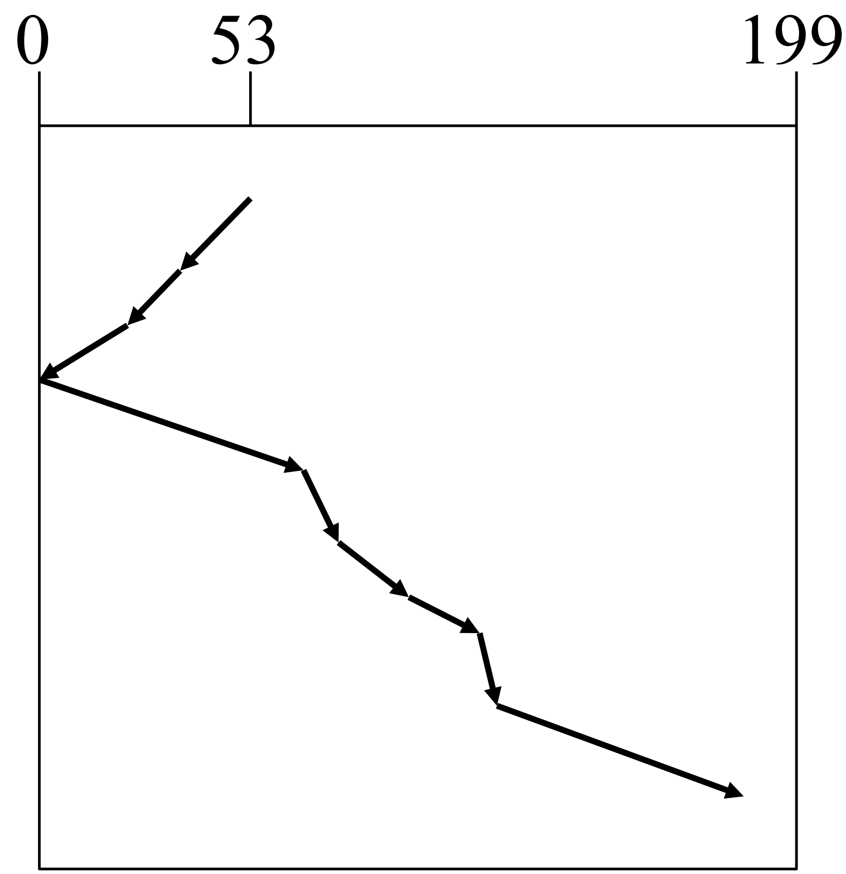
- Method
  - Pick the one closest on disk
  - Rotational delay is in calculation
- Pros
  - Try to minimize seek time
- Cons
  - Starvation
- Question
  - Is SSTF optimal?
  - Can we avoid the starvation?



98, 183, 37, 122, 14, 124, 65, 67
(65, 67, 37, 14, 98, 122, 124, 183)

Elevator (SCAN)

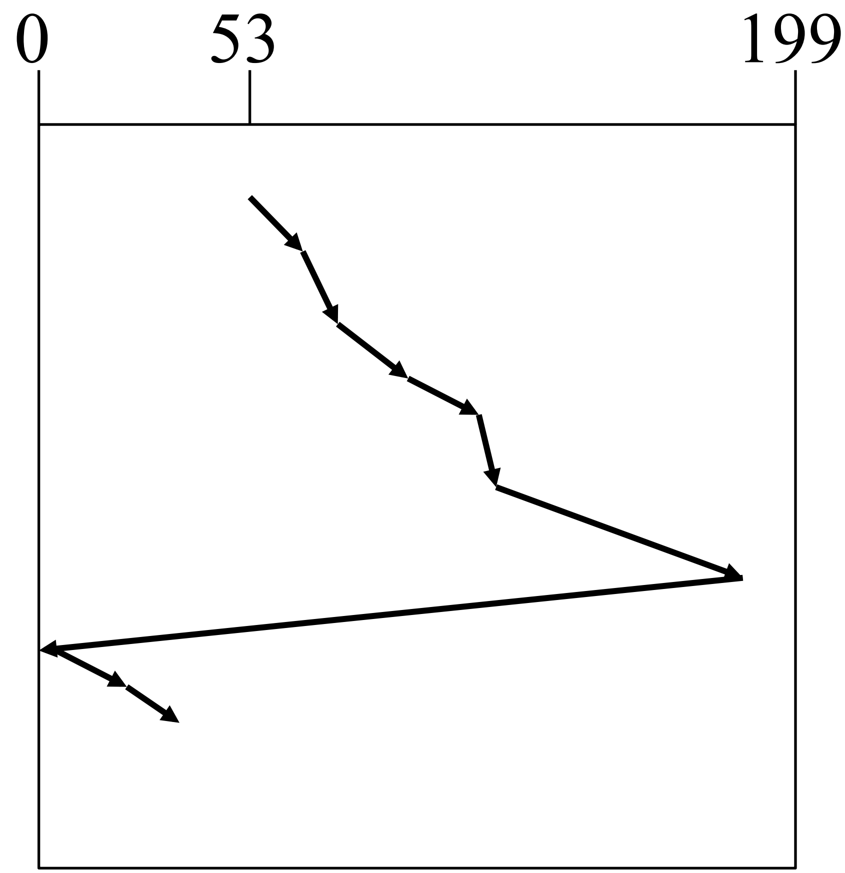
- Method
  - Take the closest request in the direction of travel
  - Real implementations do not go to the end (called LOOK)
- Pros
  - Bounded time for each request
- Cons
  - Request at the other end will take a while



98, 183, 37, 122, 14, 124, 65, 67
(37, 14, 65, 67, 98, 122, 124, 183)

C-SCAN (Circular SCAN)

- Method
  - Like SCAN
  - But, wrap around
  - Real implementation doesn't go to the end (C-LOOK)
- Pros
  - Uniform service time
- Cons
  - Do nothing on the return



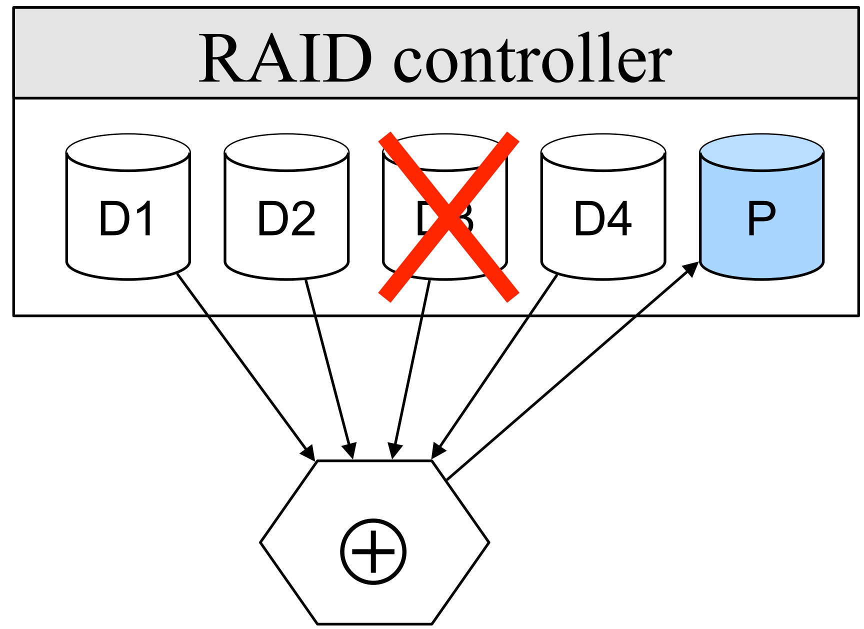
98, 183, 37, 122, 14, 124, 65, 67
(65, 67, 98, 122, 124, 183, 14, 37)

Discussions

- Which is your favorite?
  - FIFO
  - SSTF
  - SCAN
  - C-SCAN
- Disk I/O request buffering
  - Where would you buffer requests?
  - How long would you buffer requests?
- More advanced issues
  - Can the scheduling algorithm minimize both seek and rotational delays?

RAID (Redundant Array of Independent Disks)

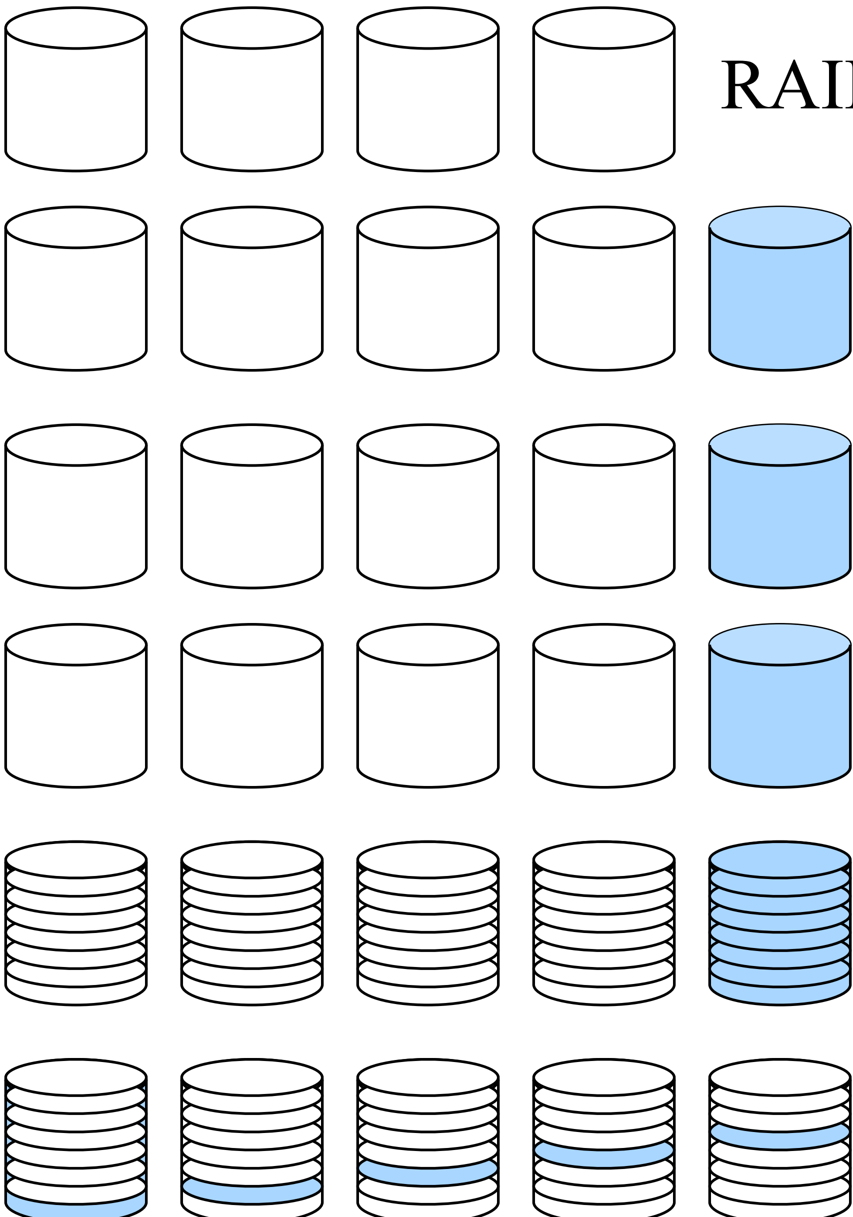
- Main idea
  - Compute XORs and store parity on disk P
  - Upon any failure, one can recover the block from using P and other disks
- Pros
  - Reliability
  - High bandwidth?
- Cons
  - Cost
  - The controller is complex



$P = D1 \oplus D2 \oplus D3 \oplus D4$

$D3 = D1 \oplus D2 \oplus P \oplus D4$

Synopsis of RAID Levels



RAID Level 0: Non redundant

RAID Level 1:
Mirroring

RAID Level 2:
Byte-interleaved, ECC

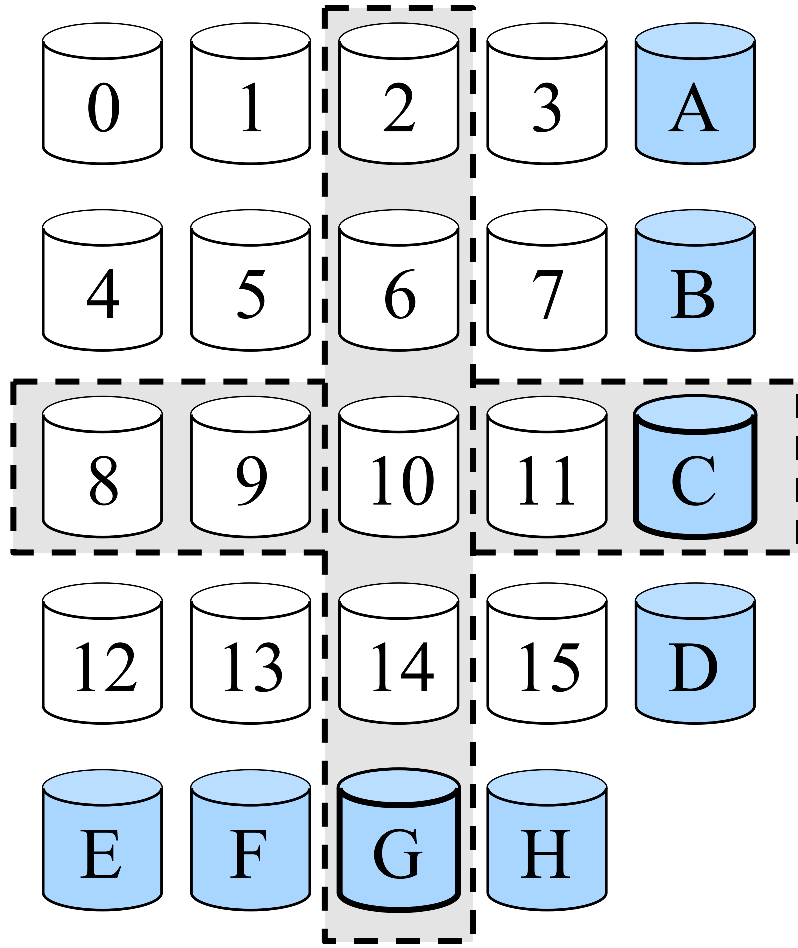
RAID Level 3:
Byte-interleaved, parity

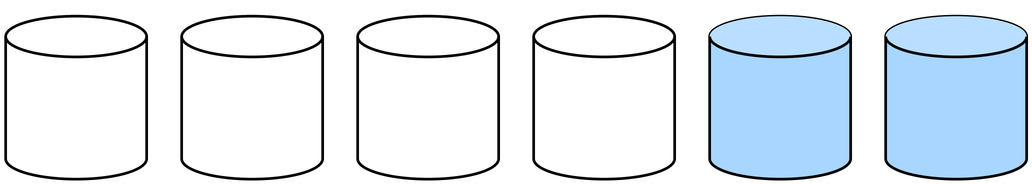
RAID Level 4:
Block-interleaved, parity

RAID Level 5:
Block-interleaved, distributed parity

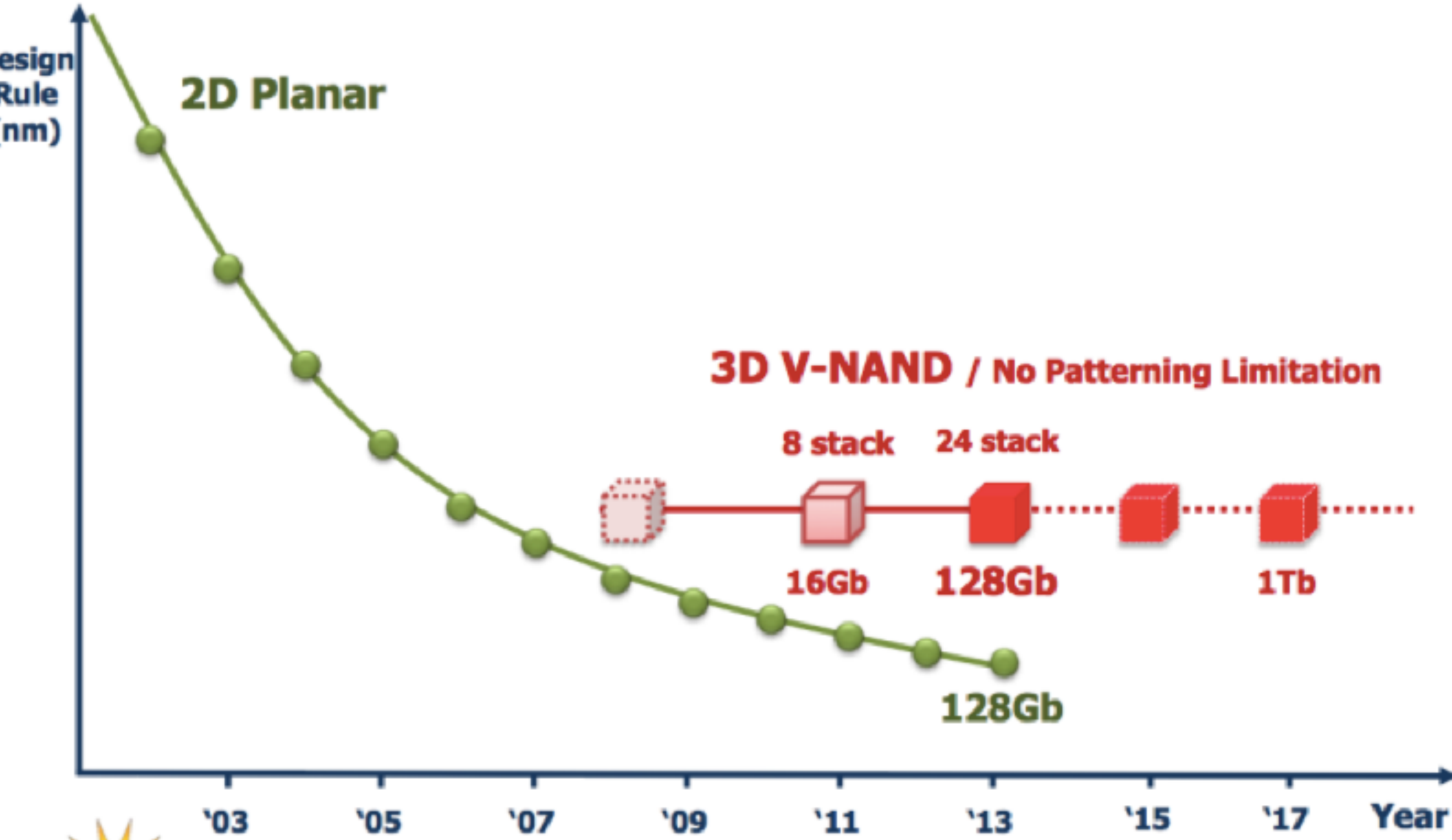
RAID Level 6 and Beyond

- Goals
  - Less computation and fewer updates per random write
  - Small amount of extra disk space
- Extended Hamming code
  - Remember Hamming code?
- Specialized Eraser Codes
  - IBM Even-Odd, NetApp RAID-DP, …
- Beyond RAID-6
  - Reed-Solomon codes, using MOD 4 equations
  - Can be generalized to deal with k (>2) disk failures



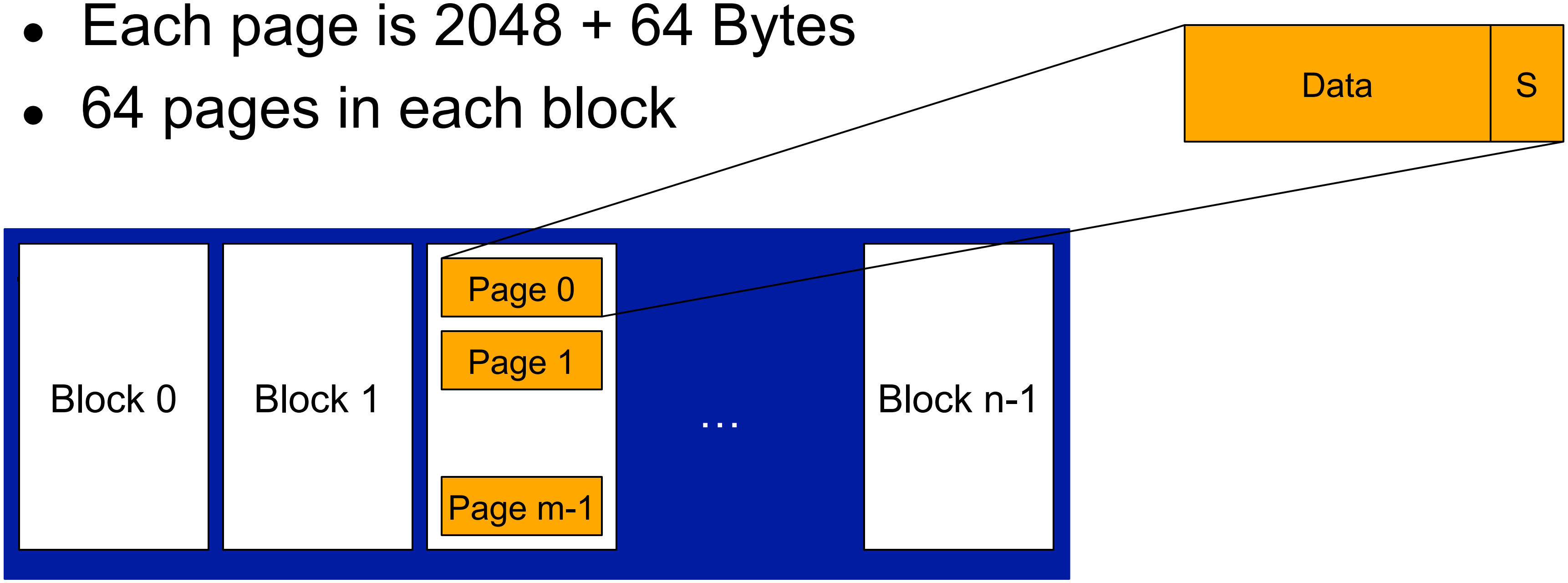


V-NAND Era for the Future



NAND Flash Memory

- High capacity
  - Single cell (more expensive, durable) vs. multiple cell
- Small block
  - Each page 512 + 16 Bytes
  - 32 pages in each block
- Large block
  - Each page is 2048 + 64 Bytes
  - 64 pages in each block



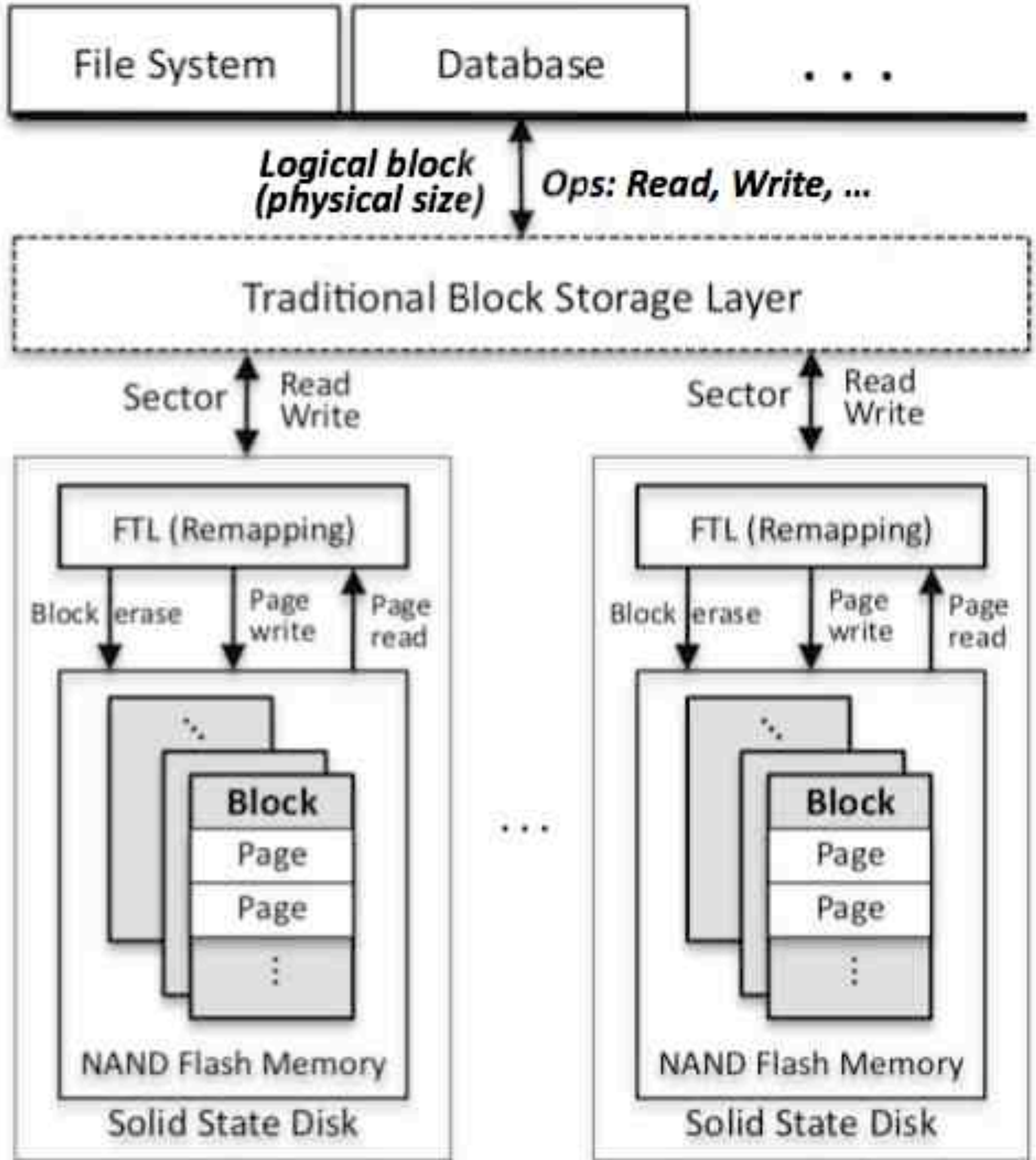
NAND Flash Memory Operations

- Speed
  - Read page: ~10-20 us
  - Write page: 20-200 us
  - Erase block: ~1-2 ms
- Limited performance
  - Can only write 0's, so erase (set all 1) then write
- Solution: Flash Translation Layer (FTL)
  - Map virtual page address to physical page address in flash controller
  - Keep erasing unused blocks
  - Remap to currently erased block to reduce latency

NAND Flash Lifetime

- Wear out limitations
  - ~50k to 100k writes / page (SLC – single level cell)
  - ~15k to 60k writes / page (MLC – multi-level cell)
- Wear Leveling:
  - Spread erases evenly across blocks, rather than using same block repeatedly
  - Question: Suppose write to cells evenly and 200,000 writes/sec, how long does it take to wear out 1,000M pages on SLC flash (50k/page)?
- Who does "wear leveling?"
  - Flash translation layer
  - File system design (later)

Flash Translation Layer



Example: Fusion I/O Flash Memory

- Flash Translation Layer (FTL) in device controller
  - Remapping
  - Wear-leveling
  - Write buffering
  - Log-structured file system (later)
- Performance
  - Fusion-IO Octal
  - 10TB
  - 6.7GB/s read
  - 3.9GB/s write
  - 45μs latency

Summary

- Disk is complex
- Disk real density has been on Moore's law curve
- Need large disk blocks to achieve good throughput
- System needs to perform disk scheduling
- RAID improves reliability and high throughput at a cost
- Flash memory has emerged at low and high ends