

# O'Caml Datatypes

COS 326

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# O'CamI So Far

- We have seen a number of basic types:
  - int
  - float
  - char
  - string
  - bool
- We have seen a few structured types:
  - pairs
  - tuples
  - options
  - lists
- In this lecture, we will see some more general ways to define our own new types and data structures

# Type Abbreviations

- We have already seen some type abbreviations:

```
type point = float * float
```

# Type Abbreviations

- We have already seen some type abbreviations:

```
type point = float * float
```

- These abbreviations can be helpful documentation:

```
let distance (p1:point) (p2:point) : float =  
  let square x = x *. x in  
  let (x1,y1) = p1 in  
  let (x2,y2) = p2 in  
  sqrt (square (x2 -. x1) +. square (y2 -. y1))
```

- But they add nothing of *substance* to the language
  - they are *equal* in every way to an existing type

# Type Abbreviations

- We have already seen some type abbreviations:

```
type point = float * float
```

- As far as O'Caml is concerned, you could have written:

```
let distance (p1:float*float)
             (p2:float*float) : float =
  let square x = x *. x in
  let (x1,y1) = p1 in
  let (x2,y2) = p2 in
  sqrt (square (x2 -. x1) +. square (y2 -. y1))
```

- Since the types are equal, you can *substitute* the definition for the name wherever you want
  - we have not added any new data structures

# DATA TYPES

# Data types

- O'CamI provides a general mechanism called a **data type** for defining new data structures that consist of many alternatives

```
type my_bool = Tru | Fal
```

a **value** with type **my\_bool** is one of two things:

- **Tru**, or
- **Fal**

read the "|" as "or"

# Data types

- O'CamI provides a general mechanism called a **data type** for defining new data structures that consist of many alternatives

```
type my_bool = Tru | Fal
```

Tru and Fal are called  
"constructors"

a **value** with type **my\_bool**  
is one of two things:

- **Tru**, or
- **Fal**

read the "|" as "or"

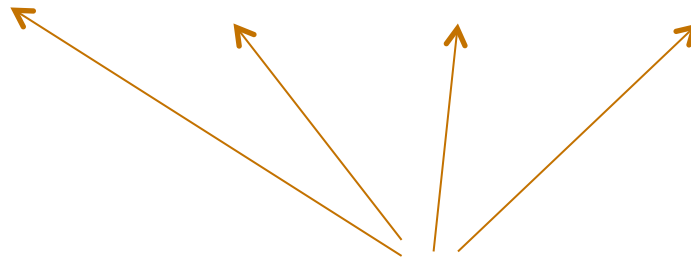


# Data types

- O'CamL provides a general mechanism called a **data type** for defining new data structures that consist of many alternatives

```
type my_bool = Tru | Fal
```

```
type color = Blue | Yellow | Green | Red
```



there's no need to stop  
at 2 cases; define as many  
alternatives as you want

# Data types

- O'CamL provides a general mechanism called a **data type** for defining new data structures that consist of many alternatives

```
type my_bool = Tru | Fal
```

```
type color = Blue | Yellow | Green | Red
```

- Creating values:

```
let b1 : my_bool = Tru
let b2 : my_bool = Fal
let c1 : color = Yellow
let c2 : color = Red
```

use constructors to create values



# Data types

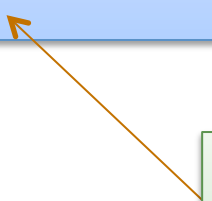
```
type color = Blue | Yellow | Green | Red
```

```
let c1 : color = Yellow
```

```
let c2 : color = Red
```

- Using data type values:

```
let print_color (c:color) : unit =  
  match c with  
  | Blue ->  
  | Yellow ->  
  | Green ->  
  | Red ->
```



use pattern matching to determine which color you have; act accordingly

# Data types

```
type color = Blue | Yellow | Green | Red
```

```
let c1 : color = Yellow
```

```
let c2 : color = Red
```

- Using data type values:

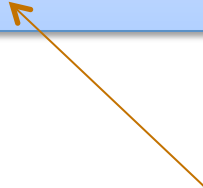
```
let print_color (c:color) : unit =  
  match c with  
  | Blue -> print_string "blue"  
  | Yellow -> print_string "yellow"  
  | Green -> print_string "green"  
  | Red -> print_string "red"
```

# Data types

```
type color = Blue | Yellow | Green | Red
```

oops!:

```
let print_color (c:color) : unit =  
  match c with  
  | Blue -> print_string "blue"  
  | Yellow -> print_string "yellow"  
  | Red -> print_string "red"
```



**Warning 8:** this pattern-matching is not exhaustive.  
Here is an example of a value that is not matched:  
Green

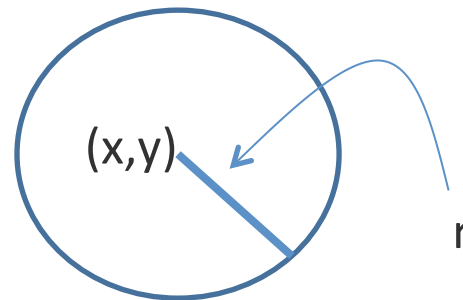
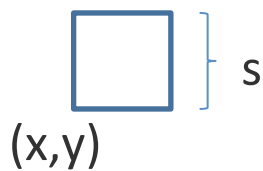
# Data Types Can Carry Additional Values

- Data types are more than just enumerations of constants:

```
type point = float * float

type simple_shape =
  | Circle of point * float
  | Square of point * float
```

- Read as: a **simple\_shape** is either:
  - a **Circle**, which contains a **pair** of a **point** and **float**, or
  - a **Square**, which contains a **pair** of a **point** and **float**



# Data Types Can Carry Additional Values

- Data types are more than just enumerations of constants:

```
type point = float * float
```

```
type simple_shape =  
  Circle of point * float  
| Square of point * float
```

```
let origin : point = (0.0, 0.0)
```

```
let circ1 : simple_shape = Circle (origin, 1.0)
```

```
let circ2 : simple_shape = Circle ((1.0, 1.0), 5.0)
```

```
let square : simple_shape = Square (origin, 2.3)
```

# Data Types Can Carry Additional Values

- Data types are more than just enumerations of constants:

```
type point = float * float

type simple_shape =
  | Circle of point * float
  | Square of point * float

let simple_area (s:simple_shape) : float =
  match s with
  | Circle (_, radius) -> 3.14 *. radius *. radius
  | Square (_, side) -> side *. side
```



# Compare

- Data types are more than just enumerations of constants:

```
type point = float * float

type simple_shape =
  | Circle of point * float
  | Square of point * float

let simple_area (s:simple_shape) : float =
  match s with
  | Circle (_, radius) -> 3.14 *. radius *. radius
  | Square (_, side) -> side *. side
```

```
type my_shape = point * float

let simple_area (s:my_shape) : float =
  (3.14 *. radius *. radius) ?? or ?? (side *. side)
```

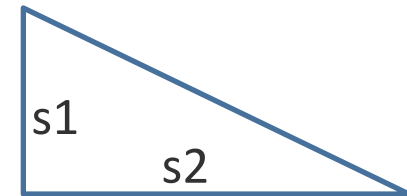
# More General Shapes

```
type point = float * float  
  
type shape =  
  Square of float  
  | Ellipse of float * float  
  | RtTriangle of float * float  
  | Polygon of point list
```

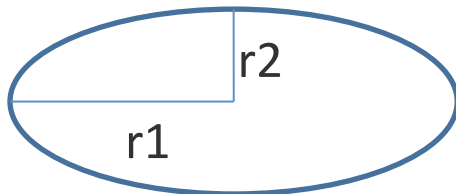
Square  $s =$



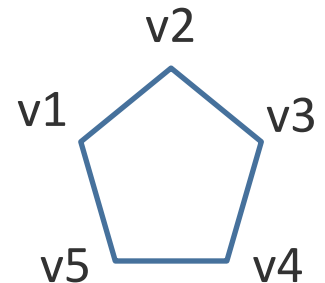
RtTriangle ( $s_1, s_2$ ) =



Ellipse ( $r_1, r_2$ ) =



RtTriangle [ $v_1; \dots; v_5$ ] =



# More General Shapes

```
type point = float * float
type radius = float
type side = float

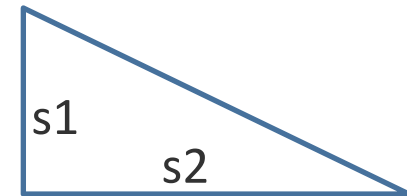
type shape =
  | Square of side
  | Ellipse of radius * radius
  | RtTriangle of side * side
  | Polygon of point list
```

Type abbreviations can aid readability

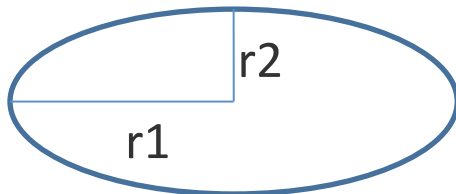
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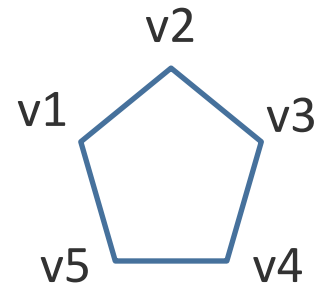
RtTriangle ( $s_1, s_2$ ) =



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RtTriangle [ $v_1; \dots; v_5$ ] =



# More General Shapes

```
type point = float * float
type radius = float
type side = float

type shape =
  Square of side
  | Ellipse of radius * radius
  | RtTriangle of side * side
  | Polygon of point list
```

Square builds a shape  
from a single side

RtTriangle builds a shape  
from a pair of sides

```
let sq    : shape = Square 17.0
let ell   : shape = Ellipse (1.0, 2.0)
let rt    : shape = RtTriangle (1.0, 1.0)
let poly  : shape = Polygon [(0., 0.); (1., 0.); (0.; 1.)]
```

they are all shapes;  
they are constructed in  
different ways

Polygon builds a shape  
from a list of points  
(where each point is itself a pair)

# More General Shapes

```
type point = float * float
type radius = float
type side = float

type shape =
  Square of side
  | Ellipse of radius * radius
  | RtTriangle of side * side
  | Polygon of point list
```

```
let area (s : shape) : float =
  match s with
  | Square s ->
  | Ellipse (r1, r2)->
  | RtTriangle (s1, s2) ->
  | Polygon ps ->
```

a data type also defines  
a pattern for matching

# More General Shapes

```
type point = float * float
type radius = float
type side = float

type shape =
  Square of side
  | Ellipse of radius * radius
  | RtTriangle of side * side
  | Polygon of point list
```

a data type also defines  
a pattern for matching

```
let area (s : shape) : float =
  match s with
  | Square s ->
  | Ellipse (r1, r2) ->
  | RtTriangle (s1, s2) ->
  | Polygon ps ->
```

**Square** carries a value  
with type **float** so **s** is  
a pattern for float values

**RtTriangle** carries a value  
with type **float \* float**  
so **(s1, s2)** is a pattern  
for that type

# More General Shapes

```
type point = float * float
type radius = float
type side = float

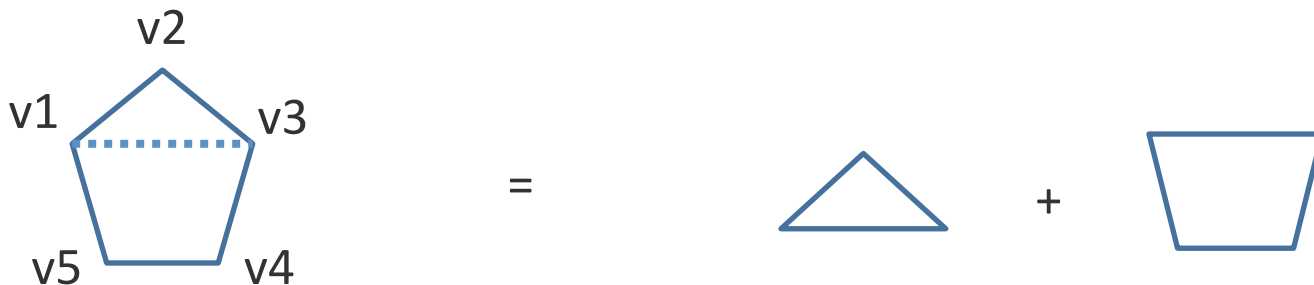
type shape =
  Square of side
  | Ellipse of radius * radius
  | RtTriangle of side * side
  | Polygon of point list
```

```
let area (s : shape) : float =
  match s with
  | Square s -> s *. s
  | Ellipse (r1, r2) -> r1 *. r2
  | RtTriangle (s1, s2) -> s1 *. s2 /. 2.
  | Polygon ps -> ???
```

a data type also defines  
a pattern for matching

# Computing Area

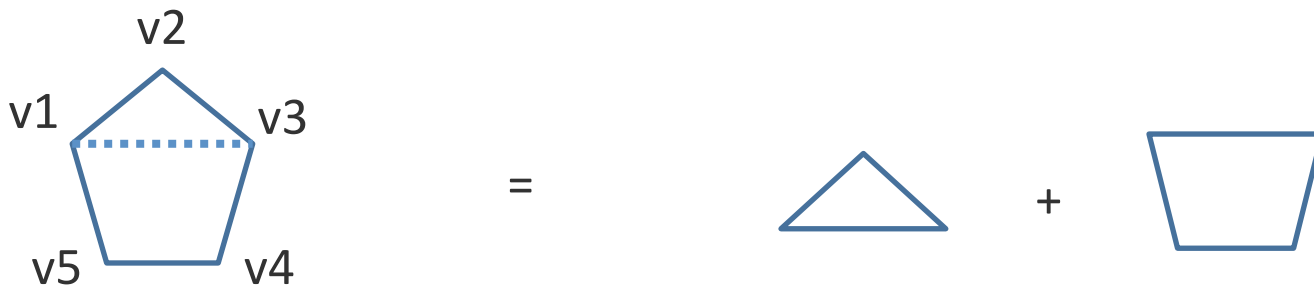
- How do we compute polygon area?
- For convex polygons:
  - Case: the polygon has fewer than 3 points:
    - it has 0 area! (it is a line or a point or nothing at all)
  - Case: the polygon has 3 or more points:
    - Compute the area of the triangle formed by the first 3 vertices
    - Delete the second vertex to form a new polygon
    - Sum the area of the triangle and the new polygon





# Computing Area

- How do we compute polygon area?
- For convex polygons:
  - **Case: the polygon has fewer than 3 points:**
    - it has 0 area! (it is a line or a point or nothing at all)
  - **Case: the polygon has 3 or more points:**
    - Compute the area of the triangle formed by the first 3 vertices
    - Delete the second vertex to form a new polygon
    - Sum the area of the triangle and the new polygon
- Note: This is a beautiful **inductive algorithm**:
  - the area of a polygon with  $n$  points is computed in terms of a smaller polygon with only  $n-1$  points!

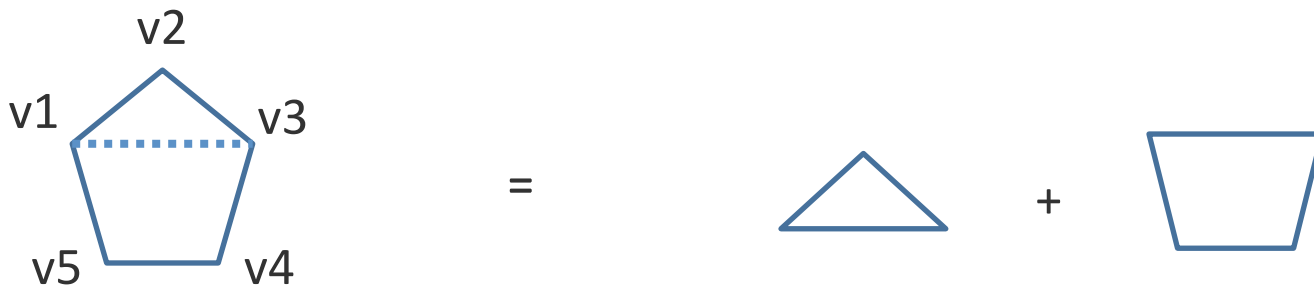


# Computing Area

```
let area (s : shape) : float =  
  match s with  
  | Square s -> s *. s  
  | Ellipse (r1, r2) -> r1 *. r2  
  | RtTriangle (s1, s2) -> s1 *. s2 /. 2.  
  | Polygon ps -> poly_area ps
```

This pattern says the list has at least 3 items

```
let poly_area (ps : point list) : float =  
  match ps with  
  | p1 :: p2 :: p3 :: tail ->  
    tri_area p1 p2 p3 +. poly_area (p1::p3::tail)  
  | _ -> 0.
```



# Computing Area

```
let tri_area (p1:point) (p2:point) (p3:point) : float =
  let a = distance p1 p2 in
  let b = distance p2 p3 in
  let c = distance p3 p1 in
  let s = 0.5 *. (a +. b +. c) in
  sqrt (s *. (s -. a) *. (s -. b) *. (s -. c))
```

```
let rec poly_area (ps : point list) : float =
  match ps with
  | p1 :: p2 :: p3 :: tail ->
    tri_area p1 p2 p3 +. poly_area (p1::p3::ps)
  | _ -> 0.
```

```
let area (s : shape) : float =
  match s with
  | Square s -> s *. s
  | Ellipse (r1, r2)-> r1 *. r2
  | RtTriangle (s1, s2) -> s1 *.s2 /.2.
  | Polygon ps -> poly_area ps
```

# **INDUCTIVE DATA TYPES**

# Inductive data types

- We can use data types to define inductive data
- A binary tree is:
  - a **Leaf** containing no data
  - a **Node** containing a **key**, a **value**, a left subtree and a right subtree

# Inductive data types

- We can use data types to define inductive data
- A binary tree is:
  - a **Leaf** containing no data
  - a **Node** containing a **key**, a **value**, a left subtree and a right subtree

```
type key = string
type value = int

type tree =
  Leaf
| Node of key * value * tree * tree
```

# Inductive data types

```
type key = int
type value = string

type tree =
  | Leaf
  | Node of key * value * tree * tree
```

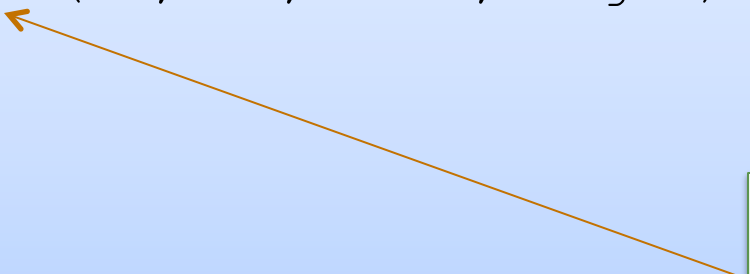
```
let rec insert (t:tree) (k:key) (v:value) : tree =
```

# Inductive data types

```
type key = int
type value = string

type tree =
  Leaf
| Node of key * value * tree * tree
```

```
let rec insert (t:tree) (k:key) (v:value) : tree =
  match t with
  | Leaf ->
  | Node (k', v', left, right) ->
```



Again, the type definition specifies the cases you must consider



# Inductive data types

```
type key = int
type value = string

type tree =
  Leaf
| Node of key * value * tree * tree
```

```
let rec insert (t:tree) (k:key) (v:value) : tree =
  match t with
  | Leaf -> Node (k, v, Leaf, Leaf)
  | Node (k', v', left, right) ->
```

# Inductive data types

```
type key = int
type value = string

type tree =
  Leaf
| Node of key * value * tree * tree
```

```
let rec insert (t:tree) (k:key) (v:value) : tree =
  match t with
  | Leaf -> Node (k, v, Leaf, Leaf)
  | Node (k', v', left, right) ->
    if k < k' then
      Node (k', v', insert left k v, right)
    else if k > k' then
      Node (k', v', left, insert right k v)
    else
      Node (k, v, left, right)
```

# Inductive data types

```
type key = int
type value = string

type tree =
  Leaf
| Node of key * value * tree * tree
```

```
let rec insert (t:tree) (k:key) (v:value) : tree =
  match t with
  | Leaf -> Node (k, v, Leaf, Leaf)
  | Node (k', v', left, right) ->
    if k < k' then
      Node (k', v', insert left k v, right)
    else if k > k' then
      Node (k', v', left, insert right k v)
    else
      Node (k, v, left, right)
```

# Inductive data types

```
type key = int
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```
let rec insert (t:tree) (k:key) (v:value) : tree =
  match t with
  | Leaf -> Node (k, v, Leaf, Leaf)
  | Node (k', v', left, right) ->
    if k < k' then
      Node (k', v', insert left k v, right)
    else if k > k' then
      Node (k', v', left, insert right k v)
    else
      Node (k, v, left, right)
```

# Inductive data types: Another Example

- Recall, we used the type "int" to represent natural numbers
  - but that was kind of broken: it also contained negative numbers
  - we had to use a dynamic test to guard entry to a function:

```
let double (n : int) : int =  
  if n < 0 then  
    raise (Failure "negative input!")  
  else  
    double_nat n
```

- it would be nice if there was a way to define the natural numbers **exactly**, and use OCaml's type system to guarantee no client ever attempts to double a negative number

# Inductive data types

- Recall, a natural number  $n$  is either:
  - zero, or
  - $m + 1$
- We use a data type to represent this definition exactly:

# Inductive data types

- Recall, a natural number  $n$  is either:
  - zero, or
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```
type nat = Zero | Next of nat
```

# Inductive data types

- Recall, a natural number  $n$  is either:
  - zero, or
  - $m + 1$
- We use a data type to represent this definition exactly:

```
type nat = Zero | Next of nat

let rec nat_to_int (n : nat) : int =
  match n with
  | Zero -> 0
  | Next n -> 1 + nat_to_int n
```



# Inductive data types

- Recall, a natural number  $n$  is either:
  - zero, or
  - $m + 1$
- We use a data type to represent this definition exactly:

```
type nat = Zero | Next of nat

let rec nat_to_int (n : nat) : int =
  match n with
  | Zero -> 0
  | Next n -> 1 + nat_to_int n

let rec double_nat (n : nat) : nat =
  match n with
  | Zero -> Zero
  | Next m -> Next (Next (double_nat m))
```

**AN EXERCISE IN TYPE DESIGN**

# Example Type Design

- A **GML document** consists of:
  - a list of **elements**
- An **element** is either:
  - a **word** or **markup** applied to an element
- **Markup** is either:
  - **italicize**, **bold**, or a **font name**

# Example Type Design

- A **GML document** consists of:
  - a list of **elements**
- An **element** is either:
  - a **word** or **markup** applied to an element
- **Markup** is either:
  - **italicize**, **bold**, or a **font name**

```
type markup = Ital | Bold | Font of string
```

```
type elt =  
  Words of string list  
| Formatted of markup * elt
```

```
type doc = elt list
```

# Example Data

```
type markup = Ital | Bold | Font of string
```

```
type elt =  
  Words of string list  
| Formatted of markup * elt
```

```
type doc = elt list
```

```
let d = [ Formatted (Bold,  
  Formatted (Font "Arial",  
  Words ["Chapter"; "One"]));  
  
  Words ["It"; "was"; "a"; "dark";  
  "&"; "stormy"; "night."; "A"];  
  
  Formatted (Ital, Words["shot"]);  
  
  Words ["rang"; "out."] ];;
```

# Challenge

- Change all of the “**Arial**” fonts in a document to “**Courier**”.
- Of course, when we program functionally, we implement *change* via a function that
  - receives one data structure as input
  - builds a new (different) data structure as an output

# Challenge

- Change all of the “**Arial**” fonts in a document to “**Courier**”.

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type markup = Ital | Bold | Font of string

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# Challenge

- Change all of the “Arial” fonts in a document to “Courier”.

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type markup = Ital | Bold | Font of string

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```

- Technique: approach the problem top down, work on **doc** first:

```
let rec chfonts (elts:doc) : doc =
```



# Challenge

- Change all of the “**Arial**” fonts in a document to “**Courier**”.

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type markup = Ital | Bold | Font of string

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```

- Technique: approach the problem top down, work on **doc** first:

```
let rec chfonts (elts:doc) : doc =
  match elts with
  | [] ->
  | hd::tl ->
```

# Challenge

- Change all of the “**Arial**” fonts in a document to “**Courier**”.

```
type markup = Ital | Bold | Font of string

type elt =
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- Technique: approach the problem top down, work on **doc** first:

```
let rec chfonts (elts:doc) : doc =
  match elts with
  | [] -> []
  | hd::tl -> (chfont hd)::(chfonts tl)
```

# Changing fonts in an element

- Change all of the “**Arial**” fonts in a document to “**Courier**”.

```
type markup = Ital | Bold | Font of string

type elt =
  Words of string list
| Formatted of markup * elt

type doc = elt list
```

- Next work on changing the font of an **element**:

```
let rec chfont (e:elt) : elt =
```

# Changing fonts in an element

- Change all of the “**Arial**” fonts in a document to “**Courier**”.

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```

- Next work on changing the font of an **element**:

```
let rec chfont (e:elt) : elt =
  match e with
  | Words ws ->
  | Formatted(m,e) ->
```

# Changing fonts in an element

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type elt =
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- Next work on changing the font of an **element**:

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let rec chfont (e:elt) : elt =
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- Next work on changing the font of an **element**:

```
let rec chfont (e:elt) : elt =
  match e with
  | Words ws -> Words ws
  | Formatted(m,e) -> Formatted(chmarkup m, chfont e)
```

# Changing fonts in an element

- Change all of the “**Arial**” fonts in a document to “**Courier**”.

```
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  Words of string list
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type doc = elt list
```

- Next work on changing a **markup**:

```
let chmarkup (m:markup) : markup =
```

# Changing fonts in an element

- Change all of the “**Arial**” fonts in a document to “**Courier**”.

```
type markup = Ital | Bold | Font of string

type elt =
  Words of string list
| Formatted of markup * elt

type doc = elt list
```

- Next work on changing a **markup**:

```
let chmarkup (m:markup) : markup =
  match m with
  | Font "Arial" -> Font "Courier"
  | _ -> m
```



# Summary: Changing fonts in an element

- Change all of the “**Arial**” fonts in a document to “**Courier**”
- Lesson: function structure follows type structure

```
let chmarkup (m:markup) : markup =
  match m with
  | Font "Arial" -> Font "Courier"
  | _ -> m

let rec chfont (e:elt) : elt =
  match e with
  | Words ws -> Words ws
  | Formatted(m,e) -> Formatted(chmarkup m, chfont e)

let rec chfonts (elts:doc) : doc =
  match elts with
  | [] -> []
  | hd::tl -> (chfont hd)::(chfonts tl)
```

# Poor Style

- Consider again our definition of markup and markup change:

```
type markup =  
  Ital | Bold | Font of string  
  
let chmarkup (m:markup) : markup =  
  match m with  
  | Font "Arial" -> Font "Courier"  
  | _ -> m
```

# Poor Style

- What if we make a change:

```
type markup =  
  Ital | Bold | Font of string | TTFont of string  
  
let chmarkup (m:markup) : markup =  
  match m with  
  | Font "Arial" -> Font "Courier"  
  | _ -> m
```

the underscore silently catches all possible alternatives

this may not be what we want -- perhaps there is an Arial TT font

it is better if we are alerted of all functions whose implementation may need to change

# Better Style

- Original code:

```
type markup =  
  Ital | Bold | Font of string  
  
let chmarkup (m:markup) : markup =  
  match m with  
  | Font "Arial" -> Font "Courier"  
  | Ital | Bold -> m
```

# Better Style

- Updated code:

```
type markup =  
  Ital | Bold | Font of string | TTFont of string  
  
let chmarkup (m:markup) : markup =  
  match m with  
  | Font "Arial" -> Font "Courier"  
  | Ital | Bold -> m
```

```
..match m with  
  | Font "Arial" -> Font "Courier"  
  | Ital | Bold -> m..
```

Warning 8: this pattern-matching is not exhaustive.  
Here is an example of a value that is not matched:  
TTFont \_

# Better Style

- Updated code, fixed:

```
type markup =  
  Ital | Bold | Font of string | TTFont of string  
  
let chmarkup (m:markup) : markup =  
  match m with  
  | Font "Arial" -> Font "Courier"  
  | TTFont "Arial" -> TTFont "Courier"  
  | TTFont s -> TTFont s  
  | Ital | Bold -> m
```

- **Lesson:** use the type checker where possible to help you maintain your code

# A couple of practice problems

- Write a function that gets rid of immediately redundant markup in a document.
  - `Formatted(Ital, Formatted(Ital,e))` can be simplified to `Formatted(Ital,e)`
  - write maps and folds over markups
- Design a datatype to describe bibliography entries for publications. Some publications are journal articles, others are books, and others are conference papers. Journals have a name, number and issue; books have an ISBN number; All of these entries should have a title and author.
  - design a sorting function
  - design maps and folds over your bibliography entries

# To Summarize

- Design recipe for writing Ocaml code:
  - write down English specifications
    - try to break problem into obvious sub-problems
  - write down some sample test cases
  - write down the signature (types) for the code
  - use the signature to guide construction of the code:
    - tear apart inputs using pattern matching
      - make sure to cover all of the cases! (Ocaml will tell you)
    - handle each case, building results using data constructor
      - this is where human intelligence comes into play
      - the “skeleton” given by types can almost be done automatically!
    - clean up your code
  - use your sample tests (and ideally others) to ensure correctness



**END**