

Recall that the diameter of our body is  $\text{poly}(n)$  since we started with a body whose diameter was  $n^2$  and then we placed a grid of size  $1/n^2$  or so. Combining all the above information, we get that

$$\phi \geq \frac{1}{\text{poly}(n)}.$$

Therefore, the mixing time is  $\mathcal{O}(\frac{1}{\phi^2} \log N) = \mathcal{O}(\text{poly}(n))$ .

## 4 Dimension Reduction

Now we describe a central result of high-dimensional geometry (at least when distances are measured in the  $\ell_2$  norm). Problem: Given  $n$  points  $z^1, z^2, \dots, z^n$  in  $\mathfrak{R}^n$ , we would like to find  $n$  points  $u^1, u^2, \dots, u^n$  in  $\mathfrak{R}^m$  where  $m$  is of low dimension (compared to  $n$ ) and the metric restricted to the points is almost preserved, namely:

$$\|z^i - z^j\|_2 \leq \|u^i - u^j\|_2 \leq (1 + \epsilon) \|z^i - z^j\|_2 \quad \forall i, j. \quad (2)$$

The following main result is by Johnson & Lindenstrauss :

**THEOREM 4**

*In order to ensure (2),  $m = \mathcal{O}(\frac{\log n}{\epsilon^2})$  suffices.*

Note: In class, we used the notation of  $k$  vectors  $z^1 \dots z^k$  in  $\mathfrak{R}^n$ , but we can always embed the  $k$  vectors in a  $k$ -dimensional space, so here I assume that  $n = k$  and use only  $n$ .

**PROOF:**

Choose  $m$  vectors  $x^1, \dots, x^m \in \mathfrak{R}^n$  at random by choosing each coordinate randomly from  $\{\sqrt{\frac{1+\epsilon}{m}}, -\sqrt{\frac{1+\epsilon}{m}}\}$ . Then consider the mapping from  $\mathfrak{R}^n$  to  $\mathfrak{R}^m$  given by

$$z \longrightarrow (z \cdot x^1, z \cdot x^2, \dots, z \cdot x^m).$$

In other words  $u^i = (z^i \cdot x^1, z^i \cdot x^2, \dots, z^i \cdot x^m)$  for  $i = 1, \dots, k$ . We want to show that with positive probability,  $u^1, \dots, u^k$  has the desired properties. This would mean that there exists at least one choice of  $u^1, \dots, u^k$  satisfying inequality 2. To show this, first we write the expression  $\|u^i - u^j\|$  explicitly:

$$\|u^i - u^j\|^2 = \sum_{k=1}^m \left( \sum_{l=1}^n (z_l^i - z_l^j) x_l^k \right)^2.$$

Denote by  $z$  the vector  $z^i - z^j$ , and by  $u$  the vector  $u^i - u^j$ . So we get:

$$\|u\|^2 = \|u^i - u^j\|^2 = \sum_{k=1}^m \left( \sum_{l=1}^n z_l x_l^k \right)^2.$$

Let  $X_k$  be the random variable  $(\sum_{l=1}^n z_l x_l^k)^2$ . Its expectation is  $\mu = \frac{1+\epsilon}{m} \|z\|^2$  (can be seen similarly to the proof of lemma 1). Therefore, the expectation of  $\|u\|^2$  is  $(1 + \epsilon) \|z\|^2$ . If we show that  $\|u\|^2$  is concentrated enough around its mean, then it would prove the theorem. More formally, we state the following Chernoff bound lemma:

LEMMA 5

There exist constants  $c_1 > 0$  and  $c_2 > 0$  such that:

1.  $Pr[\|u\|^2 > (1 + \beta)\mu] < e^{-c_1\beta^2m}$
2.  $Pr[\|u\|^2 < (1 - \beta)\mu] < e^{-c_2\beta^2m}$

Therefore there is a constant  $c$  such that the probability of a "bad" case is bounded by:

$$Pr[(\|u\|^2 > (1 + \beta)\mu) \vee (\|u\|^2 < (1 - \beta)\mu)] < e^{-c\beta^2m}$$

Now, we have  $\binom{n}{2}$  random variables of the type  $\|u_i - u_j\|^2$ . Choose  $\beta = \frac{\epsilon}{2}$ . Using the union bound, we get that the probability that any of these random variables is not within  $(1 \pm \frac{\epsilon}{2})$  of their expected value is bounded by

$$\binom{n}{2} e^{-c\frac{\epsilon^2}{4}m}.$$

So if we choose  $m > \frac{8(\log n + \log c)}{\epsilon^2}$ , we get that with positive probability, all the variables are close to their expectation within factor  $(1 \pm \frac{\epsilon}{2})$ . This means that for all  $i, j$ :

$$(1 - \frac{\epsilon}{2})(1 + \epsilon)\|z^i - z^j\|^2 \leq \|u^i - u^j\|^2 \leq (1 + \frac{\epsilon}{2})(1 + \epsilon)\|z^i - z^j\|^2$$

Therefore,

$$\|z_i - z_j\|^2 \leq \|u^i - u^j\|^2 \leq (1 + \epsilon)^2 \|z^i - z^j\|^2,$$

and taking square root:

$$\|z^i - z^j\| \leq \|u^i - u^j\| \leq (1 + \epsilon)\|z^i - z^j\|,$$

as required.

It remains to prove lemma 4. We prove the first part. Let  $\alpha^2 = \frac{1+\epsilon}{m}$ , so  $\mu = \alpha^2 m \|z\|^2$  and we get the following equation:

$$\begin{aligned} P := Pr[\|u\|^2 > (1 + \beta)\mu] &= Pr[\|u\|^2 > (1 + \beta)\alpha^2 m \|z\|^2] \\ &= Pr[\|u\|^2 - (1 + \beta)\alpha^2 m \|z\|^2 > 0] \\ &= Pr[t(\|u\|^2 - (1 + \beta)\alpha^2 m \|z\|^2) > 0] \quad \forall t > 0 \\ &= Pr[\exp(t(\|u\|^2 - (1 + \beta)\alpha^2 m \|z\|^2)) > 1] \\ &\leq E[\exp(t(\|u\|^2 - (1 + \beta)\alpha^2 m \|z\|^2))] \quad (\text{Markov}) \end{aligned} \tag{3}$$

We calculate the last expectation:

$$\begin{aligned} P &\leq E[\exp(t(\|u\|^2))] \exp(-t(1 + \beta)\alpha^2 m \|z\|^2) \quad (\text{constant goes out}) \\ &= E[\exp(t(\sum_{k=1}^m (\sum_{l=1}^n z_l x_l^k)^2))] \exp(-t(1 + \beta)\alpha^2 m \|z\|^2) \\ &= E[\exp(t(\sum_k (\sum_l z_l^2 (x_l^k)^2)) + t(\sum_k (\sum_{l \neq h} z_l z_h x_l^k x_h^k)))] \exp(-t(1 + \beta)\alpha^2 m \|z\|^2) \\ &= E[\exp(t\alpha^2 m \|z\|^2 + t(\sum_k (\sum_{l \neq h} z_l z_h x_l^k x_h^k)))] \exp(-t(1 + \beta)\alpha^2 m \|z\|^2) \end{aligned} \tag{4}$$

The last step used the fact that  $(x_l^k)^2 = \alpha^2$  and  $\sum z_l^2 = \|z\|^2$ . So continuing, we get:

$$P \leq E[\exp(t(\sum_k (\sum_{l \neq h} z_l z_h x_l^k x_h^k)))] \exp(-t\beta\alpha^2 m \|z\|^2) \quad (5)$$

The set of variables  $\{x_l^k x_h^k\}_{l \neq h}$  are pairwise independent. Therefore the above expectation can be rewritten as a product of expectations:

$$P \leq \left( \prod_k \prod_{l \neq h} E[\exp(tz_l z_h x_l^k x_h^k)] \right) \exp(-t\beta\alpha^2 m \|z\|^2) \quad (6)$$

we notice that

$$E[\exp(tz_l z_h x_l^k x_h^k)] = \frac{1}{2} \exp(tz_l z_h \alpha^2) + \frac{1}{2} \exp(-tz_l z_h \alpha^2) < \exp(t^2 z_l^2 z_h^2 \alpha^4)$$

(the last inequality is easily obtained by Taylor expanding the exponent function). Plugging that in (6), we get:

$$\begin{aligned} P &< \left( \prod_k \prod_{l \neq h} \exp(t^2 z_l^2 z_h^2 \alpha^4) \right) \exp(-t\beta\alpha^2 m \|z\|^2) \\ &= \left( \prod_{l \neq h} \exp(t^2 z_l^2 z_h^2 \alpha^4) \right)^m \exp(-t\beta\alpha^2 m \|z\|^2) \\ &= \exp(mt^2 \sum_{l \neq h} z_l^2 z_h^2 \alpha^4 - t\beta\alpha^2 m \|z\|^2) \end{aligned} \quad (7)$$

Using simple analysis of quadratic function we see that the last expression obtains its minimum when

$$t = \frac{\beta \|z\|^2}{2\alpha^2 \sum_{l \neq h} z_l^2 z_h^2}.$$

Substituting for  $t$ , we get:

$$P < \exp\left(-\beta^2 m \frac{\|z\|^4}{4 \sum_{l \neq h} z_l^2 z_h^2}\right) \quad (8)$$

Finally, the expression

$$\delta(z) = \left( \frac{\|z\|^4}{4 \sum_{l \neq h} z_l^2 z_h^2} \right)$$

is bounded below by a constant  $c_1$ . To prove this, first note that  $\delta(z) = \delta(\gamma z)$  for any  $\gamma \neq 0$ . So it is enough to consider the case  $\|z\| = 1$ . Then, using Lagrange multipliers technique, for example, we get that  $\delta(z)$  obtains its minimum when  $z_l = \frac{1}{\sqrt{n}}$  for each  $l = 1..n$ . Plugging this in the expression for  $\delta(z)$  we see that it is bounded above by a constant  $c_1$  that does not depend on  $n$ . This completes the proof.  $\square$