

- Now you have to deal with the possibility of your context switch being interrupted
 - Timer Interrupt (10 ms)
 - Page Fault
 - TLB Miss

Implement locking primitives (thread.c)

- Lock
 - Acquire & Release
- Condition
 - Signal, Wait, Broadcast
- Semaphore
 - Up & Down
- Barrier (Extra Credit)
 - Wait
- Timer Interrupt (entry.S)
 - Also handle irq7 call fake_irq7 instead of scheduler_entry



Condition Variables

- Wait operation cause the current thread to block on put itself of a wait list
- Signal causes a random thread on the wait list to start running
- Broadcast wakes up all waiting threads
- Condition Variables can "lose" signals if no thread is on the wait list to receive it

Semaphore

- Use when you need a way to store wakeup signals
 - Producer-Consumer Problem
- Two operations
 - Up -> Called to wakeup a thread, if no thread is ready to be woken up, the wakeup signal is stored
 - Down -> Called to see if a wakeup signal is stored, otherwise sleep

Mutex

- A semaphore used for mutual exclusion
 - Semaphore should only ever have values 0 and 1
- Useful when only one thread should execute a piece of code at a time

Barrier

- Used to synchronize multiple threads at a single point
 - Usually the boundary between passes of phases of an algorithm
 - Count number of threads stopped
 - Count number of threads started

Scheduling

- With preemptive kernel we can implement scheduling algorithms
 - Round-Robin (default)
 - Priority (Extra Credit)
 - Priority w/Queue
 - Random

Handling Interrupts

- Handle the timer interrupt (irq0) & irq7 plus system call interrupts (traps)
 - For irq7 call fake_irq7
 - For all interrupts
 - Save the current state
 - Interrupt controller will suspend further interrupts when generating a non-system interrupt
 - Tell the Interrupt Controller to allow more interrupts
 - Process the interrupt
 - Call scheduler to move to next process on timer interrupt

9

10

